

## THE INTEL EXASCIENCE PROJECT

#### HIGH-PERFORMANCE COMPUTING IN THE NEXT DECADE

**ROEL WUYTS EXASCIENCE LAB, IMEC** 



About Me

About Intel ExaScience Lab

About the Runtime Layer





#### Studies: Licentiaat Informatica (VUB, 1991-1995)

1995	2001	2004	07	08	09	10	11
Doctoral Researcher VUB	Postdoc University of Bern, Switzerland	Chargé de cours ULB	Princi imec	Principal Scientist imec			
				Prof KUL	fessor -	· (10%)	



# **JUGGLING HATS**

#### **IMEC**

- Embedded devices
- Runtime resource management
- Intel ExaScience Project

ULB

• AOP

Object versioning

KUL Language Design



# **ABOUT IMEC**

Research organization located in Leuven

- world-leading independent research center in nanoelectronics and nanotechnology
- More Moore research targets semiconductor scaling for the 22nm technology node and beyond.
- More than Moore research invents technology for nomadic embedded systems, wireless autonomous transducer solutions, biomedical electronics, photovoltaics, organic electronics and GaN power electronics.

#### Numbers

- Budget: ± 280 M€
- ▶ Staff: ± 1800
- Cleanroom: ± 10,000 m





# **EXASCIENCE LAB: NOVEL RESEARCH STRUCTURE**



# **ExaScience Lab** Intel Labs Europe





Vlaamse overheid 🔵





imec









# TOWARDS EXASCALE SUPERCOMPUTERS

SuperComputer ≈ 2 PetaFlops ≈ 100.000 PC's @ 20GigaFlops/PC

 $\approx$  1000 \* PetaScale  $\approx$  50.000.000 PC's of today

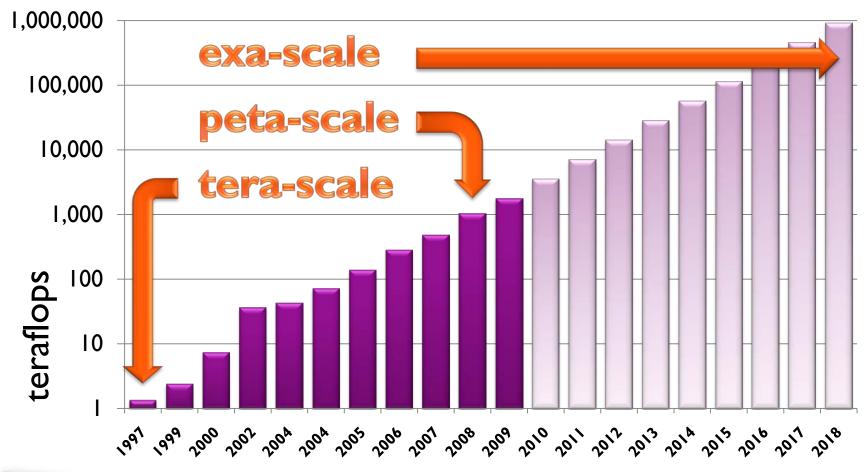
1997	Intel ASCI Red/9152	I
1999	Intel ASCI Red/9632	2
2000	IBM ASCI White	7
2002	NEC Earth Simulator	36
2004	SGI Project Columbia	43
2004	IBM Blue Gene/L	71
2005		137
2006		281
2007		478
2008	IBM Roadrunner	1.026
2009	Cray XT4/XT5 Jaguar	1.759
2010	Tianhe-I	2.5

Kilo I0 <sup>3</sup>	1. 000
Mega 106	1.000.000
Giga 109	1.000.000.000
Tera 1012	1.000.000.000.000
Peta 1015	1.000.000.000.000.000
Exa 1018	1.000.000.000.000.000.000



ExaScale

# TOWARDS EXASCALE SUPERCOMPUTERS





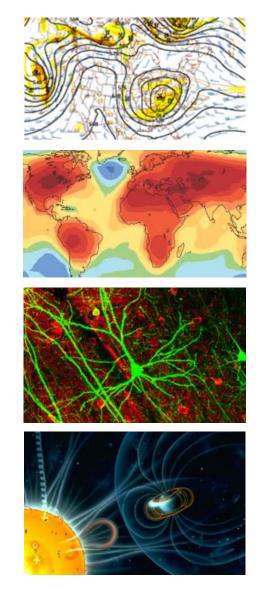
# WHY DO WE NEED THEM ?

Computer Simulations are needed to fundamentally *understand* phenomena and to accurately predict.

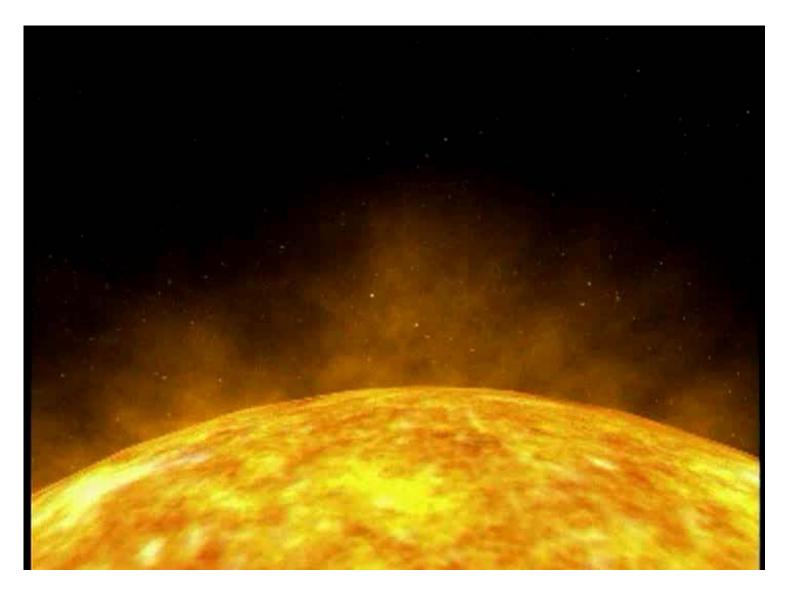
Simulation examples:

- Neural Networks
- Climate Models
- Economical Models
- Weather Prediction
- Space Weather Prediction

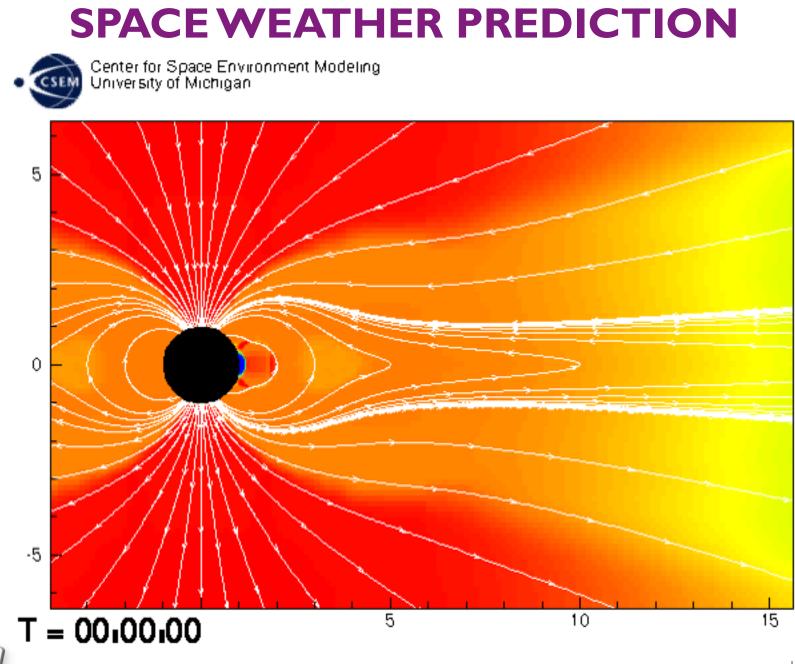
These are all limited by compute power



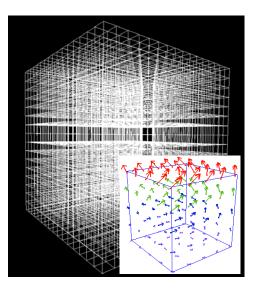




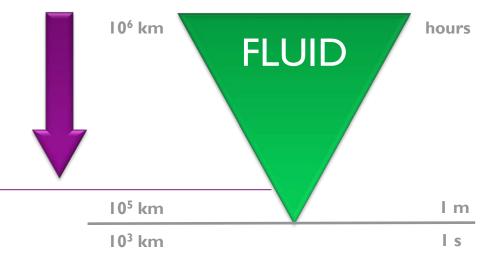






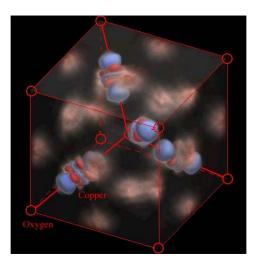


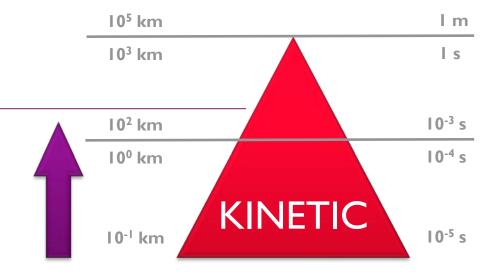
#### Petascale



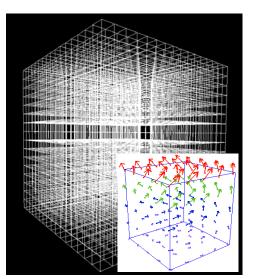


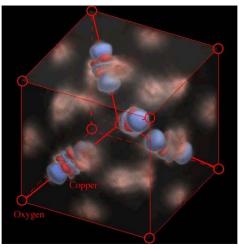
## Petascale



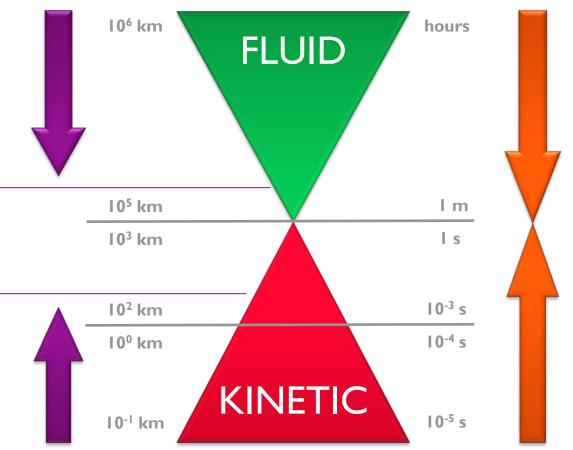








# Petascale



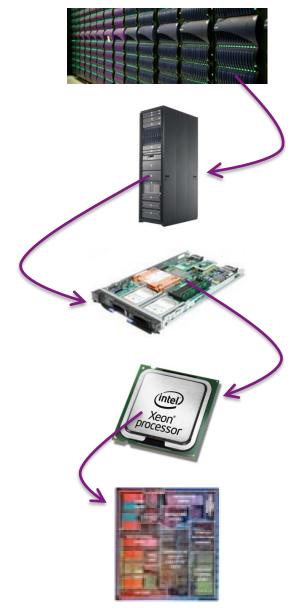


**Exascale** 

#### Applications: computationally intensive, billions of threads



# **INSIDE A HPC CENTER**



Datacenter: 10<sup>9</sup> threads

Rack: 10<sup>4</sup>-10<sup>5</sup> threads

Socket/blade: 500-5000 threads

Die: 100-1000 threads

Core/tile: 1-10 threads





energy, energy, energy





energy, energy, energy



heat



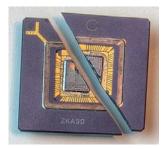




heat

hardware failures

energy, energy, energy





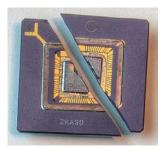




energy, energy, energy

heat

hardware failures





ExaScience Lab

programming massively parallel machines (a million cores, a billion threads!)

#### Applications: computationally intensive, billions of threads



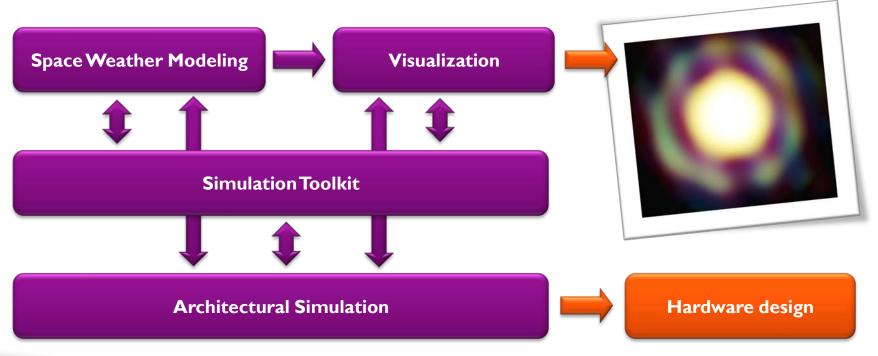
# **ExaScience Lab** Intel Labs Europe



# **RESEARCH TOPICS**

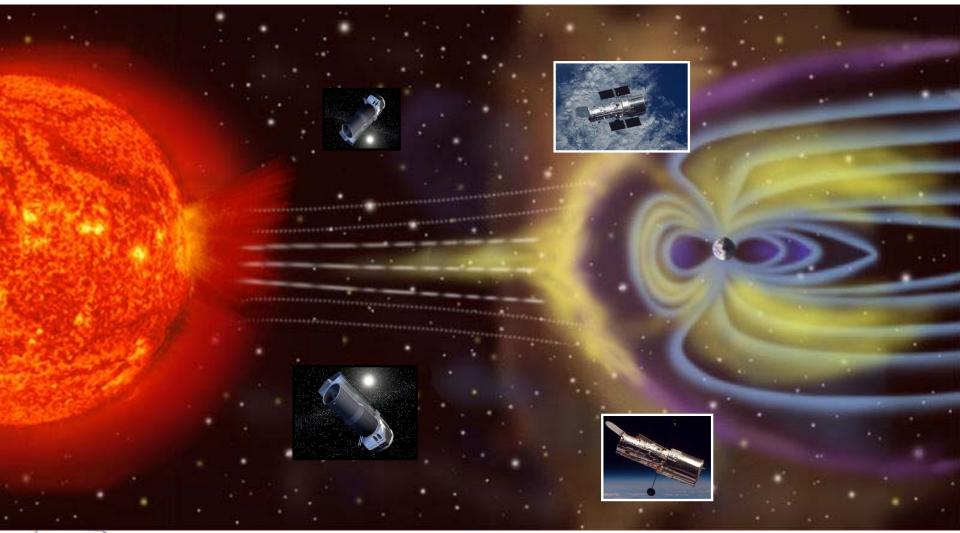
the ExaScience Lab will advance state of the art:

- More accurate space weather modeling and simulation
- Extremely scalable, fault tolerant simulation toolkit
- In-situ visualization through virtual telescopes
- Architectural simulation of large-scale systems and workloads



# VISUALIZATION

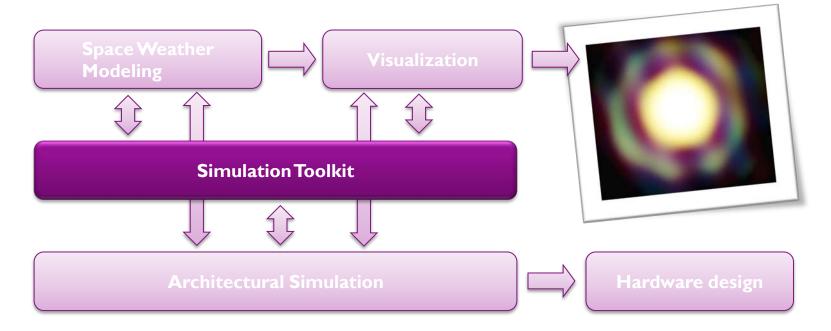
#### longer term objective = Virtual Telescopes





# **SIMULATION TOOLKIT**

#### **HOW TO PROGRAM AN EXASCALE SYSTEM ?**





# **SIMULATION TOOLKIT**

Isn't this just our current software programs x 1.000.000?

**Extreme Parallelism** 

- Optimized implementations of 'ultra' scalable numerical kernels
- Heterogeneous Architecture
- Support for exascale parallel programming models

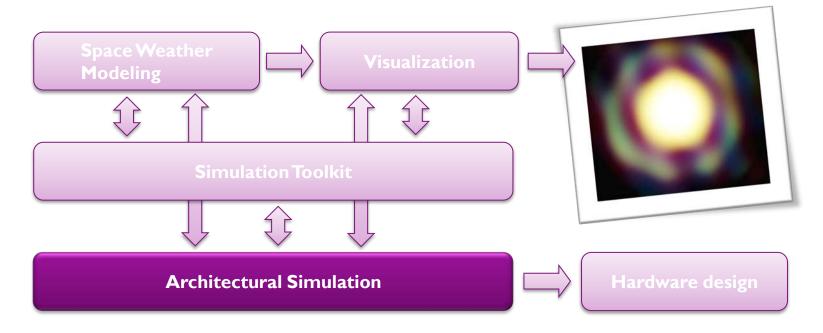
Dealing with hardware failures at the software level

- Dynamic load balancing
- Guarantee fault-tolerance



# **ARCHITECTURAL SIMULATION**

#### HOW TO SIMULATE AN EXASCALE SYSTEM?





# **ARCHITECTURAL SIMULATION**

Before running a program on an exascale system, developers need to be able to predict:

- Performance
- Power
- Reliability
- Resource utilization

Usually this happens by simulating the system on a more powerful computer system.

But this is going to be the biggest system on earth.

#### So, how will we be able to simulate this system?



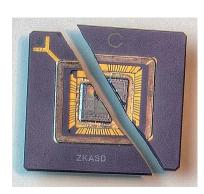
# THE RUNTIME LAYER

Software problems to tackle:

- Minimize energy consumption

   ... while remaining performant
- Deal with hardware variability and failures ... transparently for the application
- Program massively parallel systems ... without burdening the developer

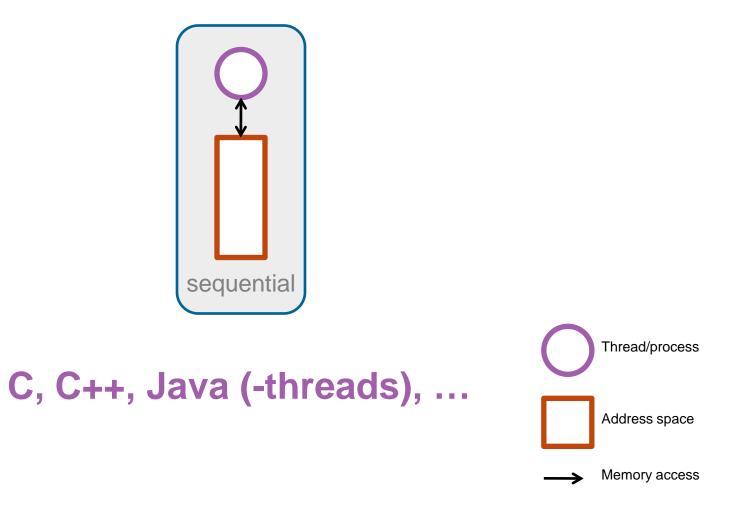








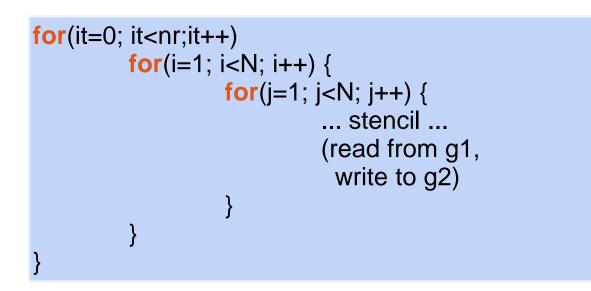
# **PROGRAMMING MODELS**

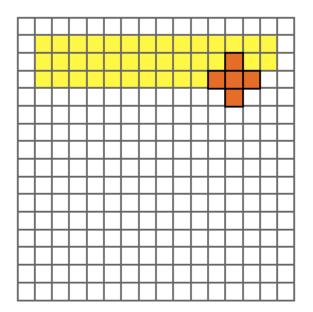




# SEQUENTIAL PROGRAMMING EXAMPLE

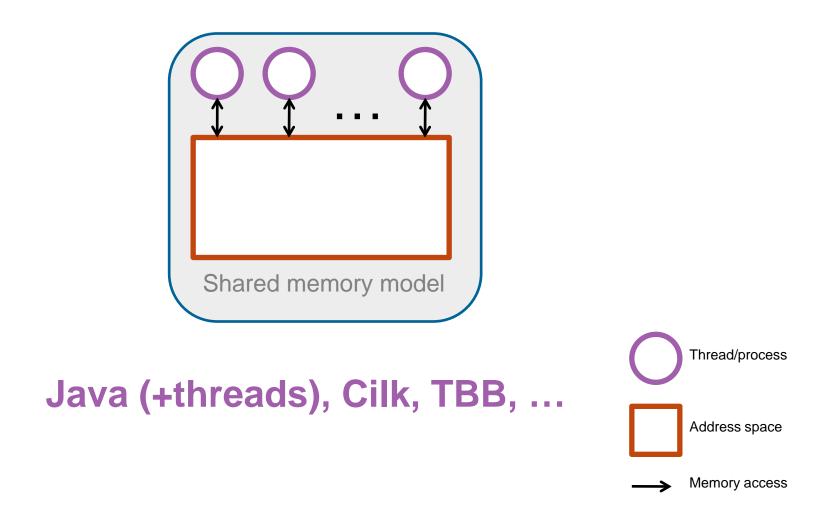
# Explicit 2-dimensional heat distribution simulation







# **PROGRAMMING MODELS**





# SHARED MEMORY PROGRAMMING: CILK

- New parallel function calling mechanism:
  - cilk\_spawn indicates a call to function that can proceed in parallel with caller
  - cilk\_sync awaits finish of spawned children
- work-stealing scheduler

#include <cilk/cilk.h>

```
...

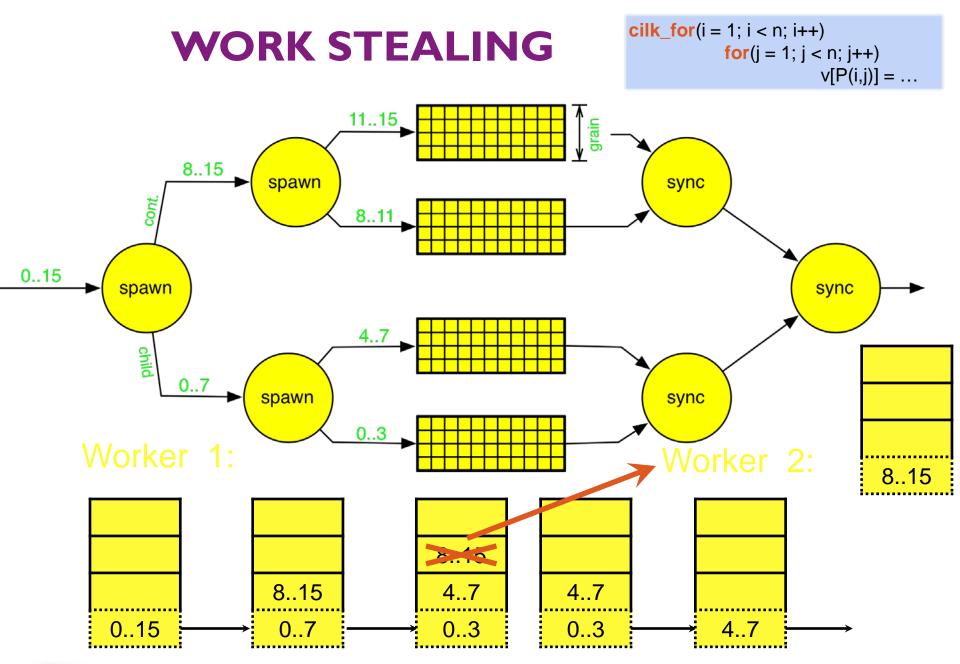
int i, j;

cilk_for(i = 1; i < n; i++)

for(j = 1; j < n; j++)

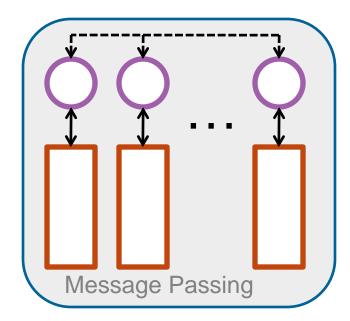
v[P(i,j)] = ...
```



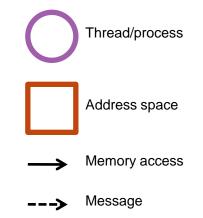




# **PROGRAMMING MODELS**



# MPI, actor models, ...





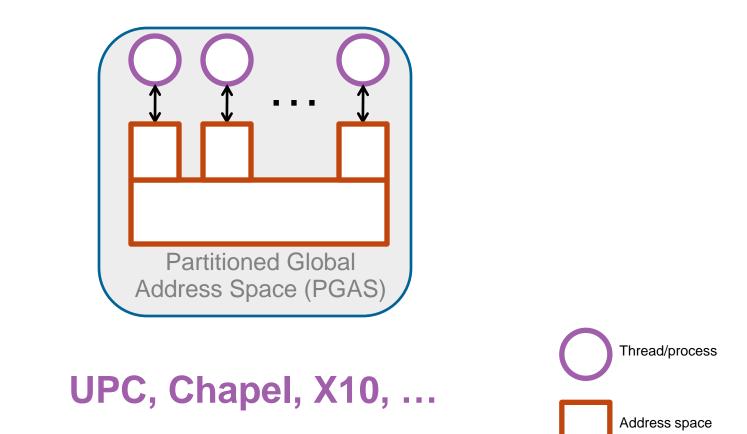
# **ACTORS IN SCALA**

- Scala blends object-oriented and functional programming
  - Pure OO
  - Classes and inheritance
  - Block closures
  - Type inference
  - Higher-order functions
  - Sophisticated static type system

```
class Worker (signal : MailBox)
            extends Runnable
  var func = () => ();
  def run ()
     var goOn = true;
     while (goOn) {
      signal.receive
        case Go() => func();
        case Stop() \Rightarrow goOn = false;
       };
      signal.send(Done());
```



#### **PROGRAMMING MODELS**





Memory access

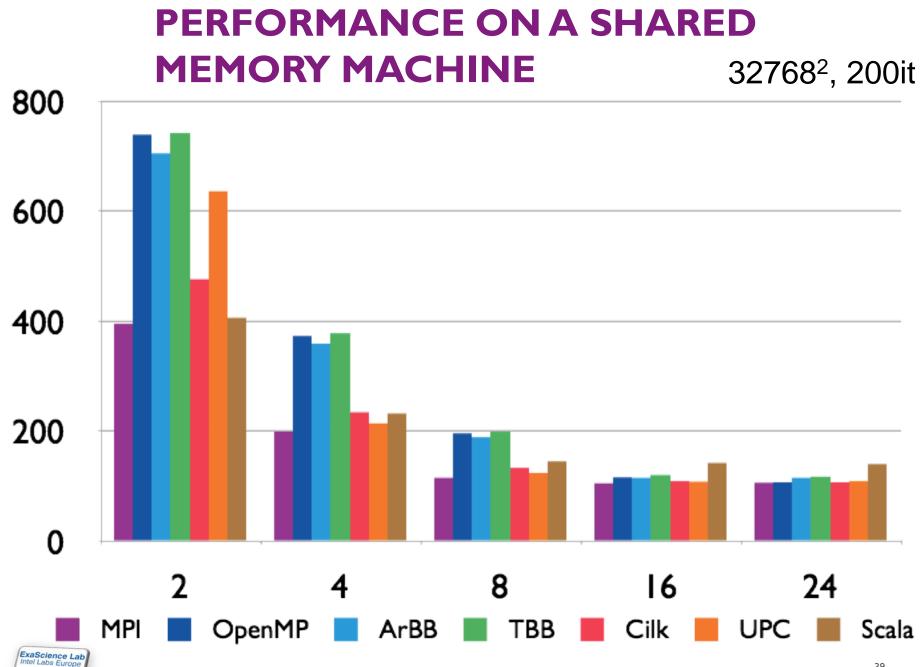
Message

#### UPC (UNIFIED PARALLEL C), A PGAS LANGUAGE

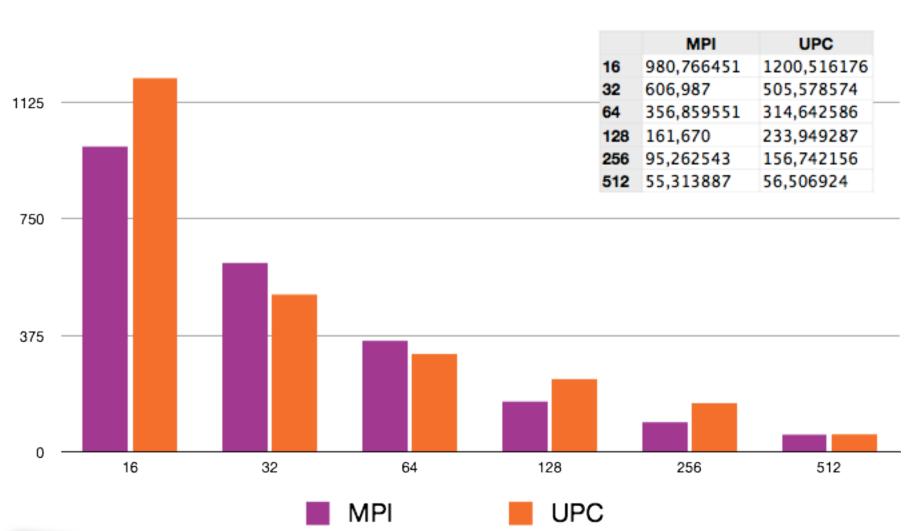
. . .

- Threads working in SPMD fashion
  - MYTHREAD specifies thread index
  - THREADS specifies number of threads (workers)
- shared keyword indicates shared scalars or arrays
  - Shared data has affinity to one thread
- Synchronization when needed (barriers, locks)





# PERFORMANCE ONSUPERCOMPUTERVic3, 327682, 2000it





#### **LESSONS LEARNED**

UPC performs quite well

- But SPMD will not work on ExaScale machines
- Impact of memory affinity is huge
- Work stealing shows promise
- Achieves automatic runtime load balancing with reasonable performance

Conclusion:

we want memory affinity aware work stealing



#### REACTIVE WORK REBALANCING EXPERIMENT ON CLUSTER

Start by giving each thread the same number of grid rows to process

Measure how long it takes each UPC thread to process these rows

If necessary: React by transferring rows to other UPC thread



#### REACTIVE WORK REBALANCING: KERNEL IMPLEMENTATION IN UPC

#### •••

. . .

for (iteration = start; iter < start+nr; iteration ++) {
 upc\_barrier;</pre>

#### //potentially rebalance rows between threads

rebalance(mgr, iteration); upc\_barrier; adjustWork(mgr, \*old, \*new, iteration); incrementIterationCounter(mgr);

## //do a heat step and measure how long it takes upc\_barrier; timeTaken = (\*step)(mgr, iteration, \*old, \*new);

//update the internal information
update(mgr, iteration, timeTaken);



#### REACTIVE WORK REBALANCING: KERNEL IMPLEMENTATION IN UPC

#### • • •

. . .

for (iteration = start; iter < start+nr; iteration ++) {
 upc\_barrier;</pre>

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## //do a heat step and measure how long it takes upc\_barrier; timeTaken = (\*step)(mgr, iteration, \*old, \*new);

//update the internal information
update(mgr, iteration, timeTaken);



#### REACTIVE WORK REBALANCING: ACTION!

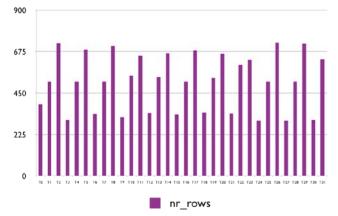
Heat distribution simulation (16384x16384, 2000 it)

Every third thread (0, 3, ...) straggles

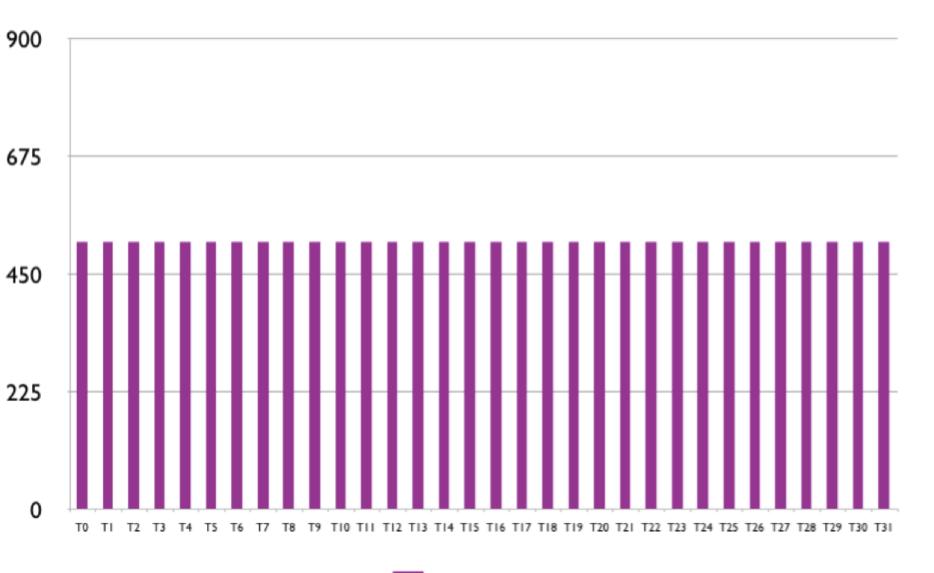
slows down with a factor of 2
 Ran distributed on 32 cores

X axis: core number

Y axis: nr of rows processed

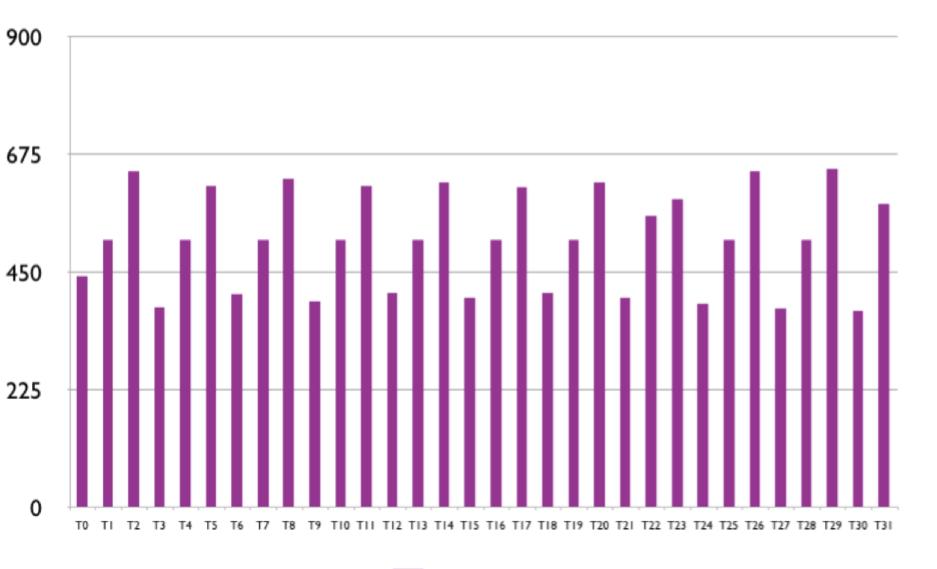






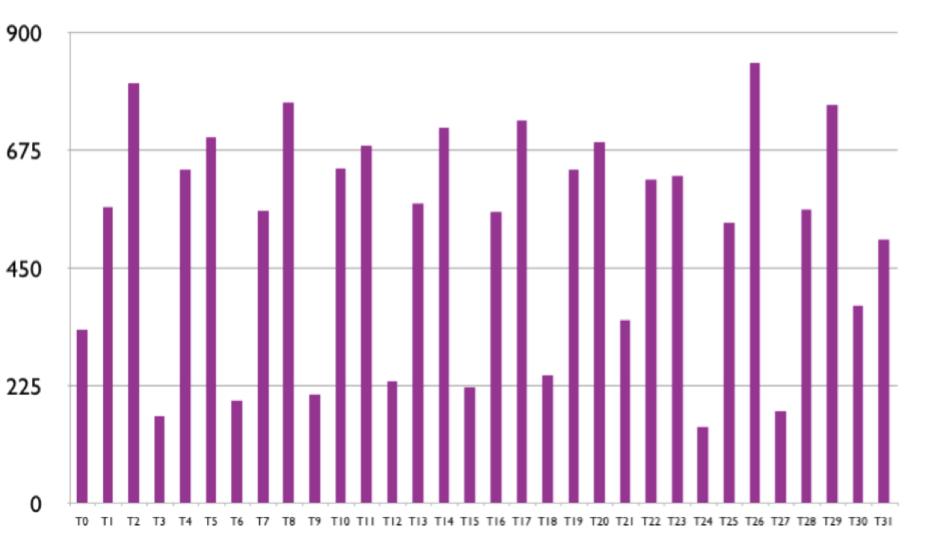


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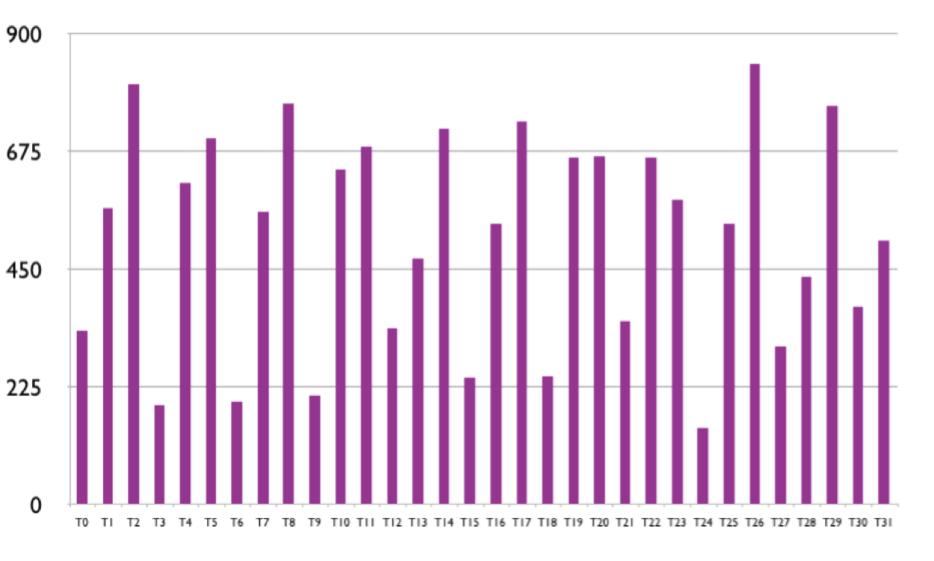


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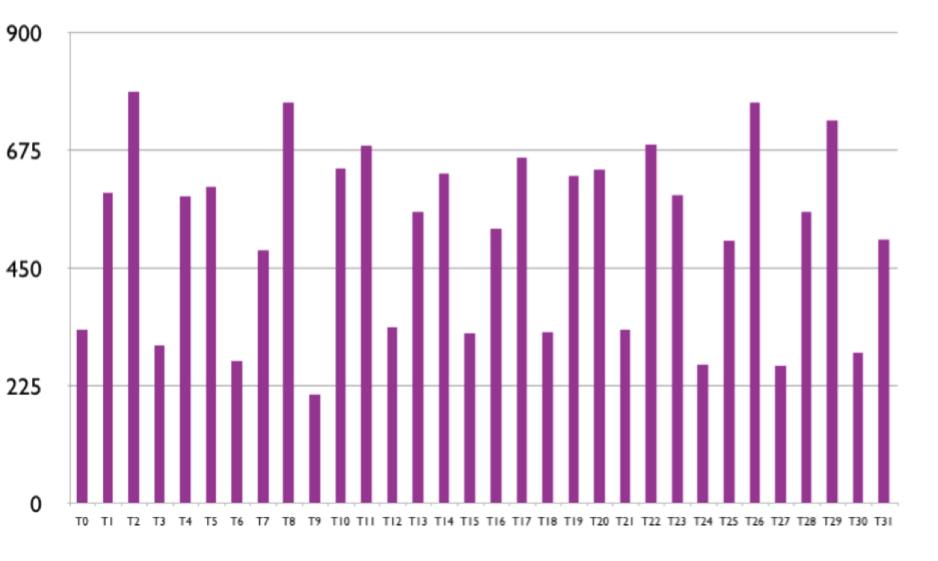


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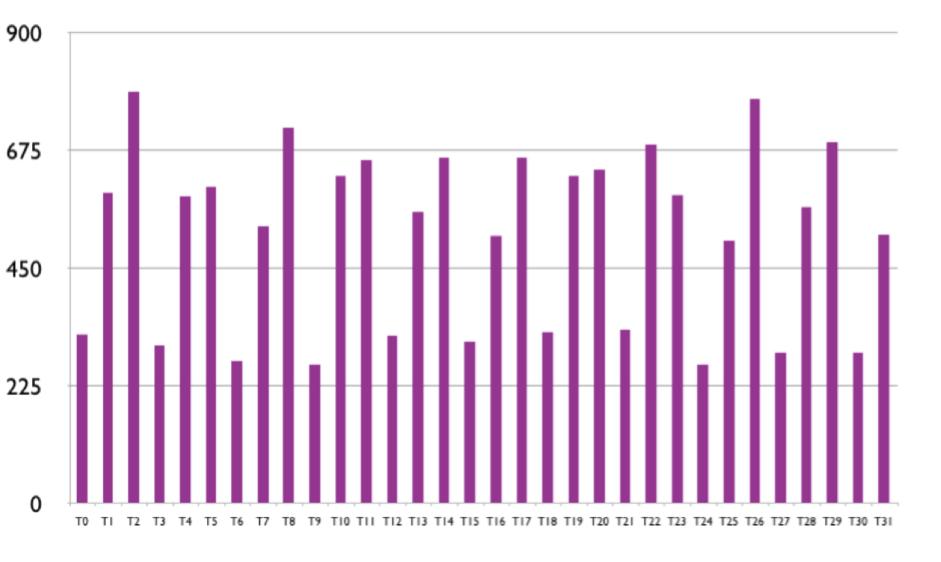


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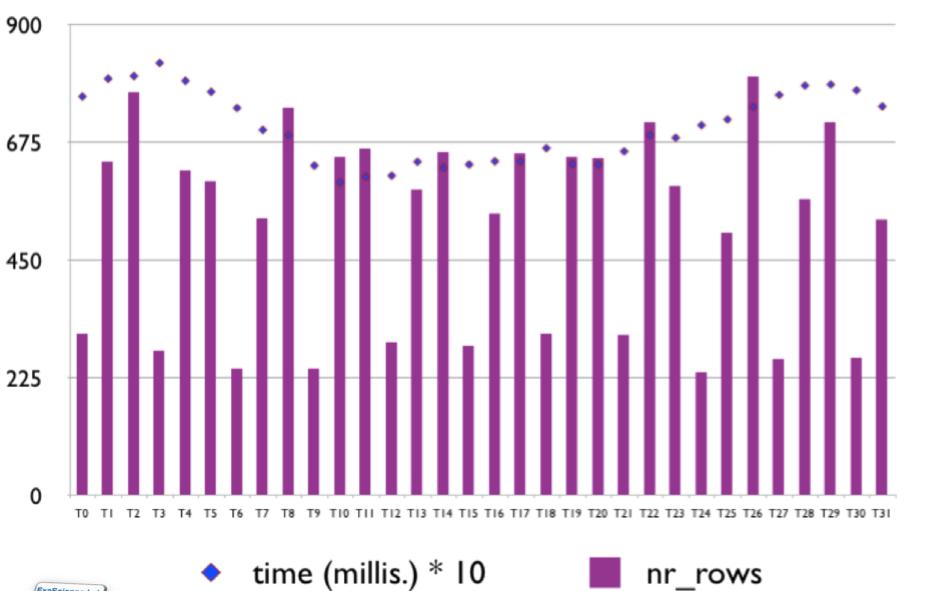


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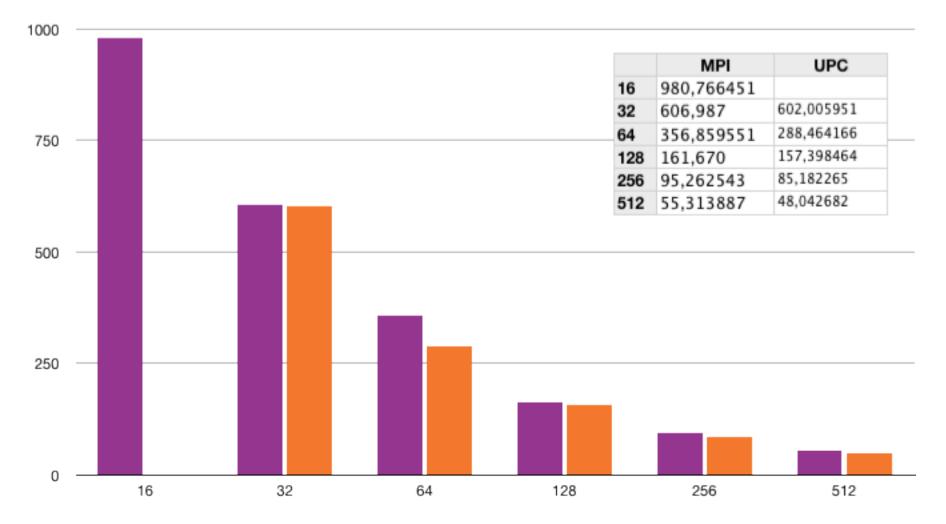
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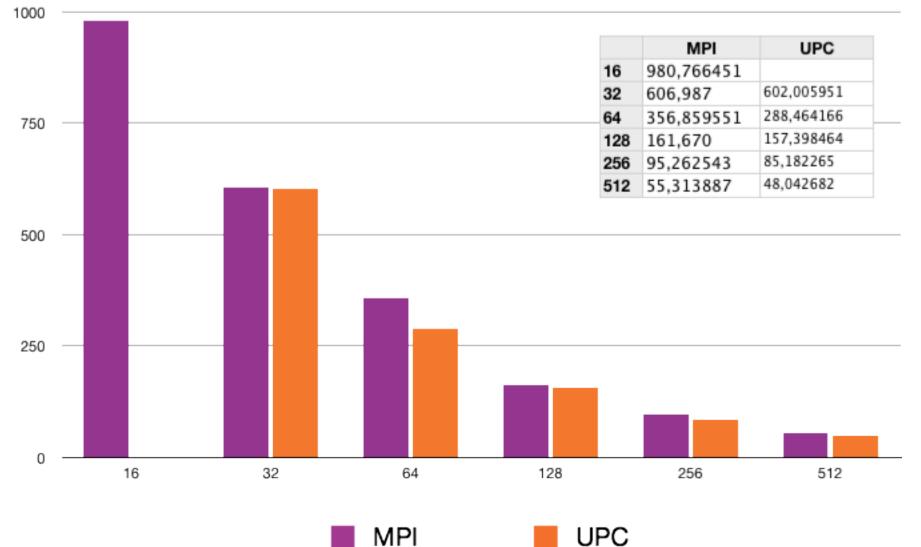
# PERFORMANCE ONSUPERCOMPUTERVic3, 327682, 2000it





#### PERFORMANCE ON SUPERCOMPUTER

#### Vic3, 32768<sup>2</sup>, 2000it





#### Applications: computationally intensive, billions of threads

#### **ExaScience Lab** Intel Labs Europe

Hardware: millions of cores, runtime variability and failures, energy

## www.exascience.com

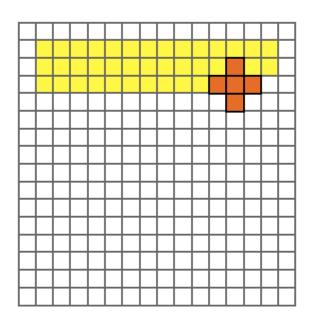
# **ExaScience Lab** Intel Labs Europe







- Grid
- Stencil Operation
- Explicit Timestepping







- New parallel function calling mechanism:
  - cilk\_spawn indicates a call to function that can proceed in parallel with caller
  - cilk\_sync awaits finish of spawned children
- work-stealing scheduler

#include <cilk/cilk.h>

```
...

int i, j;

cilk_for(i = 1; i < n; i++)

for(j = 1; j < n; j++)

v[P(i,j)] = ...
```



## TBB (THREADING BUILDING BLOCKS)

- C++ Threading library (Intel)
- Task based parallelism
- Task scheduler
- Work stealing approach similar to Cilk
- Automatic parallellization
  - parallel\_for
  - reduce/scan/sort/while
- Parallel data structures
- Scalable memory allocation
- Mutual exclusion
- Atomic operations



## ARBB (ARRAY BUILDING BLOCKS)

#include <arbb.hpp>

- C++ language extension for vector parallel programming
- Intel product
  - currently beta 2
- Computational kernel defined for the individual grid points
- Automatic parallelization over cores and SIMD units
- Can be extended to GPU
- Vector code is JIT-compiled

```
template<typename T>
void heat_stencil(T src, T& dst) {
    dst = src + D * (-4*src
                 + neighbor(src, -1, 0)
                 + neighbor(src, 1, 0)
                 + neighbor(src, 0, -1)
                 + neighbor(src, 0, 1));
template<typename T>
void heat_driver(dense<T, 2> grid,
                 dense<T, 2> next) {
   map(heat stencil<T>)(grid, next);
int main(){
    call(heat_driver<T>)(grid, next);
```





- Threads working in SPMD fashion
  - MYTHREAD specifies thread index
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- shared keyword indicates shared scalars or arrays
  - Shared data has affinity to one thread
- Synchronization when needed (barriers, locks)

int i, j;
for(i = 1; i < n; i++)
 upc\_forall(j = 1; j < n; j++; &v[i][j])
 v[i][j] = stencil(nu, u[i][j],
 u[i][j+1],
 u[i+1][j]);
...</pre>





#### "Global View" paradigm clean algorithm code

- No communication/data sharing code intermixed
- No explicit decomposition of data structures and control flow into per-task or per "node" chunks
- Data parallel features based on domains (index set)
- Data distribution customizable by introducing domain maps
  - separate from algorithm code
- Explicit Task parallellism

const MatrixSpace :

**domain**(2) = [0..size, 0..size];

const ProblemSpace:

subdomain(MatrixSpace) = [1..n, 1..n];

```
def heat(n, u, v, nu) {
  forall (i,j) in ProblemSpace do
    v(i,j) = u(i,j) + nu * ( u(i+1,j) ...
```



### **XI0**

- Type-safe parallel OO language
- Asynchronous PGAS
  - data locality through places
  - lightweight activities

```
public class HeatTransfer_v1 {
  const BigD =
    Dist.makeBlock([0..n+1, 0..n+1], 0);
  const D = BigD | ([1..n, 1..n] as Region);
  const A = DistArray.make[Real]
        (BigD,(p:Point) =>{ ...init...});
...
def run() {
```

```
tet run() {
  var iter:Int = 0;
  do {
    iter = iter + 1;
    finish ateach (p in D) Temp(p) = stencil(p);
    finish ateach (p in D) A(p) = Temp(p);
  } while (iter < iterations);</pre>
```



#### **SCALA**

- Blends object-oriented and functional programming
  - Pure OO
  - Classes and inheritance
  - Block closures
  - Type inference
  - Higher-order functions
  - Sophisticated static type system

```
class Worker (signal : MailBox)
            extends Runnable
  var func = () => ();
  def run ()
     var goOn = true;
     while (goOn) {
      signal.receive
        case Go() => func();
        case Stop() \Rightarrow goOn = false;
       };
      signal.send(Done());
```



#### HPC VERSUS...

#### Grid/Cloud Computing

- loosely coupled heterogeneous systems that are geographically dispersed
- Computing resources are not administered centrally
- Interconnected through slower networks
- Typically runs several different programs in parallel

