## Graph Database System Neo4j

Hannes Voigt

## Some Bachground

## Adjacency Matrix

Undirected graph without labels

## Directed graph with edge labels




## Adjacency Lists

## Compression of adjacency matrix

- Compression scheme:



## ADJACENCY LIST

- Source vertextogether outgoing edges

Without edge labels

$$
\begin{array}{ll}
0 & -> \\
1 & ->(1,2) \\
2 & -> \\
3 & ->(1,1,2)
\end{array}
$$

Compressed Sparse Row

source-oriented


With edge labels

$$
\left.\begin{array}{l}
0->([0, a],[2, c]) \\
1->([1, b],[3, e]) \\
2->([0, d],[1, i],[2, f]) \\
3
\end{array}\right)
$$

Compressed Sparse Column

target-oriented

Coordinate list
$(0,0, a)$
$(2,2, f)$

Almost the same!!!

With edge properties

$$
\begin{array}{ll}
0 & ->([0, a,(\text { weight }=4)],[2, c,(\text { weight }=3)]) \\
1 & ->([1, b,(\text { weight }=3)],[3, e,(\text { weight }=2)]) \\
2 & ->([0, d,(\text { weight }=5)],[1, i,(\text { weight }=2)], \ldots) \\
3 & ->([1, h,(\text { weight }=9)],[2, g,(\text { weight }=7)])
\end{array}
$$

## Basic Terminology

## Walk

- Sequence of edges connecting vertices
- $w_{v_{1}, v_{n}}=v_{1} e_{1} v_{2} \ldots v_{n-1} e_{n-1} v_{n}$ with
$e_{i}=\left(v_{i}, v_{i+1}\right) \in E, 1 \leq i<n$
- $n$ is the length of the walk



## PATH

- Walk connecting distinct vertices
- $p_{v_{1}, v_{n}}=v_{1} e_{1} v_{2} \ldots v_{n-1} e_{n-1} v_{n}$ with $\forall v_{i}, v_{j}: i \neq j \rightarrow v_{i} \neq v_{j}$ and $e_{i}=\left(v_{i}, v_{i+1}\right) \in E, 1 \leq i<n$
- Length is number of edges or sum of edge weights



## Cycle

- Walk with $v_{1}=v_{n}$ is a cycle
- Graph is acyclic iff $\nexists v \in V: w_{v, v} \in G$



## Degree (or valency)

- In/out degree of a vertex: Number of incoming/outgoing edges of that vertex
- $\operatorname{deg}_{\text {out }}(v)=|\{(v, u) \in E\}|$
- $\operatorname{deg}_{i n}(v)=|\{(u, v) \in E\}|$

- $\operatorname{deg}(v)=\operatorname{deg}_{\text {out }}(v)+\operatorname{deg}_{\text {in }}(v)$


## Distance

- Distance between two vertices in a graph is number of edges in a shortest path connecting them
- $d(v, u)=\min _{p_{v, u} \in G}\left|p_{v, u}\right|$


## Diameter

- Maximum eccentricity of any vertex in the graph
- $d(G)=\max _{v \in V} \epsilon(v)$



## Finding Shortest Paths

## UNWEIGHTED SHORTEST PATHS

- Length of path its number of edges
- Restriction to simple paths (w/o cycles)

- Two main ways of path search
- Depth-first search (DFS)
- Breadth-first search (BFS)


## DEPTH-FIRST SEARCH (DFS)

- Search tree is deepened as much as possible on each child before going to the next sibling
- Lower space complexity
- Has to examine whole graph to find shortest path between two nodes


## Breadth-First search (BFS)

- Search tree is broadened as much as possible on each depth before going to the next depth
- Potential large space required
- Find shortest path between two nodes first (before finding a longer one)




## Bidirectional BFS

## IDEA: SEARCH FROM START AND END VERTEX

- Alternatingly explore vertices on both sides
- Optimization: explore vertices on the side with smaller frontier
- Algorithm stops when both BFS meet
- When discovering a new vertex, each BFS check if that vertex is in frontier of other side

Forward (598 vertices explored)


Backward (860 vertices explored)


Bidirectional (448 vertices explored)


## Centrality Measures

Question: Who are The key players in a graph

- Most social contacts (vaccination schedules)
- Most influential thinkers/papers (reading lists)
- Most important website (web search)
- Most important distributers (supply network)
- etc.
- Can we measure that?


## Yes! With centrality Measures!

- Centrality measures identify the most important verticeswithin a graph







































 [http://brendangriffen.com/blog/gow-influential-thinkers]


## Centrality Measures

Dresden Database
Systems Group

## Various centrality measure have been defined

- Betweenness centrality (A)
- Number of shortest paths between all other vertices that pass through that vertex
- Closeness centrality (B)
- Reciprocal of the sum of distances to all other vertices
- Harmonic centrality (E) uses the sum of reciprocal of distances instead
- Eigenvector centrality/Eigencentrality (C)
- Score of a vertex contributes to score of neighboring vertices
- Page rank is variant of eigenvector centrality
- Katz centrality (F)
- Number of all vertices that can be connected through a path
- Contributions of distant nodes are penalized
- Degree centrality (D) only considers direct neighbors


Neo4j

## Database Systems Landscape



451 Research
Data Platforms Map
October 2014

## Kev: General purpose <br>  <br> -as-a-Service <br> ${ }^{\text {Gigraph }}$ <br> $=\begin{aligned} & \text { Document } \\ & \text { Key value store } \\ & \text { Kever }\end{aligned}$ <br> $\underbrace{\text { Key value direct }}_{\text {Key value stores }}$ <br> Key value access Hadoop <br> $\longrightarrow_{\text {MySQL ecosystem }}^{\text {Hadoop }}$ <br> Advanced clustering/sharding New SOL databases <br>  <br> $\qquad$ <br> $\square$ Search <br> In-memory <br> https:// <br> 451research.com/ dashboard/dpa <br> © 2014 by 451 Research LLC

## Neo4j Terminology

Label


## Match

## MATCH-CLAUSE

- Primary way of getting data from a Neo4j database
- Allows you to specify the patterns
- Named pattern element, e.g. (p:Person), will be bound to the match instance
- Query can have multiple MATCH-clauses


## RETURN-CLAUSE

- Projects to the result set
- Allows projection to nodes, edges, and properties

MATCH (p:Person)-[:Likes]->(f:Person) RETURN p.name, f.sex

MATCH (p:Person)-[:Likes]->(:Person) -[:Likes]->(fof:Person) RETURN p.name, fof.name

Dresden Database
Systems Group

| p.name | fof.name |
| :--- | :--- |
| Lucy | Jen |
| Peter | Lucy |
| Peter | Peter |
| Jen | Peter |
| Lucy | Lucy |

## Pattern Syntax

## Vertex pattern

- ()
- (matrix)
- (:Movie)
- (matrix:Movie:Action)
- (matrix:Movie \{title: "The Matrix"\})
- (matrix: Movie \{title: "The Matrix", released: 1997\})


## Edge pattern

- -->
-     - [role]->
-     - [:ACTED_IN]->
-     - [role:ACTED_IN]->
-     - [role:ACTED_IN \{roles: ["Neo"]\}]->
unidentified vertex
vertex identified by variable matrix
unidentified vertex with label Movie
vertex with labels Movie and Action identified by variable matrix
+ property title equal the string "The Matrix"
+ property released equal the integer 1997
unidentified edge
edge identified by variable role
unidentified edge with label ACTED_IN
edge with label $A C T E D_{-} I N$ identified by variable role
+ property roles contains the string "Neo"


## Pattern Syntax

## Path Patterns

- String of alternating vertex pattern and edge pattern
- Starting and ending with a vertex pattern
- (a)-->(b)<--(c)--(d)-->(a)-->(e)

- (keanu:Person:Actor \{name: "Keanu Reeves"\}) - [role:ACTED_IN \{roles: ["Neo"]\}]-> (matrix:Movie \{title: "The Matrix"\})


## GRAPH PATTERNS

- One or multiple path patterns
- Path patterns should have at least one shared variable
- Without shared variable graph pattern is disconnected
- Results in a cross-product of the results for connected sub patterns
- Quadrating blow up in result size and computational complexity
- (a) $-->(b)<--(c)--(d)-->(a)-->(e),(e)-->(b)-->(d),(a)-->(a)$



## Return

## RETURN-CLAUSE

- Defines what to include in the query result set
- Comparable with relational projection
- Only once per query
- Allows to return nodes, edges, properties, or any expressions
- Column can be rename using AS <new name>


## Example

- MATCH (n)

RETURN n, "node " $+\mathrm{id}(\mathrm{n})+$ " is " +
$\begin{array}{ll}\text { CASE } & \text { WHEN n.title IS NOT NULL THEN "a Movie" } \\ & \text { WHEN EXISTS(n.name) THEN "a Person" } \\ & \text { ELSE "something unknown" }\end{array}$


END AS about

## Optional Match \& Where

[http://neo4i.com/docs/developer-manual/current/\#query-optional-match]
[http://neo4j.com/docs/developer-manual/current/\#query-where]

## OPTIONAL MATCH-clause

- Matches patterns against your graph database, just like MATCH
- Matches the complete pattern or not
- If no matches are found, OPTIONAL MATCH will use NULLs as bindings
- Like relational outer join
- Example: MATCH (a:Movie)

$$
\text { OPTIONAL MATCH }(a)<-[\text { :WROTE]-(x) }
$$

RETURN a.title, x.name

## WHERE

- After an (OPTIONAL) MATCH, it adds constraints to the (optional) match
- WHERE becomes part of the pattern
- After a WITH, it just filters the result
- Syntax: WHERE <expression>
- Example: MATCH (n)

WHERE n.name $=$ 'Peter' XOR (n.age $<30$ AND n.name $=$ 'Tobias') OR NOT (n.name ~= 'Tob.*' OR n.name CONTAINS 'ete')

| \$ MATCH (a:Movie) OPTIO... |  | + | \# | ${ }^{*}$ | $\boldsymbol{*}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $\underset{\text { Rows }}{\text { 且 }}$ | a.title | x.name |  |  |  |
|  | The Matrix | null |  |  |  |
| $\begin{gathered} \langle/\rangle \\ \text { Code } \end{gathered}$ | The Matrix Reloaded | null |  |  |  |
|  | The Matrix Revolutions | null |  |  |  |
|  | The Devil's Advocate | null |  |  |  |
|  | A Few Good Men | Aaron Sorkin |  |  |  |
|  | Top Gun | Jim Cash |  |  |  |
|  | Jerry Maguire | Cameron Crowe |  |  |  |
|  | Stand By Me | null |  |  |  |
|  | As Good as It Gets | null |  |  |  |
|  | What Dreams May Come | null |  |  |  |
|  | Snow Falling on Cedars | null |  |  |  |
|  | You've Got Mail | null |  |  |  |
|  | Sleepless in Seattle | null |  |  |  |
|  | Joe Versus the Volcano | null |  |  |  |
|  | When Harry Met Sally |  | Nora Ephron |  |  |
|  | That Thing You Do |  | null |  |  |
|  | Returned 41 rows in 40 ms . |  |  |  |  |

RETURN n

## Matching Paths

## Variable lengTh Path Patterns

- Repetitive edge types can be expressed by specifying a length with lower and upper bounds
- Example:
(a) $-[: x * 2]->(b)$
is equal to
(a) $-[: x]->()-[: x]->(b)$
- More examples:
(a) $-[* 3 . .5]->(b)$
(a) $-[* 3 .]-.>$ (b)
(a) $-[* . .5]->(b)$
(a) $-[\star]->$ (b)
- Complete example: MATCH (me)-[:KNOWS*1..2]-(remote_friend)

WHERE me.name = "Filipa" RETURN remote_friend.name

- Matches unique paths (relationship uniquness), not unique reachable nodes!!!
- Particularly the unbounded [*] easily matches larger numbers of paths -> exponential blowup!!!


## Path variables

- Assign matched paths to variable or further processing
- Example:

$$
p=((a)-[\star 3 . .5]->(b))
$$

## Matching Shortest Paths

## SHORTEST PATHS

- Path between two nodes with minimum number of edges
- Apply the shortestPath/allShortestPath function to a path pattern to match single/all shortest paths
- Additional filter predicates can be given with WHERE clause
- Universal (NONE/ALL) predicates can be evaluated during shortest path search
- Other predication can be evaluated only after shortest path has been discovered
- Fast evaluation algorithm
- Bidirectional BFS
- Standard for paths without additional predicates and path with universal predicates
- Slow evaluation algorithm
- DFS
- Fallback for paths with non-universal predicates
- Example (fast evaluation):

MATCH (m \{ name:"Martin Sheen" \}),(o \{ name:"Oliver Stone" \}), p = shortestPath((m)-[*..15]-(o)) WHERE NONE(r IN rels(p) WHERE type(r)= "FATHER") RETURN p

- Example (fast evaluation):


MATCH (m \{ name:"Martin Sheen" \}),(o \{ name:"Oliver Stone" \}), p = shortestPath((m)-[ぇ..15]-(o)) WHERE length(p) > 1 RETURN $p$

## Aggregation

## In RETURN-CLAuse

- Implicit group by
- Expressions without an aggregation function will grouping keys
- Expressions with an aggregation function will produce aggregates
- DISTINCT within the aggregation function removes duplicates in a group before the aggregation
- Aggregation function: COUNT, SUM, AVG, MIN, MAX, STDEV, STDEVP, PERCENTILEDISC, PERCENTILECONT, and COLLECT - collects all the values into a list


## In WITH-CLAUSE

- Like a process pipe
- Chains query parts together, piping the results from one to be used as starting points in the next
- Like RETURN, WITH defines - including aggregation - the output before it is passed on
- Allows to
- Filter on aggregates
- Aggregation of aggregates
- Limit search space based on order of properties or aggregates


## Aggregation

## In RETURN-CLAUSE

- Implicit group by
- Expressions without an aggregation function will be group keys
- Expressions with an aggregation function will produce aggregates
- DISTINCT within the aggregation function removes duplicates in a group before the aggregation
- Aggregation function: COUNT, SUM, AVG, MIN, MAX, STDEV, STDEVP, PERCENTILEDISC, PERCENTILECONT, and COLLECT - collects all the values into a list
- Example: MATCH (me:Person \{name:'Ann'\})-->(friend:Person)-->(friend_of_friend:Person)

RETURN me.name, count(DISTINCT friend_of_friend), count(friend_of_friend)


## In WITH-clause

- See next slide


Result

| me | COUNT <br> DISTINCT | COUNT |
| :--- | :--- | :--- |
| Ann | 3 | 4 |

## Query Composition

## WITH-cLAUSE

- Like a process pipe
- Chains query parts together, piping the results from one to be used as starting points in the next
- Like RETURN, WITH defines - including aggregation - the output before it is passed on
- Filter on aggregates

Example: Soccer team on average younger than 25
MATCH (p)-[:PLAYS]->(t) WITH t, AVG(p.age) AS a WHERE a < 25 RETURN t

- Aggregation of aggregates

Example: Average age of the youngest player in each team
MATCH (p)-[:PLAYS]->(t) WITH t, MIN(p.age) AS a RETURN AVG(a)

- Limit search space based on order of properties or aggregates

Example: Friends of five best friends
MATCH (p)-[f:FRIENDS]->(p2)
WITH f.p2 ORDER BY f.rating DESC LIMIT 5
MATCH (p2)-[f:FRIENDS]->(p3) RETURN DISTINCT p3

## Exercise

## Exercise

## Preparation

- Download and install neo4j community edition:
- For installation follow standard download procedure: http://neo4j.com/download/
- For portable usage without installation download archive (tar/zip): https://neo4j.com/download/other-releases/ and follow OS-specific installing instructions at download page
- Import the movie database $>$ : play movie graph $>2^{\text {nd }}$ page $>$ click on code $>$ execute
- Try out query: MATCH (n) WITH COUNT(n) AS numVertices

MATCH (a)-[e]->(b)
RETURN numVertices, COUNT(e) AS numEdges

- Try out query: MATCH (n) RETURN n

Expected Result

| numVertices | numEdges |
| :--- | :--- |
| 171 | 253 |

## Add Data

- Add movie, actor (three main characters), director as vertices and ACTED_IN/DIRECTED edges for the movie The Bridges of Madison County http://www.imdb.com/title/tt0112579/
- Do not insert vertices that already exist in the database!!!
$\qquad$


## Exercise

## Simple patterns

- Find all actors that directed a movie they also acted in and return actor and movie nodes
- Find all reviewer pairs, one following the other and both reviewing the same movie, and return entire subgraphs
- Find all reviewer pairs, one following the other, and return the two reviewers and a movie the may have reviewed both
- Restrict previous query so that the name of the followed reviewer is not 12 characters long
- Try a different position for the where clause. Explain why this gives a different result.
- Find all actors that acted in a movie together after 2010 and return the actor names and movie node
- By extending the previous query, find all movies that the cast of the movies found before also acted in


## Matching semantics of Neo4J

- Which matching semantics does Neo4j implement? Homomorphism, Isomorphism, Induced subgraph isomorphism?
- Remove duplicates for pattern (x)--(y)
- Match pattern (a1)-[:REVIEWED]->(m)<-[:REVIEWED]-(a2) as induced subgraph
- Find all actor pairs that acted in multiple movies together
- Find all pairs of actor-movie subgraphs with equal roles (on ACTS_IN edges), return actors names, roles, and movie titles


## Exercise

## Paths

- Match all reviewers and the one they are following directly or via another a third reviewer
- Count the number of paths of at most length 4 starting from Clint Eastwood ignoring edge direction
- Count the number of paths of at most length 10 starting from Clint Eastwood ignoring edge direction
- Count the number of paths of at most length 11 starting from Clint Eastwood ignoring edge direction
- Count the number of nodes reachable in at most 4 hops starting from Clint Eastwood ignoring edge direction
- Count the number of nodes reachable in at most 10 hops starting from Clint Eastwood ignoring edge direction
- Count the number of nodes reachable in at most 11 hops starting from Clint Eastwood ignoring edge direction


## Exercise

## YOUNG AND OLD MOVIES

- Determine the average age of the Apollo 13 cast at the time of the movie's release
- Find the movies with the top-10 oldest cast at the time of the movie's release
- Return movie and average age rounded to two decimal ordered by descending age
- Find average age of youngest actors in movie casts at time of release
- Find ACTED_IN subgraph of the movie with the youngest cast at the time of the movie's release
- Determine the movie with youngest and movie with oldest cast and their age difference rounded to two decimal points


## ADJACENCY LIST AND DISTRIBUTIONS

- Return the whole graph a simple adjacency list of vertex ids ordered by decreasing vertex degree
- Return out degree distribution ordered by ascending degree
- Return degree distribution ordered by ascending degree
- Return edge types with number of instances order by decreasing instances number


## Exercise

## SIX DEGREES OF KEVIN BACON Inttps:/en.wikipediborgywidisi_degees_of_kevin_bacon)

- Determine the Bacon number of Clint Eastwood
- Count for each Bacon number the number of actor
- Return degree and number of actors ordered by ascending degree


## KATZ CENTRALITY [nttps://en.wikipedia.org/wiki/Katz_centrality]

- Find actors with top 10 Katz centrality along ACTED_IN edges
- Distance penalty is reciprocal of path length (e.g. 3-hop neighbor gets a penalty of 1/3)
- Return actor vertex and Katz centrality


## Exercise

Dresden Database
Systems Group

## Hints

- Use the neo4j browser (web frontend) http://neo4j.com/docs/stable/toolswebadmin.html
- Use the Cypher documentation: http://neo4j.com/docs/stable/cypher-query-lang.html
- Use your preferred search engine
- Try out! Explorer! Have fun!!!


