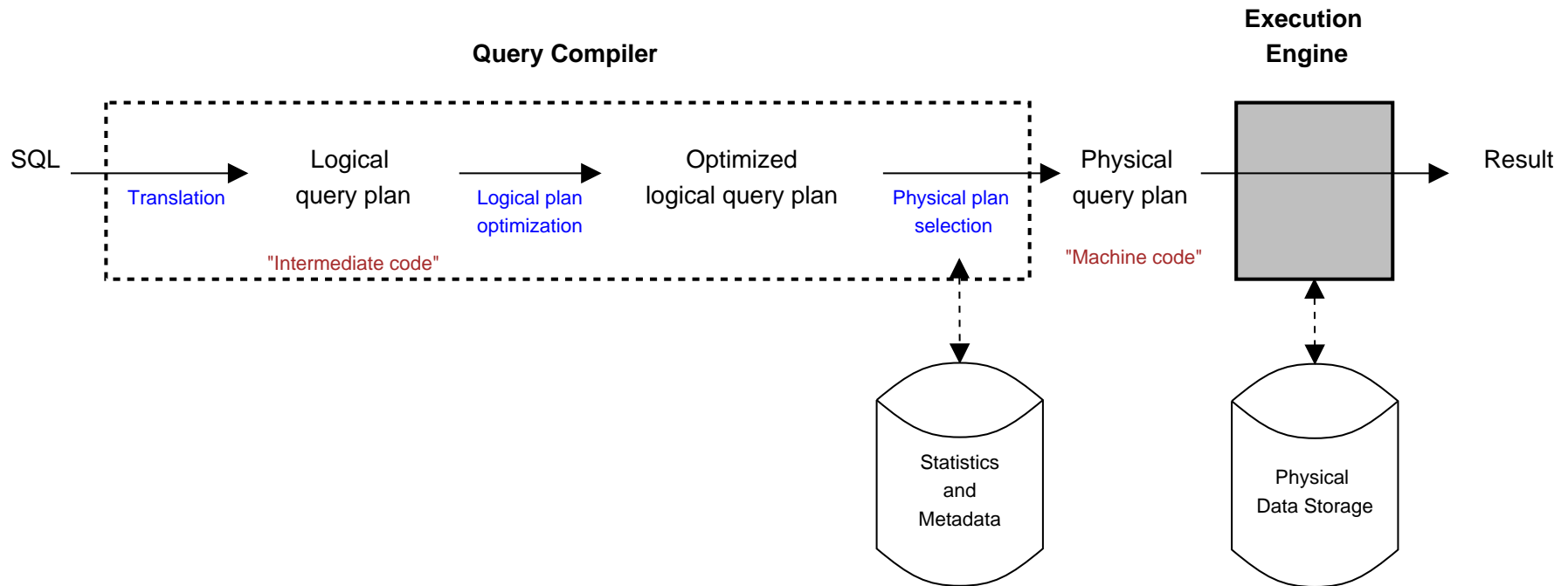


# Optimization of logical query plans

## Eliminating redundant joins

# Optimization of logical query plans



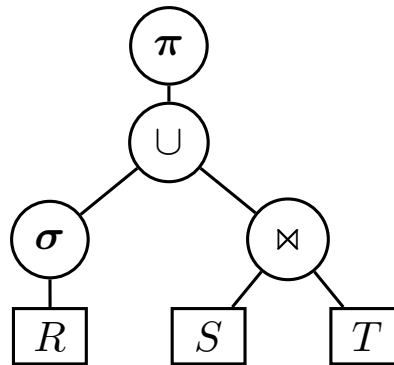
# Optimization of logical query plans

**Logical plans = execution trees**

A logical query plan like

$$\pi_{A,B}(\sigma_{A=5}(R) \cup (S \bowtie T))$$

**is essentially an execution tree**



**A physical query plan is a logical query plan in which each node is assigned the algorithm that should be used to evaluate the corresponding relational algebra operator. (There are multiple ways to evaluate each algebra operator)**

# Optimization of logical query plans

## The need to optimize

- Every internal node (i.e., each occurrence of an operator) in the plan must be executed.
- Hence, the fewer nodes we have, the faster the execution

## Definition

A relational algebra expression  $e$  is **optimal** if there is no other expression  $e'$  that is both (1) equivalent to  $e$  and (2) “shorter” than  $e$  (i.e., has fewer occurrence of operators than  $e$ ).

## The optimization problem:

**Input:** a relational algebra expression  $e$

**Output:** an optimal relational algebra expression  $e'$  equivalent to  $e$ .

# Optimization of logical query plans

## The optimization problem is undecidable:

It is known that the following problem is undecidable:

**Input:** relational algebra expressions  $e_1$  and  $e_2$  over a single relation  $R(A, B)$

**Output:**  $e_1 \equiv e_2$ ?

**Proof:** Suppose that we can compute  $e'$ , the optimal expression for  $(e_1 - e_2) \cup (e_2 - e_1)$ . Note that  $e_1 \equiv e_2$  if, and only if,  $e'$  is either  $\sigma_{\text{false}}(R)$  or  $R - R$ .

**Conclusion:** the optimization problem is undecidable

**Question:** can we optimize plans that are of a particular form?

## Optimization of select-project-join expressions

In practice, most SQL queries are of the following form:

```
SELECT ...  
FROM R1, R2, ..., Rm  
WHERE A1 = B1 AND A2 = B2 AND ... AND An = Bn
```

The corresponding logical query plans are of the form:

$$\pi_{\dots} \sigma_{A_1=B_1 \wedge A_2=B_2 \wedge \dots \wedge A_n=B_n} (R_1 \times R_2 \times \dots \times R_m).$$

We call such relational algebra expressions **select-project-join expressions (SPJ expressions for short)**

# Optimization of select-project-join expressions

## Removing redundant joins:

A careless SQL programmer writes the following query:

```
SELECT movieTitle FROM StarsIn S1
WHERE starName IN (SELECT name
                   FROM MovieStar, StarsIn S2
                   WHERE birthdate = 1960
                   AND S2.movieTitle = S1.movieTitle)
```

This query is equivalent to the following one, which has one join less to execute!

```
SELECT movieTitle FROM StarsIn
WHERE starName IN (SELECT name
                   FROM MovieStar
                   WHERE birthdate = 1960)
```

# Optimization of select-project-join expressions

Here are the corresponding logical query plans:

$$\pi_{S_1.movieTitle}(\rho_{S_2}(\text{StarsIn}) \bowtie_{S_2.movieTitle=S_1.movieTitle} \rho_{S_1}(\text{StarsIn}) \bowtie_{S_1.starName=name} \sigma_{birthdate=1960}(\text{MovieStar}))$$

versus

$$\pi_{S_1.movieTitle}(\rho_{S_1}(\text{StarsIn}) \bowtie_{S_1.starName=name} \sigma_{birthdate=1960}(\text{MovieStar})).$$

Redundant joins may also be introduced because of **view expansion**

Can we optimize select-project-join expressions? (And hence remove redundant joins?)



# Optimization of select-project-join expressions

## Definition

A **conjunctive query** is an expression of the form

$$Q(\underbrace{x_1, \dots, x_n}_{\text{head}}) \leftarrow \underbrace{R(t_1, \dots, t_m), \dots, S(t'_1, \dots, t'_k)}_{\text{body}}$$

Here  $t_1, \dots, t'_k$  denote variables and constants, and  $x_1, \dots, x_n$  must be variables that occur in  $t_1, \dots, t'_k$ . We call an expression like  $R(t_1, \dots, t_m)$  an **atom**. If an atom does not contain any variables, and hence consists solely of constants, then it is called a **fact**.

# Optimization of select-project-join expressions

## Semantics of conjunctive queries

Consider the following toy database  $D$ :

$R$	$S$
1 2	2
2 3	7
2 5	
6 7	
7 5	
5 5	

as well as the following conjunctive query over the relations  $R(A, B)$  and  $S(C)$ :

$$Q(x, y) \leftarrow R(x, y), R(y, 5), S(y).$$

Intuitively,  $Q$  wants to retrieve all pairs of values  $(x, y)$  such that (1) this pair occurs in relation  $R$ ; (2)  $y$  occurs together with the constant 5 in a tuple in  $R$ ; and (3)  $y$  occurs as a value in  $S$ . The formal definition is as follows.

# Optimization of select-project-join expressions

## Semantics of conjunctive queries

Consider the following toy database  $D$ :

$R$	$S$
1 2	2
2 3	7
2 5	
6 7	
7 5	
5 5	

as well as the following conjunctive query over the relations  $R(A, B)$  and  $S(C)$ :

$$Q(x, y) \leftarrow R(x, y), R(y, 5), S(y).$$

A **substitution**  $f$  of  $Q$  into  $D$  is a function that maps variables in  $Q$  to constants in  $D$ . For example:

$$f: \begin{array}{l} x \mapsto 1 \\ y \quad 2 \end{array}$$

# Optimization of select-project-join expressions

## Semantics of conjunctive queries

Consider the following toy database  $D$ :

$R$	$S$
1 2	2
2 3	7
2 5	
6 7	
7 5	
5 5	

as well as the following conjunctive query over the relations  $R(A, B)$  and  $S(C)$ :

$$Q(x, y) \leftarrow R(x, y), R(y, 5), S(y).$$

A **matching** is a substitution that maps the body of  $Q$  into facts in  $D$ . For example:

$$f: \begin{array}{l} x \mapsto 1 \\ y \quad 2 \end{array}$$

# Optimization of select-project-join expressions

## Semantics of conjunctive queries

Consider the following toy database  $D$ :

$R$	$S$
1 2	2
2 3	7
2 5	
6 7	
7 5	
5 5	

as well as the following conjunctive query over the relations  $R(A, B)$  and  $S(C)$ :

$$Q(x, y) \leftarrow R(x, y), R(y, 5), S(y).$$

The **result** of a conjunctive query is obtained by applying all possible matchings to the head of the query. In our example:

$$Q(D) = \{(1, 2), (6, 7)\}.$$

# Optimization of select-project-join expressions

## Translation of SPJ expressions into conjunctive queries

SPJ:  $\pi_{\text{title}}(\text{Movie} \bowtie_{\text{title=movieTitle}} \text{StarsIn} \bowtie_{\text{starName=name}} \sigma_{\text{birthdate=1960}}(\text{MovieStar}))$

CQ:  $Q_1(t) \leftarrow \text{Movie}(t, y, \ell, i, s, p), \text{StarsIn}(t, y_2, n), \text{MovieStar}(n, a, g, 1960)$

## Translation of conjunctive queries into SPJ expressions

CQ:  $Q(x, y) \leftarrow R(x, y), R(y, 5), S(y)$

SPJ:  $\pi_{R_1.A, R_1.B} \sigma_{R_1.B=R_2.A} \sigma_{R_2.B=5} \sigma_{R_1.B=C} (\rho_{R_1}(R) \times \rho_{R_2}(R) \times S)$

## Conclusion:

Select-project-join expressions and conjunctive queries are two separate syntaxes for the same class of queries.

# Optimization of select-project-join expressions

## In-class exercise

Consider the following SQL expression over the relations  $R(A, B)$  and  $S(B, C)$ :

```
SELECT R1.A, S1.B
FROM R R1, R R2, R R3, S S1, S S2
WHERE
    R1.A = R2.A AND R2.B =4 AND R2.A = R3.A
AND R3.B = S1.B AND S1.C = S2.C AND S2.B=4
```

**Exercise:** Translate this SQL expression into a logical query plan. If this plan is an SPJ-expression, then translate this plan into a conjunctive query.

# Optimization of select-project-join expressions

## In-class exercise

Consider the following SQL expression over the relations  $R(A, B)$  and  $S(B, C)$ :

```
SELECT R1.A, S1.B
FROM R R1, R R2, R R3, S S1, S S2
WHERE
    R1.A = R2.A AND R2.B = 4 AND R2.A = R3.A
    AND R3.B = S1.B AND S1.C = S2.C AND S2.B=4
```

Here is the logical query plan:

$$\pi_{R1.A, S1.B} \sigma_{R1.A=R2.A \wedge R2.B=4 \wedge R2.A=R3.A \wedge R3.B=S1.B \wedge S1.C=S2.C \wedge S2.B=4} (\rho_{R1}(R) \times \rho_{R2}(R) \times \rho_{R3}(R) \times \rho_{S1}(S) \times \rho_{S2}(S))$$



# Optimization of select-project-join expressions

## In-class exercise

Consider the following SQL expression over the relations  $R(A, B)$  and  $S(B, C)$ :

```
SELECT R1.A, S1.B
FROM R R1, R R2, R R3, S S1, S S2
WHERE
    R1.A = R2.A AND R2.B = 4 AND R2.A = R3.A
    AND R3.B = S1.B AND S1.C = S2.C AND S2.B = 4
```

Here is the logical query plan:

$$\pi_{R_1.A, S_1.B} \sigma_{R_1.A=R_2.A \wedge R_2.B=4 \wedge R_2.A=R_3.A \wedge R_3.B=S_1.B \wedge S_1.C=S_2.C \wedge S_2.B=4} (\rho_{R_1}(R) \times \rho_{R_2}(R) \times \rho_{R_3}(R) \times \rho_{S_1}(S) \times \rho_{S_2}(S))$$

Constructing the conjunctive query: first step

$$Q(x_{R_1.A}, y_{S_1.B}) \leftarrow R(x_{R_1.A}, y_{R_1.B}), R(x_{R_2.A}, y_{R_2.B}), R(x_{R_3.A}, y_{R_3.B}), \\ S(y_{S_1.B}, z_{S_1.C}), S(y_{S_2.B}, z_{S_2.C})$$

# Optimization of select-project-join expressions

## In-class exercise

Consider the following SQL expression over the relations  $R(A, B)$  and  $S(B, C)$ :

```
SELECT R1.A, S1.B
FROM R R1, R R2, R R3, S S1, S S2
WHERE
    R1.A = R2.A AND R2.B = 4 AND R2.A = R3.A
    AND R3.B = S1.B AND S1.C = S2.C AND S2.B = 4
```

Here is the logical query plan:

$$\pi_{R_1.A, S_1.B} \sigma_{R_1.A=R_2.A \wedge R_2.B=4 \wedge R_2.A=R_3.A \wedge R_3.B=S_1.B \wedge S_1.C=S_2.C \wedge S_2.B=4} (\rho_{R_1}(R) \times \rho_{R_2}(R) \times \rho_{R_3}(R) \times \rho_{S_1}(S) \times \rho_{S_2}(S))$$

Constructing the conjunctive query: forcing equalities

$$Q(x_{R_1.A}, y_{S_1.B}) \leftarrow R(x_{R_1.A}, y_{R_1.B}), R(x_{R_1.A}, 4), R(x_{R_1.A}, y_{S_1.B}), \\ S(y_{S_1.B}, z_{S_1.C}), S(4, z_{S_1.C})$$

# Optimization of select-project-join expressions

## In-class exercise

Consider the following SQL expression over the relations  $R(A, B)$  and  $S(B, C)$ :

```
SELECT R1.A, S1.B
FROM R R1, R R2, R R3, S S1, S S2
WHERE
    R1.A = R2.A AND R2.B = 4 AND R2.A = R3.A
    AND R3.B = S1.B AND S1.C = S2.C AND S2.B = 4
```

Here is the logical query plan:

$$\pi_{R_1.A, S_1.B} \sigma_{R_1.A=R_2.A \wedge R_2.B=4 \wedge R_2.A=R_3.A \wedge R_3.B=S_1.B \wedge S_1.C=S_2.C \wedge S_2.B=4} (\rho_{R_1}(R) \times \rho_{R_2}(R) \times \rho_{R_3}(R) \times \rho_{S_1}(S) \times \rho_{S_2}(S))$$

Constructing the conjunctive query: renaming variables (optional)

$$Q(x, y) \leftarrow R(x, u), R(x, 4), R(x, y), S(y, z), S(4, z)$$

# Optimization of select-project-join expressions

## Another in-class exercise

Consider the following conjunctive query over the relations

- `MovieStar`(name: string, address: string, gender: char, birthdate: date)
- `StarsIn`(movieTitle: string, movieYear: string, starName: string)

$$Q(t) \leftarrow \text{MovieStar}(n, a, g, 1940), \text{StarsIn}(t, y, n)$$

Translate this conjunctive query into an SPJ expression. What is the corresponding SQL query?

# Optimization of select-project-join expressions

## Another in-class exercise

Consider the following conjunctive query over the relations

- `MovieStar`(name: string, address: string, gender: char, birthdate: date)
- `StarsIn`(movieTitle: string, movieYear: string, starName: string)

$$Q(t) \leftarrow \text{MovieStar}(n, a, g, 1940), \text{StarsIn}(t, y, n)$$

Translate this conjunctive query into an SPJ expression. What is the corresponding SQL query?

SPJ expression:

$$\pi_{S.\text{movieTitle}}(\sigma_{M.\text{starName}=S.\text{starName}} \sigma_{M.\text{birthdate} = 1940}(\rho_M(\text{MovieStar}) \times \rho_S(\text{StarsIn})))$$

# Optimization of select-project-join expressions

## Another in-class exercise

Consider the following conjunctive query over the relations

- `MovieStar`(name: string, address: string, gender: char, birthdate: date)
- `StarsIn`(movieTitle: string, movieYear: string, starName: string)

$$Q(t) \leftarrow \text{MovieStar}(n, a, g, 1940), \text{StarsIn}(t, y, n)$$

Translate this conjunctive query into an SPJ expression. What is the corresponding SQL query?

SPJ expression:

$$\pi_{S.\text{movieTitle}}(\sigma_{M.\text{starName}=S.\text{starName}}\sigma_{M.\text{birthdate} = 1940}(\rho_M(\text{MovieStar}) \times \rho_S(\text{StarsIn})))$$

SQL query:

```
SELECT S.movieTitle FROM StarsIn S, MovieStar M
WHERE S.starName = M.name AND M.birthdate = 1960
```

# Optimization of select-project-join expressions

## Containment of conjunctive queries

$Q_1$  is contained in  $Q_2$  if  $Q_1(D) \subseteq Q_2(D)$ , for every database  $D$ .

### Example:

$$\begin{aligned} A(x, y) &\leftarrow R(x, w), G(w, z), R(z, y) \\ B(x, y) &\leftarrow R(x, w), G(w, w), R(w, y) \end{aligned}$$

Then  $B$  is contained in  $A$ . Proof:

1. Let  $D$  be an arbitrary database, and let  $t \in B(D)$ .
2. Then there exists a matching  $f$  of  $B$  into  $D$  such that  $t = (f(x), f(y))$ . We need to show that  $(f(x), f(y)) \in A(D)$ .
3. Let  $h$  be the following substitution:

$$x \rightarrow f(x) \quad y \rightarrow f(y) \quad w \rightarrow f(w) \quad z \rightarrow f(w)$$

4. Then  $h$  is a matching of  $A$  into  $D$ , and hence

$$t = (f(x), f(y)) = (h(x), h(y)) \in A(D)$$

# Optimization of select-project-join expressions

**Containment of conjunctive queries is decidable**

$$A(x, y) \leftarrow R(x, w), G(w, z), R(z, y)$$

$$B(x, y) \leftarrow R(x, w), G(w, w), R(w, y)$$

**Golden method to check whether  $B \subseteq A$ :**

1. First calculate the **canonical database**  $D$  for  $B$ :

$$\begin{array}{c} R \\ \boxed{\begin{array}{cc} x & w \\ w & y \end{array}} \end{array} \quad \begin{array}{c} G \\ \boxed{\begin{array}{cc} w & w \end{array}} \end{array}$$

2. Then check whether  $(x, y) \in A(D)$ . If so,  $B \subseteq A$ , otherwise  $B \not\subseteq A$ .



# Optimization of select-project-join expressions

## Containment of conjunctive queries is decidable

$$A(x, y) \leftarrow R(x, w), G(w, z), R(z, y)$$
$$B(x, y) \leftarrow R(x, w), G(w, w), R(w, y)$$

**Fact:**  $B \subseteq A \Leftrightarrow (x, y) \in A(D)$  with  $D$  the canonical database for  $B$ .

**First possibility:**  $(x, y) \notin A(D)$

In this case we have just constructed a counter-example because  $(x, y) \in B(D)$ .

$R$	$G$						
<table border="1"><tr><td><math>x</math></td><td><math>w</math></td></tr><tr><td><math>w</math></td><td><math>y</math></td></tr></table>	$x$	$w$	$w$	$y$	<table border="1"><tr><td><math>w</math></td><td><math>w</math></td></tr></table>	$w$	$w$
$x$	$w$						
$w$	$y$						
$w$	$w$						

# Optimization of select-project-join expressions

## Containment of conjunctive queries is decidable

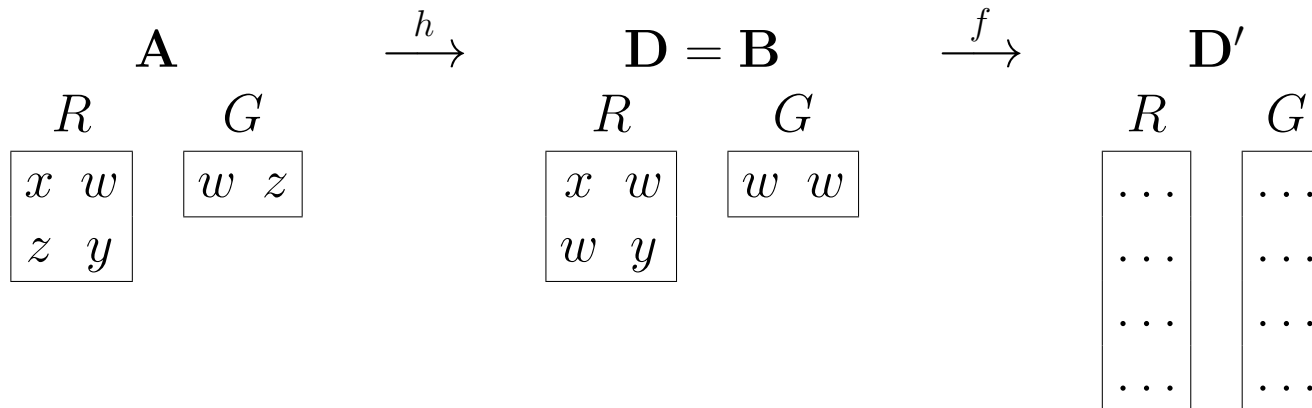
$$A(x, y) \leftarrow R(x, w), G(w, z), R(z, y)$$

$$B(x, y) \leftarrow R(x, w), G(w, w), R(w, y)$$

**Fact:**  $B \subseteq A \Leftrightarrow (x, y) \in A(D)$  with  $D$  the canonical database for  $B$ .

**Second possibility:**  $(x, y) \in A(D)$

- There hence exists a matching  $h$  of  $A$  into  $D$  such that  $h(x) = x$  and  $h(y) = y$
- Let  $D'$  be an arbitrary other database, and pick  $t \in B(D')$ . There hence exists a matching  $f$  such that  $t = (f(x), f(y))$ .
- Then  $f \circ h$  is a matching of  $A$  on  $D'$ :



# Optimization of select-project-join expressions

## Containment of conjunctive queries is decidable

$$\begin{aligned} A(x, y) &\leftarrow R(x, w), G(w, z), R(z, y) \\ B(x, y) &\leftarrow R(x, w), G(w, w), R(w, y) \end{aligned}$$

**Fact:**  $B \subseteq A \Leftrightarrow (x, y) \in A(D)$  with  $D$  the canonical database for  $B$ .

### Second possibility: $(x, y) \in A(D)$

- There hence exists a matching  $h$  of  $A$  into  $D$  such that  $h(x) = x$  and  $h(y) = y$
- Let  $D'$  be an arbitrary other database, and pick  $t \in B(D')$ . There hence exists a matching  $f$  such that  $t = (f(x), f(y))$ .
- Then  $f \circ h$  is a matching of  $A$  on  $D'$ :
- And hence

$$t = (f(x), f(y)) = (f(h(x)), f(h(y))) \in A(D')$$

# Optimization of select-project-join expressions

## Conclusion:

- Containment of conjunctive queries is decidable
- Consequently the **equivalence** of conjunctive queries is also decidable

# Optimization of select-project-join expressions

## Optimizing conjunctive queries

**Input:** A conjunctive query  $Q$

**Output:** A conjunctive query  $Q'$  equivalent to  $Q$  that is optimal (i.e., has the least number of atoms in its body).

**For each conjunctive query we can obtain an equivalent, optimal query, by removing atoms from its body**

- Let  $Q$  be a CQ and let  $P$  be an arbitrary optimal and equivalent query.
- Then  $Q \subseteq P$  and hence  $(x, y) \in P(D_Q)$  with  $D_Q$  the canonical database for  $Q$ . Let  $f$  be the matching that ensures this fact.
- Let  $Q'$  be obtained by removing from  $Q$  all atoms that are not in the range of  $f$
- Then  $Q \subseteq Q'$
- Moreover, also  $Q' \subseteq P$  (because  $(x, y) \in P(D_{Q'})$  still holds) and  $P \subseteq Q$ . Hence  $Q' \equiv Q$ .
- Note that  $Q'$  contains at most the same number of atoms as  $P$ . Hence  $Q'$  is optimal.

# Optimization of select-project-join expressions

## Optimization of conjunctive queries

**Input:** A conjunctive query  $Q$

**Output:** A conjunctive query  $Q'$  equivalent to  $Q$  that is optimal (i.e., has the least number of atoms in its body).

## Optimization algorithm

- A conjunctive query is given. Consider for example:

$$Q(x) \leftarrow R(x, x), R(x, y)$$

- We check, atom by atom, what atoms in its body are redundant.

In our example we first try to delete  $R(x, x)$ :

$$Q_1(x) \leftarrow R(x, y)$$

Note that  $Q \subseteq Q_1$  but  $Q_1 \not\subseteq Q$ . We hence cannot remove this atom.

# Optimization of select-project-join expressions

## Optimization of conjunctive queries

**Input:** A conjunctive query  $Q$

**Output:** A conjunctive query  $Q'$  equivalent to  $Q$  that is optimal (i.e., has the least number of atoms in its body).

## Optimization algorithm

- A conjunctive query is given. Consider for example:

$$Q(x) \leftarrow R(x, x), R(x, y)$$

- We check, atom by atom, what atoms in its body are redundant.

We next try to remove  $R(x, y)$ :

$$Q_2(x) \leftarrow R(x, x)$$

Note that  $Q \subseteq Q_2$  and  $Q_2 \subseteq Q$ .

$Q_2$  is certainly shorter than  $Q$  and hence closer to the optimal query. Since there remain no other atoms to test, our result is  $Q_2$ .

# Optimization of select-project-join expressions

## Optimization of select-project-join expressions

1. Translate the select-project-join expression  $e$  into an conjunctive query  $Q$ .
2. Optimize  $Q$ .
3. Translate  $Q$  back into a select-project-join expression.



# Optimization of logical query plans

## Optimization of complete, arbitrary query plans

1. Detect and optimize the select-project-join sub-expressions in the plan
2. Then use [heuristics](#) to further optimize the modified plan.

## Heuristic: rewriting through algebraic laws

Consider relations  $R(A, B)$  and  $S(C, D)$  and the following expression that we want to optimize

$$\pi_A \sigma_{A=5 \wedge B < D}(R \times S)$$

### Pushing selections:

$$\pi_A \sigma_{B < D}(\sigma_{A=5}(R) \times S)$$

### Recognizing joins:

$$\pi_A(\sigma_{A=5}(R) \bowtie_{B < D} S)$$

### Introduce projections where possible:

$$\pi_A(\sigma_{A=5}(R) \bowtie_{B < D} \pi_D(S))$$

# Optimization of logical query plans

## An in-class integrated exercise

Recall the relational schema.

- MovieStar(name: string, address: string, gender: char, birthdate: date)
- StarsIn(movieTitle: string, movieYear: string, starName: string)

A careless SQL programmer writes the following query:

```
SELECT movieTitle FROM StarsIn S1
WHERE starName IN (SELECT name
                   FROM MovieStar, StarsIn S2
                   WHERE birthdate = 1960
                   AND S2.movieTitle = S1.movieTitle)
```

Translate this query into a logical query plan, remove redundant joins from it, and then use heuristics to further optimize the obtained logical plan.

# Optimization of logical query plans

## An in-class integrated exercise

### SQL

```
SELECT movieTitle FROM StarsIn S1
WHERE starName IN (SELECT name
                   FROM MovieStar, StarsIn S2
                   WHERE birthdate = 1960
                   AND S2.movieTitle = S1.movieTitle)
```

### Translated logical query plan

$$\pi_{S1.movieTitle} \pi_{S1.*,name}$$
$$\sigma_{S2.movieTitle=S1.movieTitle \wedge birthdate=1960 \wedge S1.starName=MovieStar.name}$$
$$(Moviestar \times \rho_{S2}(StarsIn) \times \rho_{S1}(StarsIn))$$

# Optimization of logical query plans

## An in-class integrated exercise

- MovieStar(name: string, address: string, gender: char, birthdate: date)
- StarsIn(movieTitle: string, movieYear: string, starName: string)

## Translated logical query plan

$\pi_{S1.movieTitle} \pi_{S1.*,name}$

$\sigma_{S2.movieTitle=S1.movieTitle \wedge birthdate=1960 \wedge S1.starName=MovieStar.name}$

$(MovieStar \times \rho_{S2}(StarsIn) \times \rho_{S1}(StarsIn))$

## Conjunctive query:

$Q(t) \leftarrow MovieStar(n, a, g, 1960), StarsIn(t, y_2, n_2), StarsIn(t, y, n)$

# Optimization of logical query plans

## An in-class integrated exercise

Conjunctive query:

$$Q(t) \leftarrow \text{MovieStar}(n, a, g, 1960), \text{StarsIn}(t, y_2, n_2), \text{StarsIn}(t, y, n)$$

- the atom  $\text{MovieStar}(n, a, g, 1960)$  cannot be removed (it is the only atom for relation  $\text{MovieStar}$ ).
- the atom  $\text{StarsIn}(t, y_2, n_2)$  can be removed since  $Q \subseteq Q_1$  and  $Q_1 \subseteq Q$  where:

$$Q_1(t) \leftarrow \text{MovieStar}(n, a, g, 1960), \text{StarsIn}(t, y, n)$$

# Optimization of logical query plans

## An in-class integrated exercise

We proceed with the conjunctive query:

$$Q_1(t) \leftarrow \text{MovieStar}(n, a, g, 1960), \text{StarsIn}(t, y, n)$$

- the only remaining atom  $\text{StarsIn}(t, y, n)$  cannot be removed (it is the only atom for relation `StarsIn` and it is the only atom containing the head variable  $t$ ).

# Optimization of logical query plans

## An in-class integrated exercise

The minimal conjunctive query is hence:

$$Q_1(t) \leftarrow \text{MovieStar}(n, a, g, 1960), \text{StarsIn}(t, y, n)$$

Corresponding relational algebra expression:

$$\pi_{S.\text{movieTitle}} \sigma_{M.\text{birthdate}=1960 \wedge S.\text{starName}=M.\text{name}} (\rho_M(\text{Moviestar}) \times \rho_S(\text{StarsIn}))$$

# Optimization of logical query plans

## An in-class integrated exercise

The minimal relational algebra expression is hence:

$$\pi_{S.movieTitle} \sigma_{M.birthdate=1960 \wedge S.starName=M.name} (\rho_M(\text{Moviestar}) \times \rho_S(\text{StarsIn}))$$

Recognize joins:

$$\pi_{S.movieTitle} \sigma_{M.birthdate=1960} (\rho_M(\text{Moviestar})) \bowtie_{S.starName=M.name} \rho_S(\text{StarsIn})$$

Push selections:

$$\pi_{S.movieTitle} (\sigma_{M.birthdate=1960} (\rho_M(\text{Moviestar})) \bowtie_{S.starName=M.name} \rho_S(\text{StarsIn}))$$

Add projections:

$$\pi_{S.movieTitle} (\sigma_{M.birthdate=1960} (\pi_{M.birthdate, M.name} \rho_M(\text{Moviestar})) \bowtie_{S.starName=M.name} \pi_{S.movieTitle, S.starName} \rho_S(\text{StarsIn}))$$