



Fourth European Business Intelligence Summer School (eBISS 2014)

July 6 - July 11, 2014 Berlin, Germany

eBISS summer school 2014 in Berlin

From the Web of Data
To the Linked Data

**The Web of Data, understanding the
technological keys and publishing Linked Data**

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Atos



- The Data, open or not, always raise the problems of heterogeneity: many data formats, many metadata schema descriptions, complexity to mix... Using the Web of Data technologies, these problems may be overcome. Semantic Web technologies allow to move from raw data to linked data after elevation and interlinking phases. However, lifting raw Data to Linked Data is far from being straightforward. We will see the whole process and the results of a concrete implementation.

 1. The **basic principles** of the Web of Data: from raw and heterogeneous Data to the lingua franca of the Web of Data
 2. The **RDF** semantic and its main syntaxes
 3. **Handling and querying** Data (triple store, end point, SPARQL)
 4. Understanding the basics of **knowledge modeling** (ontology, OWL)
 5. The **lifting and linking** of Data in order to publish 5 Stars Data (ex. DataLift plateform).



1

The basic principles of the Web of Data

from raw and heterogeneous
Data to the lingua franca of
the Web of Data

1. The basic principles of the Web of Data: from raw and heterogeneous Data to the lingua franca of the Web of Data
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4. Understanding the basics of knowledge modeling (ontology, OWL)
5. The lifting and linking of Data in order to publish 5 stars Data (DataLift).

A data, alone,
has no value...

Data

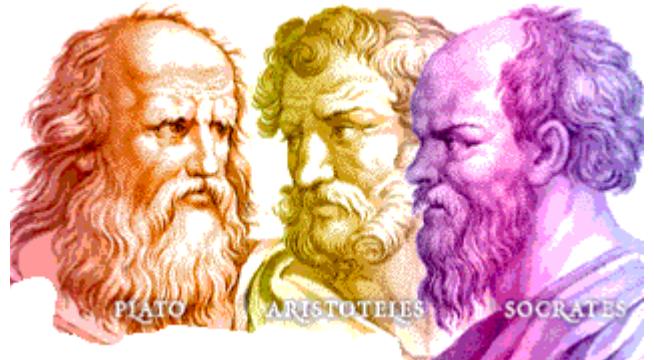
Datasets

Sources of data

Data stores



From data to information
From information to knowledge



\wedge	y
x	0 1
0	0 0
1	0 1

\vee	y
x	0 1
0	0 0
1	1 1

\rightarrow	y
x	0 1
0	1 1
1	0 1

\oplus	y
x	0 1
0	0 1
1	1 0

Figure 1. Truth tables



Figure 2. Logic gates

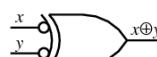
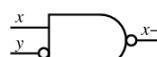
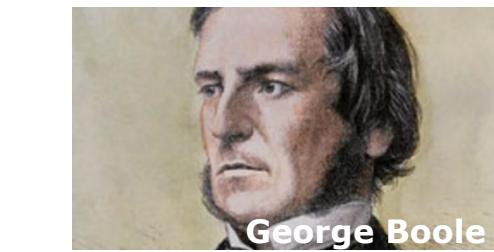


Figure 3. De Morgan equivalents



Figure 4. Venn diagrams



George Boole

*5443. $\vdash \alpha, \beta \in L. \Box : \alpha \wedge \beta = \Lambda . \equiv . \alpha \vee \beta \in L$

Dem.

$\vdash *5426. \Box \vdash . \alpha = t^i x . \beta = t^j y . \Box : \alpha \vee \beta \in L . \equiv . x \neq y .$

$\equiv . t^i x \wedge t^j y = \Lambda .$

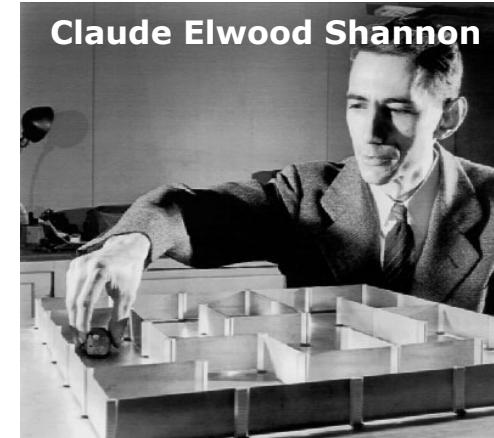
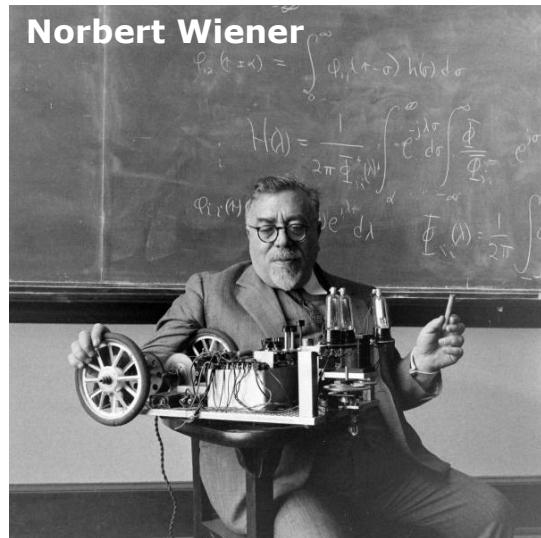
$\equiv . \alpha \wedge \beta = \Lambda .$ (1)

$\vdash . (1) . *11\text{-}11\text{-}35. \Box$

$\vdash . (\Diamond x, y) . \alpha = t^i x . \beta = t^j y . \Box : \alpha \vee \beta \in L . \equiv . \alpha \wedge \beta = \Lambda .$ (2)

$\vdash . (2) . *11\text{-}54 . *521. \Box \vdash . \text{Prop}$

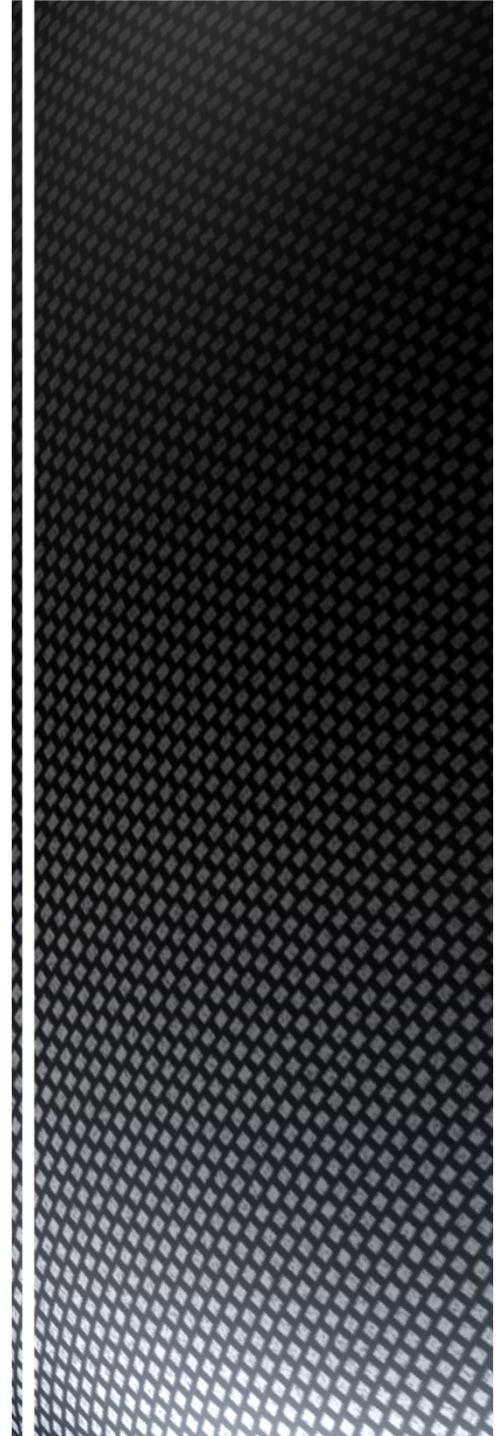
From this proposition it will follow, when arithmetical addition has been defined, that $1 + 1 = 2.$

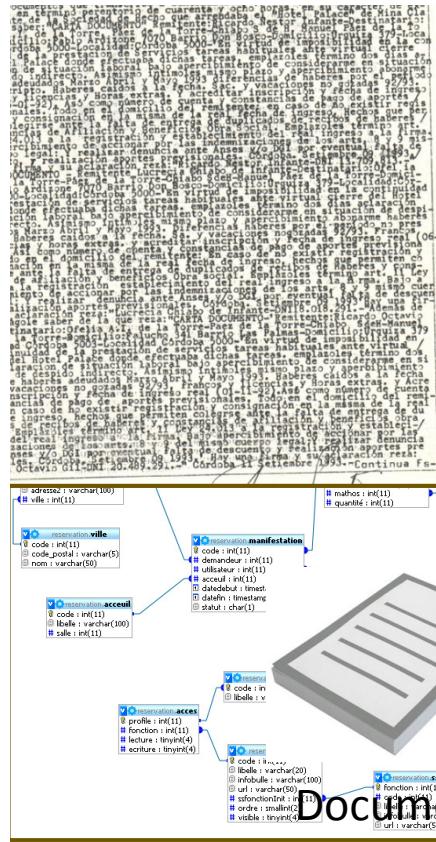




2007, 2008, ...

The Data





Document (CMS)

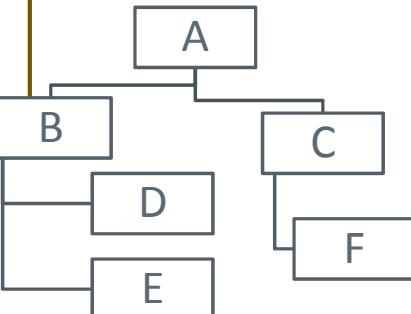
	x	y	z
a	1	2	4
b	0	9	6
c	9	0	7

Data base (SQL)

<Books>

```

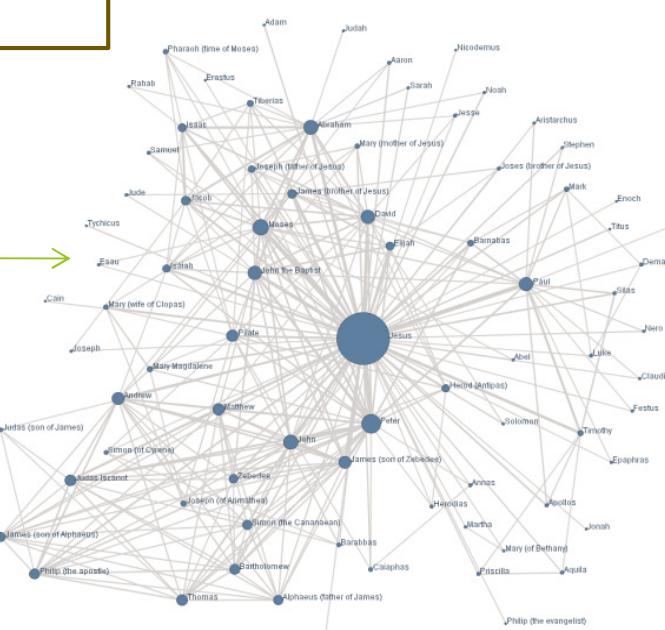
        <Book ISBN="0553212419">
            <title>Sherlock Holmes: Complete Novels...
            <author>Sir Arthur Conan Doyle</author>
        </Book>
        <Book ISBN="0743273567">
            <title>The Great Gatsby</title>
            <author>F. Scott Fitzgerald</author>
        </Book>
        <Book ISBN="0684826976">
            <title>Undaunted Courage</title>
            <author>Stephen E. Ambrose</author>
        </Book>
        <Book ISBN="0743203178">
            <title>Nothing Like It In the World</title>
            <author>Stephen E. Ambrose</author>
        </Book>
    </Books>
    
```



Hierarchy (XML)

	animal	hair	feathers	eggs	milk	airborne	aquatic	predator	toothed
1	String	String	String	String	String	String	String	String	String
2	aardvark	true	false	false	true	false	false	true	true
3	antelope	true	false	false	true	false	false	false	true
4	bass	false	false	true	false	false	true	true	true
	true	false	false	true	false	false	true	true	true
	true	false	false	true	false	false	true	true	true
	true	false	false	true	false	false	false	true	true
	false	false	true	false	false	true	true	false	true
	false	false	true	false	false	true	true	false	true
	true	false	false	true	false	false	false	true	true
	true	false	false	true	false	false	false	true	true
	false	true	true	false	true	false	false	false	false

Graph (RDF)





Do we talk about data without giving them meaning?

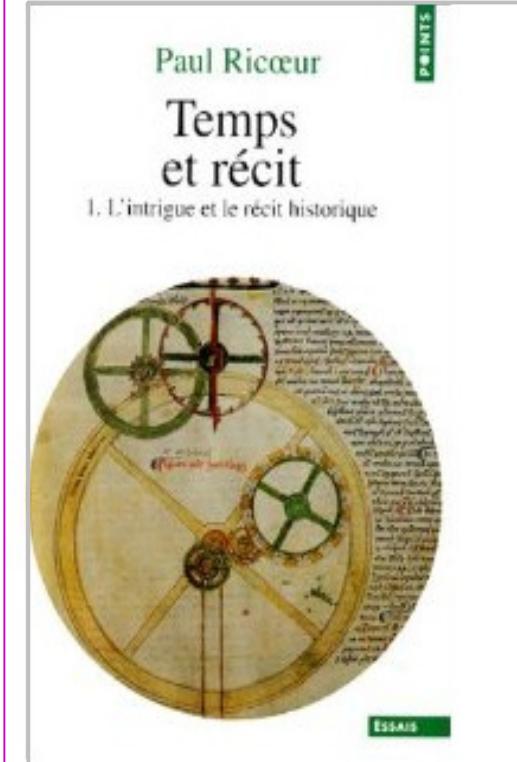
- Simple information
 - The name of Voltaire
 - The number of legs of a spider
 - The fact that the door is closed
- A grouping of data, document
 - A biography, a text, a newspaper
 - A table of values
 - A web page
- A collection of data
 - A list of composers
 - A set of test scores
 - The names of the months of the year

The elementary data

Many limitations

- Research possibilities (lexical)
- The presentation is summary
- The freshness of data
- The reliability of data
- The comparison of available sources
- The evaluation of sources
- Etc.

There is always much more data than we think



Paul Ricoeur
Temps et récit
1. L'intrigue et le récit historique
ESSAIS

Temps et récit, tome 1 [Poche]

Paul Ricoeur (Auteur)

 (2 commentaires client)  J'aime (0)

Prix conseillé : EUR 8,50

Prix : **EUR 8,07** **LIVRAISON GRATUITE** [En savoir plus.](#)

Économisez : EUR 0,43 (5%)

En stock.

Expédié et vendu par **Amazon.fr**. Emballage cadeau disponible.

Plus que 8 ex (réapprovisionnement en cours). Commandez vite !

Voulez-vous le faire livrer le mercredi 22 juin ?

Commandez-le dans les **14 h et 53 min** et choisissez la **livraison Éclair** sur votre bon de commande. [En savoir plus.](#)

9 neufs à partir de EUR 8,07 **5 d'occasion** à partir de EUR 2,90

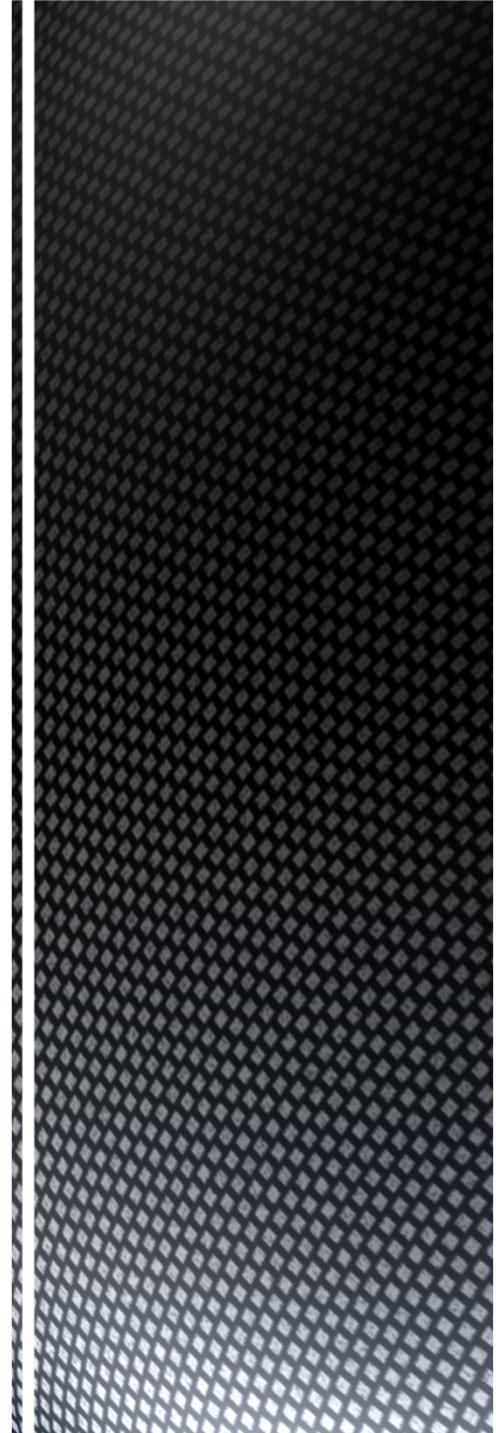
- The book is titled "Temps et récit"
- The book was written by Paul Ricoeur
- The discount is 5%
- ETC.

Understanding,
knowledge?

Structuring data

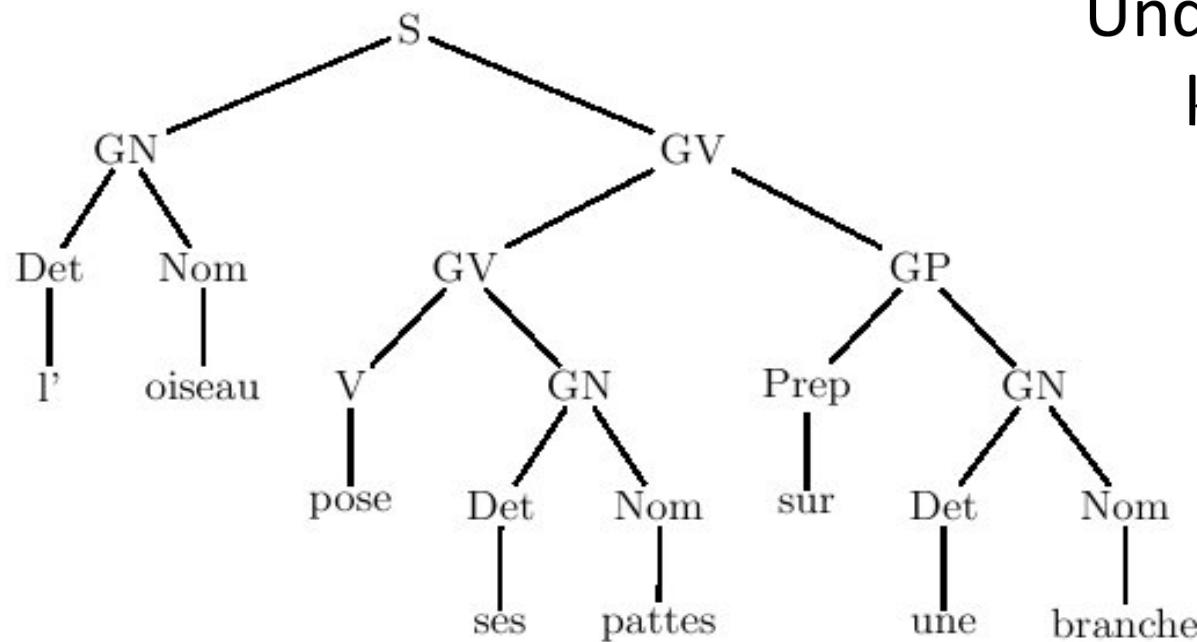
Representing informations

**They are ways to gain in understanding
and knowledge.**





Understanding,
knowledge?



WEB = A technological mean to make available, to link and to share **documents** on a network of connected **machines**.

The four building blocks of the Web

A protocole



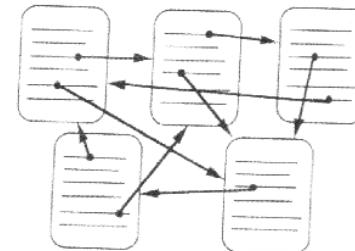
HTTP

A name
(identification's
mecanism)



URL/URI

A principle



L'hypertexte

A langague

```
<!DOCTYPE html PUBLIC "-//W3C//DTD XHTML 1.0 Transitional//EN"
"http://www.w3.org/1999/xhtml">
<html xmlns="http://www.w3.org/1999/xhtml">
<head>
  <title>Ma première page Web</title>
</head>
<body>
  <h1>Titre de niveau 1</h1>
  <p>Paragraphe avec une <em>épithète</em></p>
  <p>Une autre paragraphe avec un lien s</p>
  <ul>
    <li>
      Item de listes à puces
    </li>
    <li>
      Un autre item
    </li>
  </ul>
</body>
</html>
```

HTML

WEB = A technological mean to make available, to link and to share **ressources** on a network of connected **devices**.

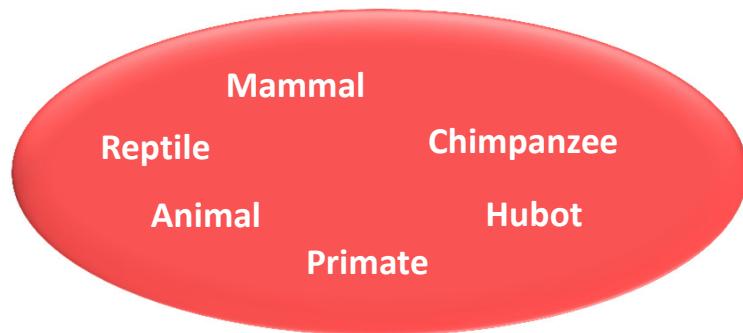
Describing the data

Classification schemes

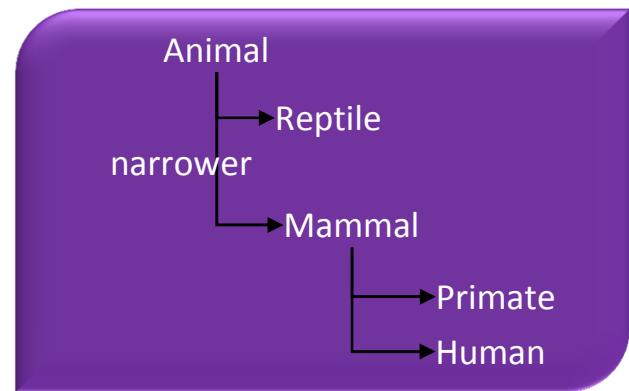
Formats to represent them

(Major) Knowledge Organization's Systems

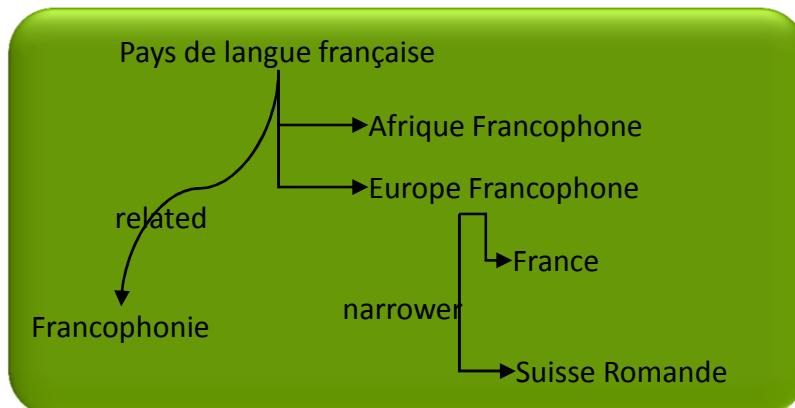
Controlled vocabulary



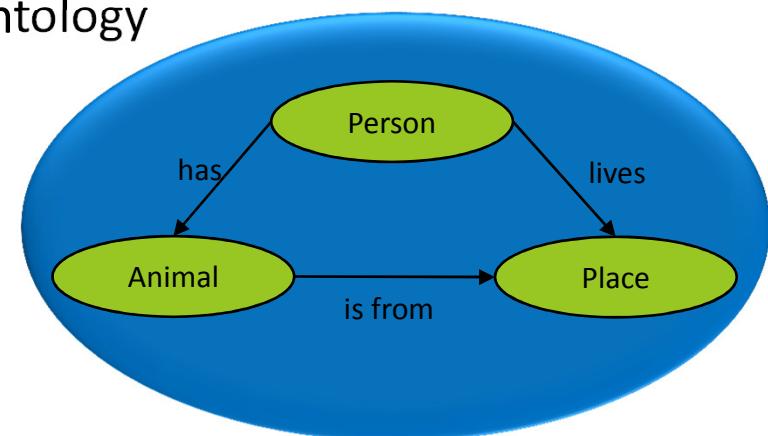
Taxonomy



Thesaurus



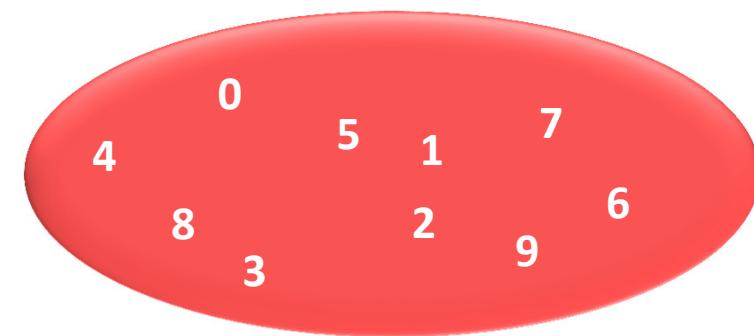
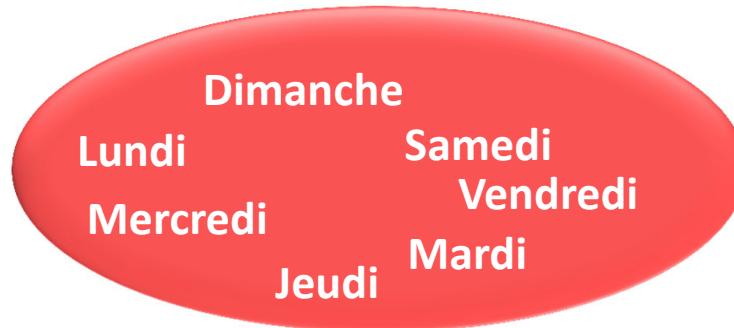
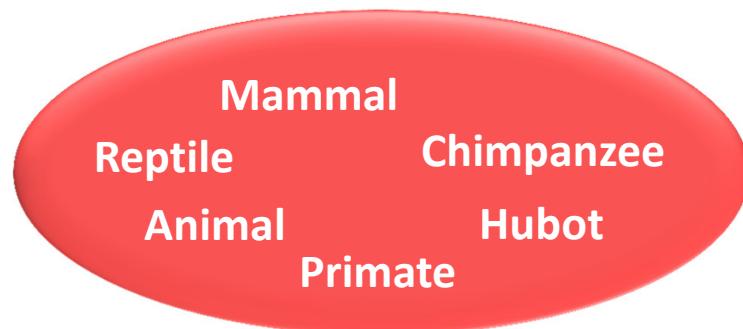
Ontology



Controlled vocabulary

Controlled vocabulary

or controlled lexicon or
list of standardized terms



A controlled vocabulary is a set of recognized terms, fixed, unalterable, standardized and validated by a group (a community of practices) used to index and analyze contents and search for information in a defined area of information.

```
<?xml version="1.0" encoding="UTF-8" ?>
<xs:schema xmlns:xs="http://www.w3.org/2001/XMLSchema"
  xmlns="http://purl.org/gem/instance/level/" targetNamespace="http://purl.org/gem/instance/level/"
  elementFormDefault="qualified" attributeFormDefault="unqualified">
  <xs:annotation>
    <xs:documentation xml:lang="en">GEM controlled vocabulary for designating the education level of the
      audience for a resource. Created in 1996 by the GEM Consortium.</xs:documentation>
  </xs:annotation>
  <xs:simpleType name="Level">
    <xs:union>
      <xs:simpleType>
        <xs:restriction base="xs:Name">
          <xs:enumeration value="Preschool education">
          <xs:enumeration value="Kindergarten">
          <xs:enumeration value="Grade 1">
          <xs:enumeration value="Grade 2">
          <xs:enumeration value="Grade 11">
          <xs:enumeration value="Grade 12">
          <xs:enumeration value="Adult/continuing education">
          <xs:enumeration value="Higher education">
          <xs:enumeration value="Vocational education">
        </xs:restriction>
      </xs:simpleType>
    </xs:union>
  </xs:simpleType>
</xs:schema>
```

Controlled vocabulary an XML sample

GEM educationLevel Vocabulary (XML)

Controlled vocabulary

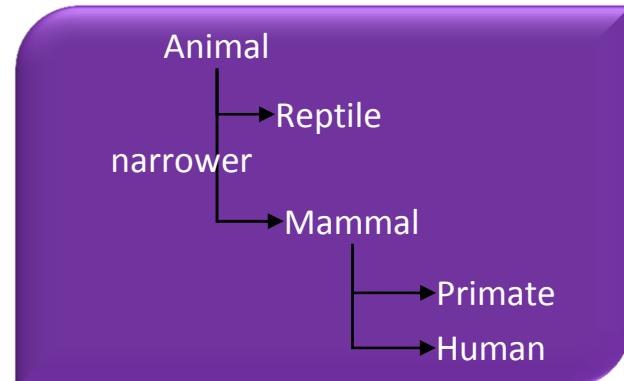
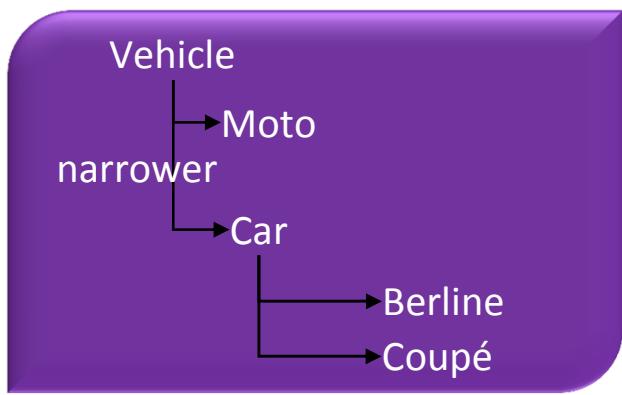
an JSON sample

```
{ "countries": [  
    {"country": {"country_id": "AD", "country_name": "ANDORRA"}},  
    {"country": {"country_id": "AE", "country_name": "UNITED ARAB EMIRATES"}},  
    {"country": {"country_id": "AF", "country_name": "AFGHANISTAN"}},  
    {"country": {"country_id": "AG", "country_name": "ANTIGUA AND BARBUDA"}},  
    {"country": {"country_id": "AI", "country_name": "ANGUILLA"}},  
    {"country": {"country_id": "AL", "country_name": "ALBANIA"}},  
    {"country": {"country_id": "AM", "country_name": "ARMENIA"}},  
    {"country": {"country_id": "AN", "country_name": "NETHERLANDS ANTILLES"}},  
    {"country": {"country_id": "AO", "country_name": "ANGOLA"}},  
    {"country": {"country_id": "AQ", "country_name": "ANTARCTICA"}},  
    {"country": {"country_id": "AR", "country_name": "ARGENTINA"}},  
    {"country": {"country_id": "AS", "country_name": "AMERICAN SAMOA"}},  
    {"country": {"country_id": "AT", "country_name": "AUSTRIA"}},  
    Etc.  
    {"country": {"country_id": "ZM", "country_name": "ZAMBIA"}},  
    {"country": {"country_id": "ZW", "country_name": "ZIMBABWE"}]} ] }
```

countries.json

Taxonomy



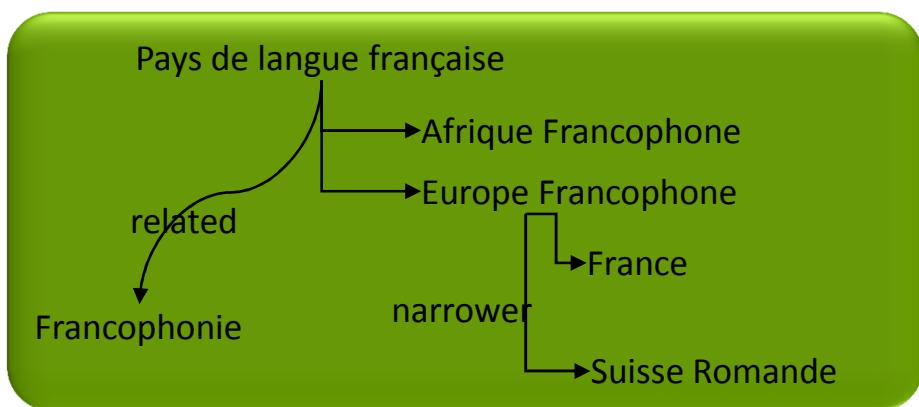
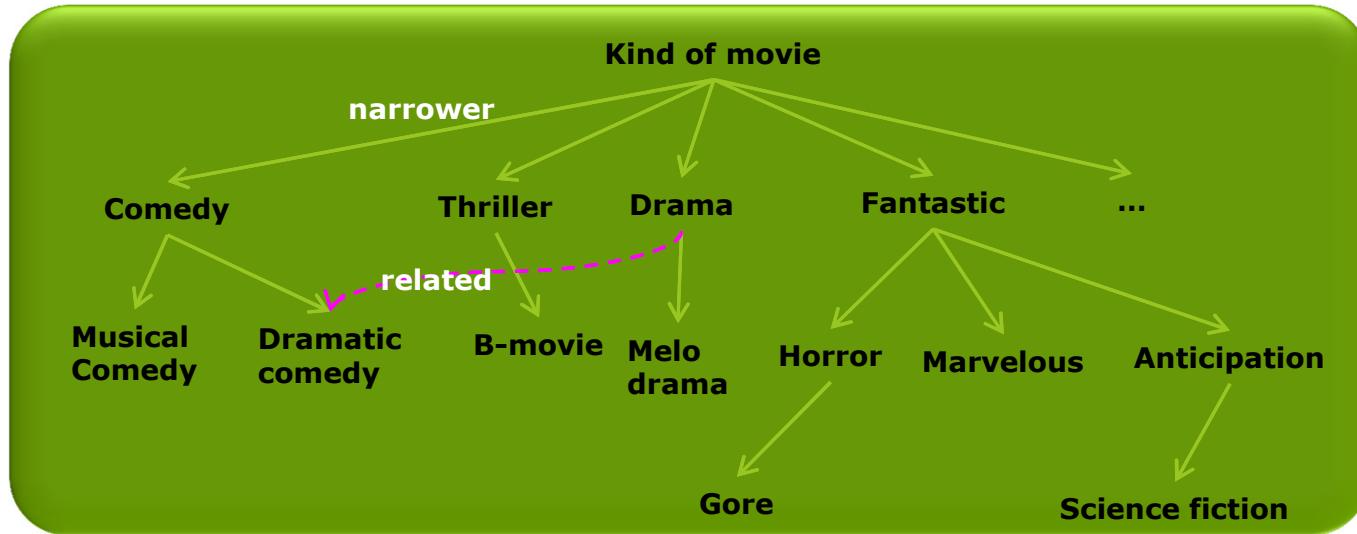


Taxonomy

A taxonomy is a classification method in a hierarchical structured and scalable architecture of information.

Thesaurus





Thesaurus

A thesaurus is like a taxonomy with more information about each concept including alternative and preferred terms. It may contain links to related concepts.

Thesaurus

- It is a set of standard terms based on a hierarchical structure.
- The terms are conceptually organized and interconnected by semantic relations.
- The thesaurus is alphabetically organized, it forms an alphabetical list of standardized terms for
 - content analysis,
 - classification
 - and indexing documents of information
- In many cases, the thesaurus also proposes a definition of the terms used.

Relations

- **BT** : Broader Term
- **NT** : Narrower Term
- **TT** : Top Term
- **RT** : Related Term (other than BT, NT, TT, etc.)



UNESCO Thesaurus

Thesaurus sample

Hierarchical display

List of the seven major subject domains, broken down into microthesauri which allow you to gain a quick overview of the subject matter.

> Select a field > Then select a microthesaurus Search Need help?

> Select a field

1 Education
2 Science
Y 3 Culture
C 4 Social and human sciences
5 Information and communication
6 Politics, law and economics
7 Countries and country groupings

You can launch a query in the online catalogue (UNESDOC) directly from the thesaurus. Click on the number next to the desired descriptor to obtain the records indexed with that term in unesdoc/uneshib.

Hierarchical display

List of the seven major subject domains, broken down into microthesauri which allow you to gain a quick overview of the subject matter.

2 Science > Then select a microthesaurus Search Need help?

> Then select a microthesaurus

2.05 Scientific approach
2.10 Science and research management
2.15 Mathematics and statistics
2.20 Physical sciences
2.25 Chemical sciences
2.30 Space sciences
2.35 Earth sciences
2.40 Geography and oceanography
2.45 Meteorology
2.50 Hydrology
2.55 Environmental sciences and engineering
2.60 Pollution, disasters and safety
2.65 Natural resources
2.70 Biology
2.75 Natural sciences
2.80 Medical sciences
2.85 Pathology

Governing Bodies | Director-General | Secretariat
UNESDOC@unesco | Information centres | Library
Terms of use | Contact us | Help

Thesaurus sample

5 records found for: Space sciences Click on the [number] to display the records indexed with that descriptor in unesdoc/unesbib.	
Term:	Astronomical systems [21] MT 2.30 Space sciences UF Celestial bodies NT Cosmic matter [7] NT Galaxies [9] NT Interstellar space [11]NT2 Black holes [3] NT Meteorites [10] NT Solar system [24]NT2 Planets [16]NT3 Earth (planet) [74]NT2 Satellites [54]NT3 Moon [27]NT2 Sun [19]NT3 Solar activity [7] NT Stars [19]NT2 Quasars [3] NT Universe [67]
Term:	Astronomy [153] MT 2.30 Space sciences NT Astrophysics [50]NT2 Gravitation [24] NT Celestial mechanics [4] NT Cosmology [113]
Term:	Space [129] MT 2.30 Space sciences UF Outer space
Term:	Space sciences [55] MT 2.30 Space sciences
Term:	Time [108] MT 2.30 Space sciences

 United Nations Educational, Scientific and Cultural Organization

UNESCO Thesaurus

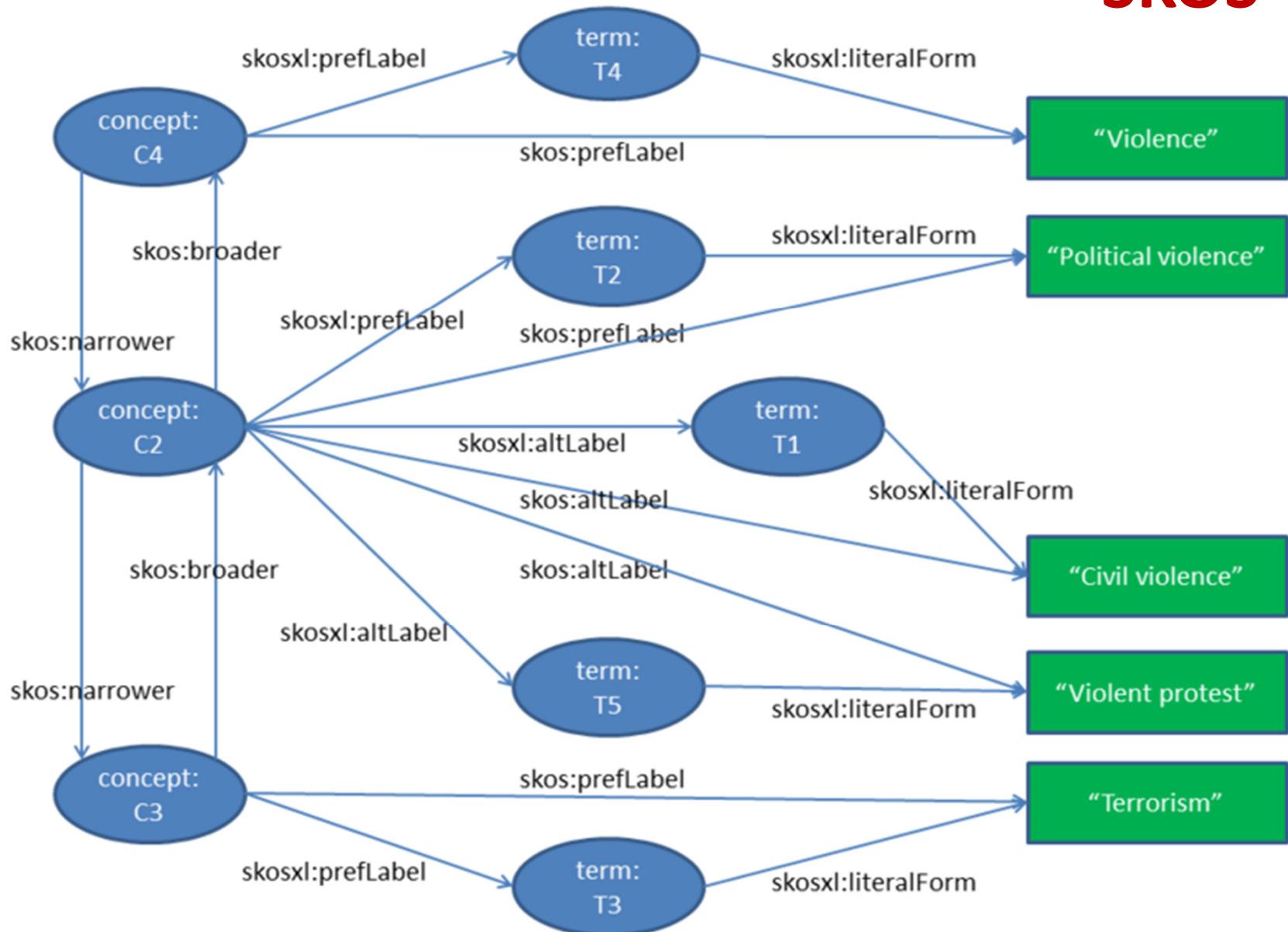
1 record found for: **Satellites\$**
Click on the [number] to display the records indexed with that descriptor in unesdoc/unesbib.

Term:	Satellites [54] Terme français: Satellite Término español: Satélite Русский термин : Спутники MT 2.30 Space sciences BT Solar system [24]BT2 Astronomical systems [21] NT Moon [27] RT Artificial satellites [90]
-------	---

[Print](#) [New query](#)

Une notice complète contient le microthésaurus (**MT**) auquel le terme appartient; la version anglaise, espagnole et russe; une note explicative le cas échéant (**NE**); les relations entre les termes (**EMP:EP; TG:TS; TA**)

SKOS



SKOS

```
@prefix rdfs: <http://www.w3.org/2000/01/rdf-schema#> .  
@prefix xsd: <http://www.w3.org/2001/XMLSchema#> .  
@prefix skos: <http://www.w3.org/2004/02/skos/core#> .  
@prefix skosxl: <http://www.w3.org/2008/05/skos-xl#> .  
@prefix dcterms: <http://purl.org/dc/terms/> .  
@prefix owl: <http://www.w3.org/2002/07/owl#> .  
@prefix con: <http://example.org/id/concept/polthes/> .  
@prefix term: <http://example.org/id/term/polthes/> .
```

```
term:T1 a skosxl:Label;  
skosxl:literalForm "Civil violence"@en.
```

```
con:C1 a skos:Concept;  
skos:prefLabel "Civil violence"@en;  
skosxl:prefLabel term:T1;  
skos:altLabel "Political violence"@en;  
skosxl:altLabel term:T2;  
skos:altLabel "Violent protest"@en;  
skosxl:altLabel term:T5;  
skos:broader con:C4;  
skos:narrower con:C3;  
dcterms:replaces con:C2 .
```

(...)

```
term:T3 a skosxl:Label;  
skosxl:literalForm "Terrorism"@en.
```

```
con:C4 a skos:Concept;  
skos:prefLabel "Violence"@en;  
skosxl:prefLabel term:C4;  
skos:narrower con:C1 .
```

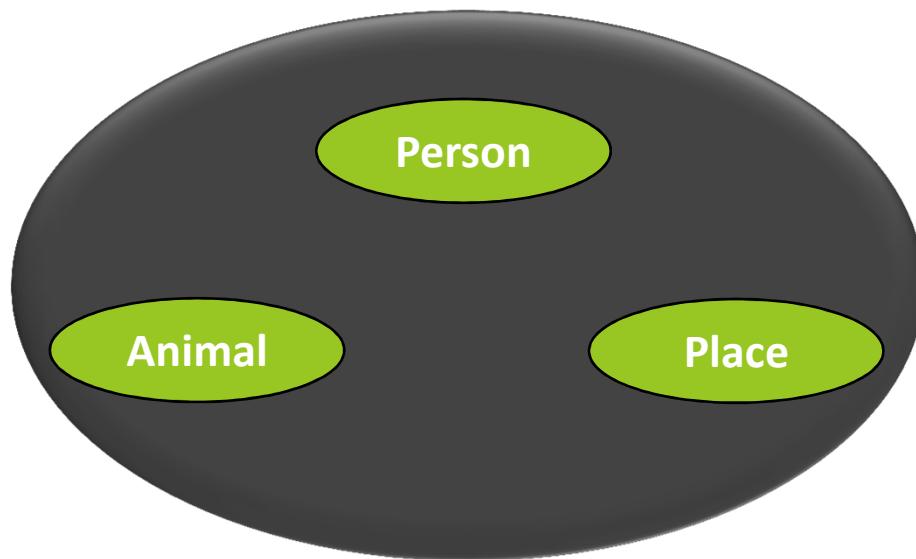
(...)

Ontology

Ontology

An ontology is a model of a domain of knowledge, describing the **concepts** within that domain

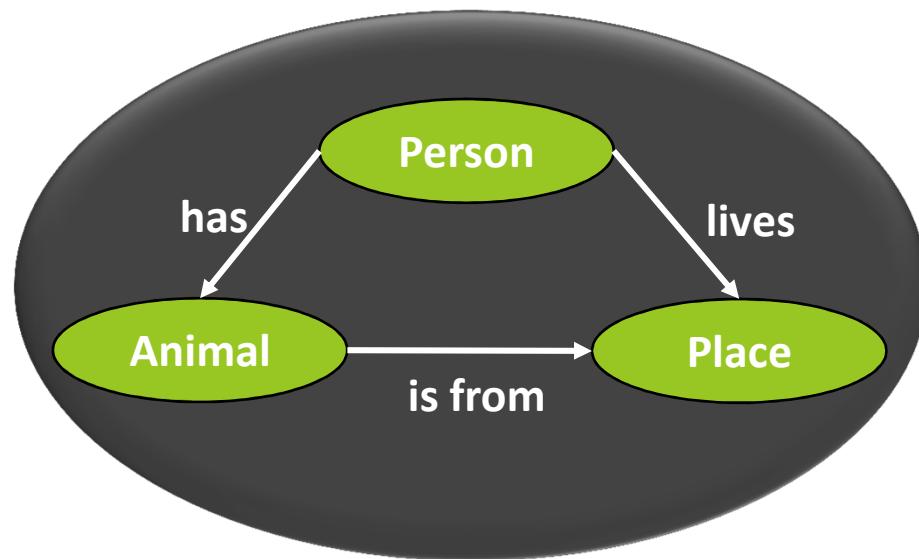
Ontologies are often also called vocabularies.



Ontology

An ontology is a model of a domain of knowledge, describing the **concepts** within that domain and the **relations** between concepts in the domain.

Ontologies are often also called vocabularies.



Studer et al (1998) distinguish four types of ontology.

Ontology

- **Representation ontologies**
 - don't belong to any field
 - organize the primitives of the logical theory (Charlet 2002)

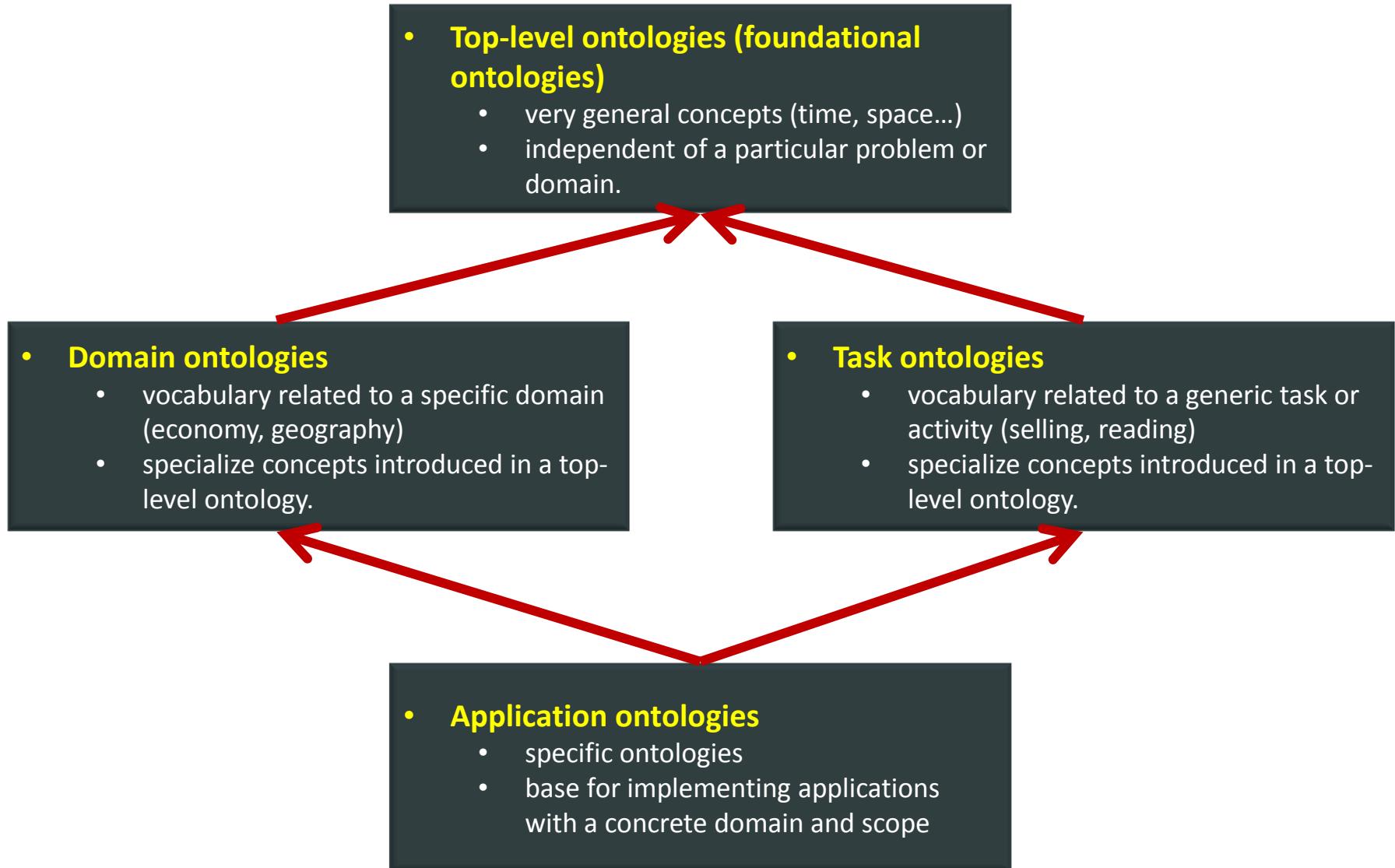
- **Generic ontologies (abstraction)**
 - abstract enough to be valid regardless of the field of study
 - upper ontology

- **Task ontologies (problem solving)**
 - provide concepts modeling a generic activity
 - explain the role played by each concept in reasoning

- **Application ontologies**
 - double specialisation : modelise a specific activity in a given domain

A classification system that uses the subject of conceptualization as a main criterion has been introduced by Guarino (1998).

Ontology



Ontology

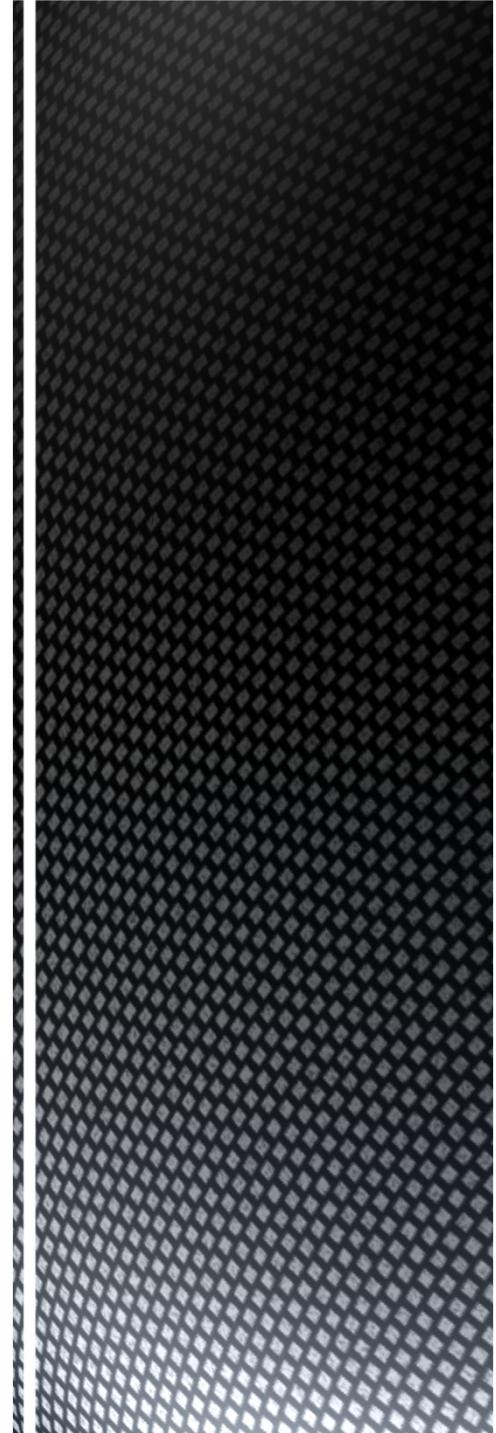
Concepts in relation

Ontology

- The term comes from philosophy
- A study of what *is*

Continuum : data, information, knowledge

The knowledge enables the production of new data, informations, knowledges : inference





2

The RDF semantic and its main syntaxes

1. [The basic principles of the Web of Data: from raw and heterogeneous Data to the lingua franca of the Web of Data](#)
2. [The RDF semantic and its main syntaxes](#)
3. [Handling and querying Data \(triple store, end point, SPARQL\)](#)
4. [Understanding the basics of knowledge modeling \(ontology, OWL\)](#)
5. [The lifting and linking of Data in order to publish 5 stars Data \(DataLift\).](#)

We can decompose...

- **Data**

- A base element, an atom, a quantity

16

- **Information**

- A data that agreements structure, for example a unit of measure

T= 16°C

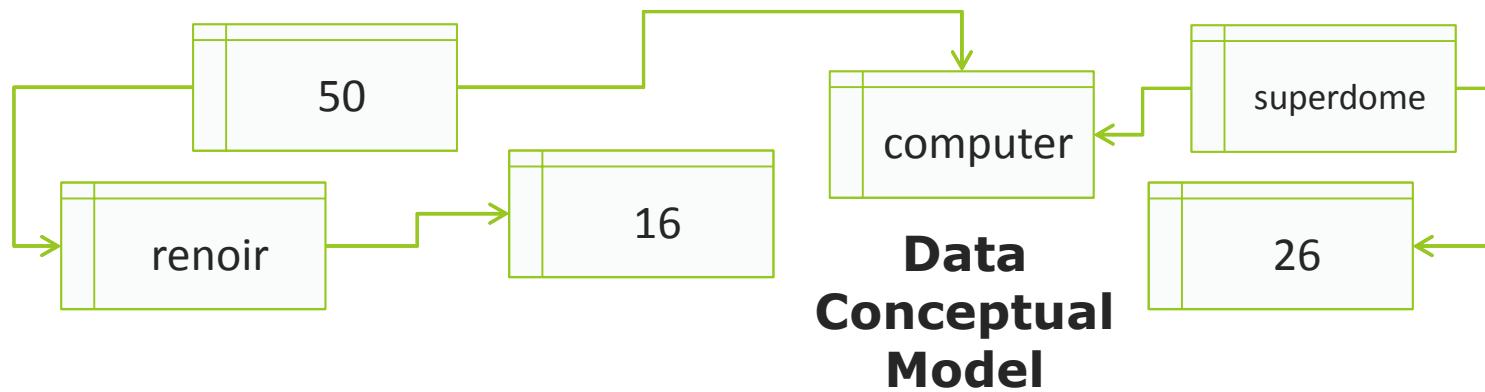
- **Knowledge**

- A (re)usable and interpretable information with context

The superdome 50 is in the Renoir room whose temperature is 16°C and the air conditioning is on, is an HP computer that the manufacturer guarantees to be operational until 26°C.

Rewriting complex knowledge

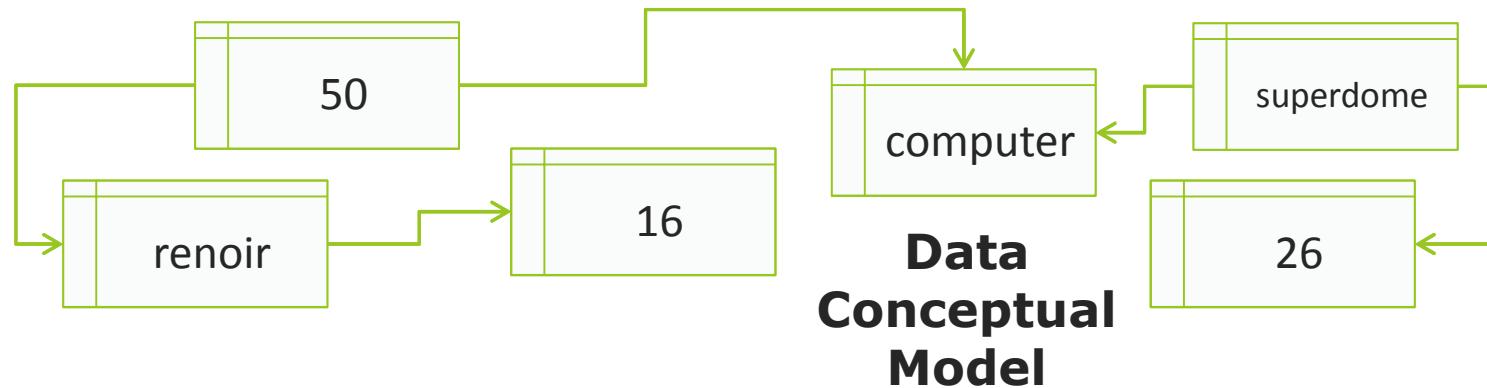
The superdome 50 is in the Renoir room whose temperature is 16°C and the air conditioning is on, is an HP computer that the manufacturer guarantees to be operational until 26°C.



Rewriting complex knowledge

The superdome 50 is in the Renoir room whose temperature is 16°C and the air conditioning is on, is an HP computer that the manufacturer guarantees to be operational until 26°C.

- The rewriting
 - **The superdome 50 is in the Renoir room**
 - **The Renoir room has the temperature 16°C**
 - **The air conditioning of the Renoir room is on**
 - **The superdome 50 is an HP computer**
 - **The superdome computer is guaranteed operational by HP until 26°C**



Rewriting complex knowledge

The superdome 50 is in the Renoir room whose temperature is 16°C and the air conditioning is on, is an HP computer that the manufacturer guarantees to be operational until 26°C.

- **The rewriting**
 - **The superdome 50 is in the Renoir room**
 - **The Renoir room has the temperature 16°C**
 - **The air conditioning of the Renoir room is on**
 - **The superdome 50 is an HP computer**
 - **The superdome computers are guaranteed operational by HP until 26°C**

These are regularly formed sentences

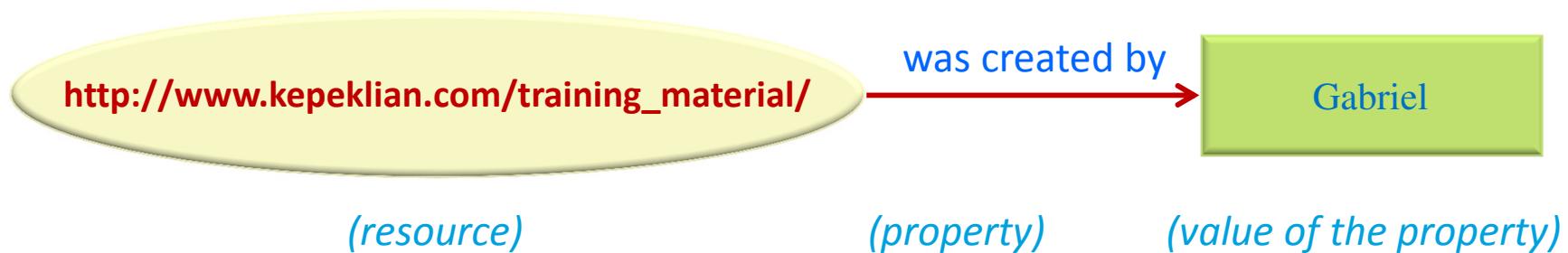
Subject	Predicat	Object
The superdome 50	is in	the Renoir room
The Renoir room	has the temperature of	16°C
The air conditioning of the Renoir room	is	on
The superdome 50	is	HP computer
The superdomes computers	are guaranteed operational by HP until	26°C

- A shared grammar is needed to ensure that metadata are interpreted consistently

- A framework to declare properties and relationships of items in the Web
- A basic model to declare resources:
 - **Resource**: anything with an URI
 - **Description**: states the properties of the resource using terms named by URIs
 - **Framework**: a common model or grammar for declarations
- Uses XML as a serialisation syntax

RDF =
**Resource
Description
Framework**

RDF = Resource Description Framework



```
<?xml version="1.0"?>
<RDF xmlns="http://www.w3.org/1999/02/22-rdf-syntax-ns#">
    <Description about="http://www.kepeklan.com/training\_material/">
        <Creator>Gabriel</Creator>
    </Description>
</RDF>
```

- ▶ That declaration can be read as:
 - *The training material identified by http://www.kepeklan.com/training_material/, was **created** by **Gabriel***



3

Handling and querying Data (triple store, end point, SPARQL)

1. The basic principles of the Web of Data: from raw and heterogeneous Data to the lingua franca of the Web of Data
2. The RDF semantic and its main syntaxes
3. Handling and querying Data (triple store, end point, SPARQL)
4. Understanding the basics of knowledge modeling (ontology, OWL)
5. The lifting and linking of Data in order to publish 5 stars Data (DataLift).

The limits of the CSV files



The data structure limit

Separation between the data structure and the data itself

idcity	label-fr	Label.de	label-en	longitude	latitude
1	Athènes	Athen	Athens	23.716667	37.966667
2	Pékin	Pekin	Beijing	116.400002	39.900002
3	Londres	London	London	-0.121137	51.515161

The data structure limit

Separation between the data structure and the data itself

idcity	label-fr	Label.de	label-en	longitude	latitude
1	Athènes	Athen	Athens	23.716667	37.966667
2	Pékin	Pekin	Beijing	116.400002	39.900002
3	Londres	London	London	-0.121137	51.515161

If we extract the data from the file, we have to extract also the structure to understand each data

The data structure limit

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2	Pékin	Pekin	Beijing	116.400002	39.900002
3	Londres	London	London	-0.121137	51.515161

23.716667

Alone :
no mining
no value

longitude

Alone :
no mining
no value

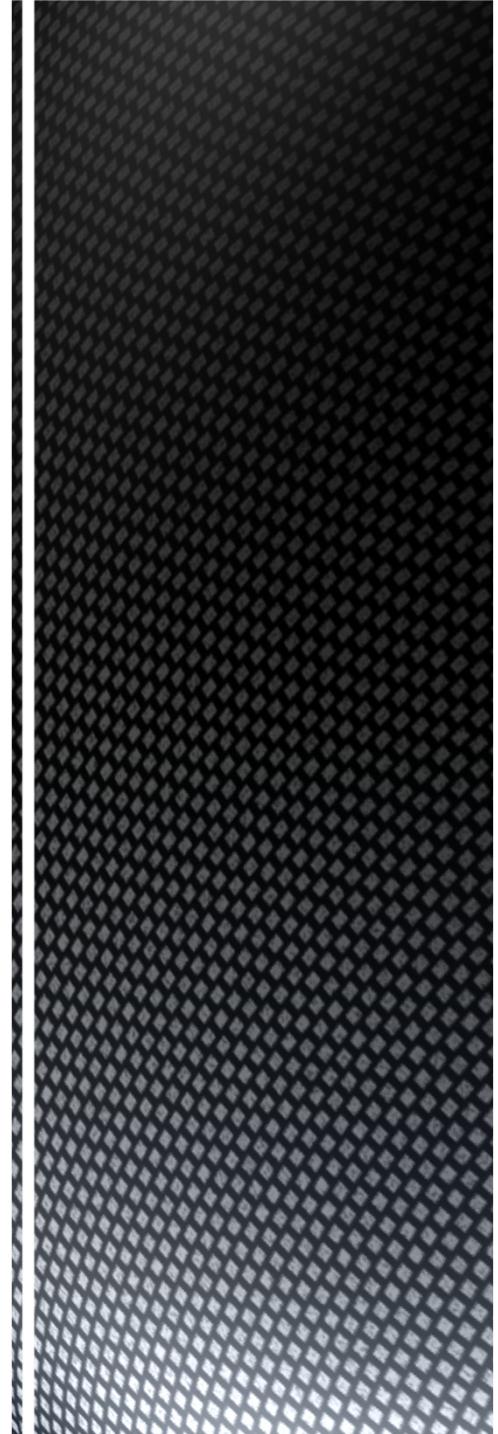
1

longitude

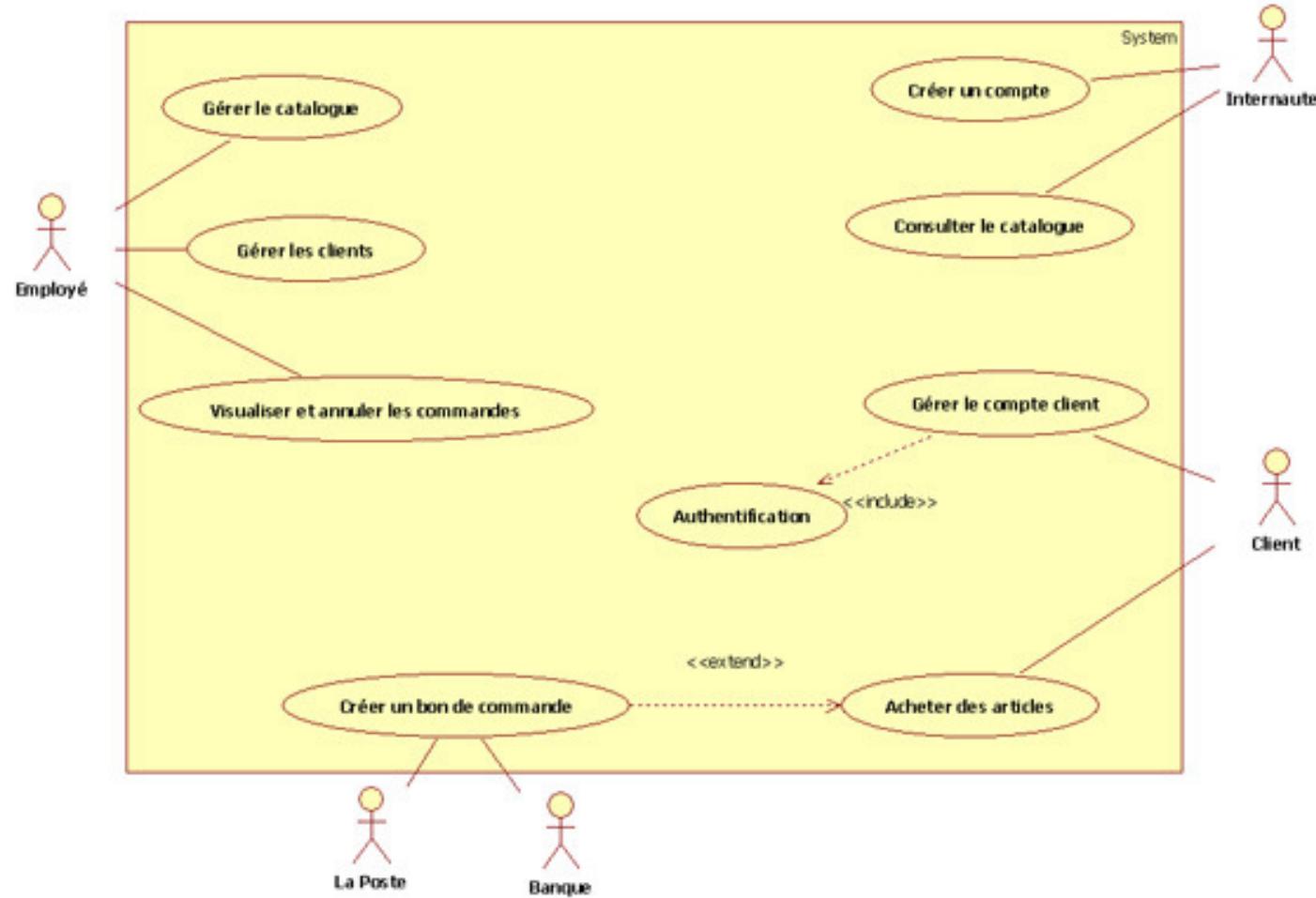
23.716667

If we extract the data from the file, we have to extract also the structure to understand each data

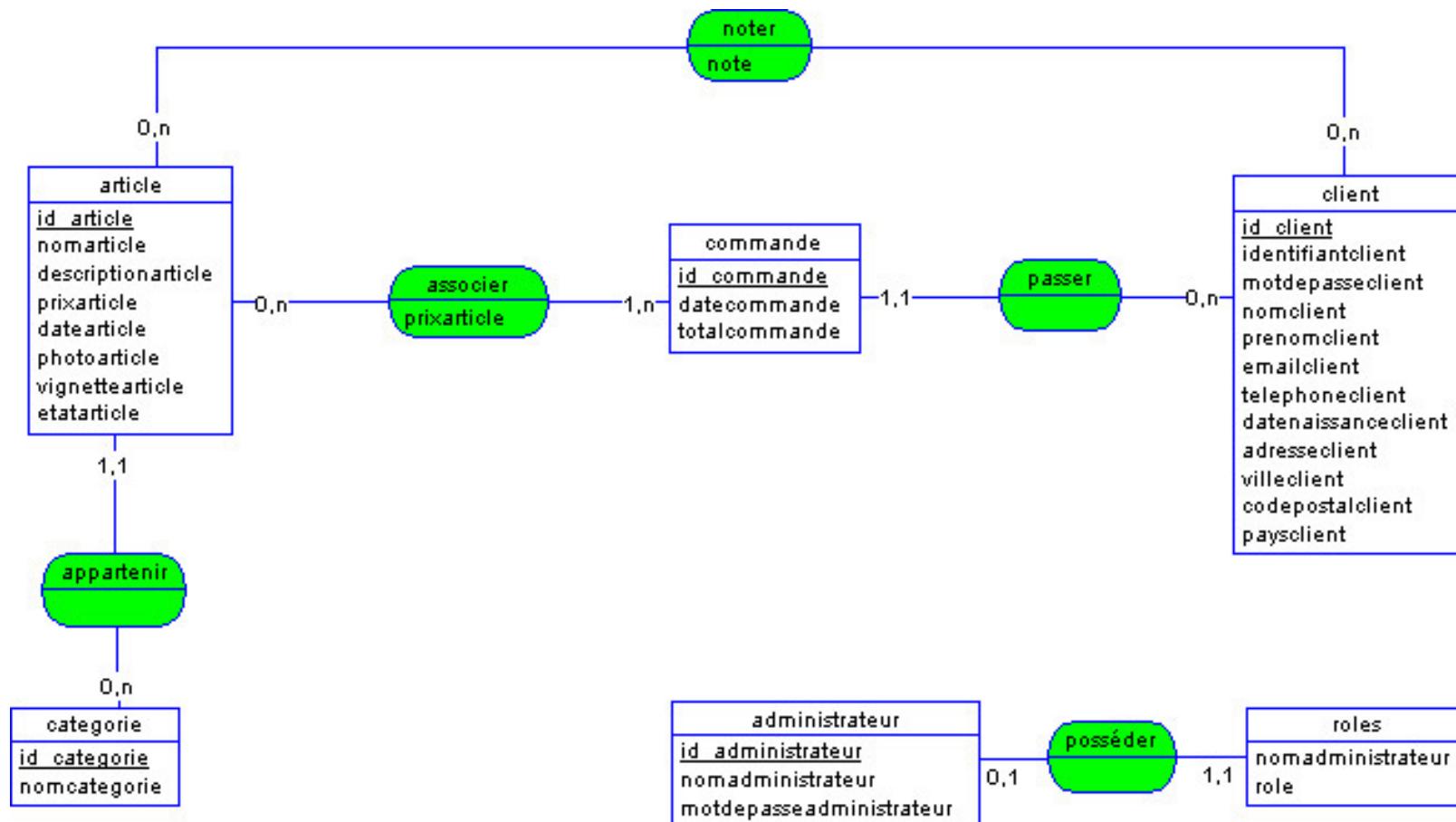
The limits of the relationnal model



UML



CDM



Here is a classic relational model describing the olympic games (a part of)

idcity	label-fr	label-en	longitude	latitude
1	Athènes	Athens	23.716667	37.966667
2	Pékin	Beijing	116.400002	39.900002

id	id city	year	opened Ceremony Date	closed Ceremony Date	id athletes Representant	athletes Number	motto
1	1	1896	1896/04/06	1896/04/15	NULL	241	NULL
2	2	2008	2008/08/08	2008/08/24	1	11028	One World one Dream

idathlete	label-fr	label-en	label-cn	birthdate	deathdate
1	Zhang Yining	Zhang Yining	張怡寧	1981/10/05	NULL
2	John Doe	John Doe	NULL	1982/07/08	NULL

We will now observe the induced problems

A - The relational limits – the data structure

1 – Separation between the data structure and the data itself

idcity	label-fr	label-en	longitude	latitude
1	Athènes	Athens	23.716667	37.966667
2	Pékin	Beijing	116.400002	39.900002

id	id city	year	opened Ceremony Date	closed Ceremony Date	id athletes Representant	athletes Number	motto
1	1	1896	1896/04/06	1896/04/15	NULL	241	NULL
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1	Zhang Yining	Zhang Yining	張怡寧	1981/10/05	NULL
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If we extract the data from the database, we have to extract also the structure to understand each data

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idcity	label-fr	label-en	longitude	latitude
1	Athènes	Athens	23.716667	37.966667
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id	id city	year	opened Ceremony Date	closed Ceremony Date	id athletes Representant	athletes Number	motto
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idathlete	label-fr	label-en	label-cn	birthdate	deathdate
1	Zhang Yining	Zhang Yining	張怡寧	1981/10/05	NULL
2	John Doe	John Doe	NULL	1982/07/08	NULL

A data applies to a record, because it is associated with a field.

The relationship is induced by the structure of the table.

A - The relational limits – the data structure

1 – Separation between the data structure and the data itself

idcity	label-fr	label-en	longitude	latitude
1	Athènes	Athens	23.716667	37.966667
2	Pékin	Beijing	116.400002	39.900002

id	id city	year	opened Ceremony Date	closed Ceremony Date	id athletes Representant	athletes Number	motto
1	1	1896	1896/04/06	1896/04/15	NULL	241	NULL
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idathlete	label-fr	label-en	label-cn	birthdate	deathdate
1	Zhang Yining	Zhang Yining	張怡寧	1981/10/05	NULL
2	John Doe	John Doe	NULL	1982/07/08	NULL

The data are not independent from each other.

They are understandable in the context of the database, of a record and of a field.

A - The relational limits – the data structure

2 – The structure of a database is rigid

idcity	label-fr	label-en	longitude	latitude
1	Athènes	Athens	23.716667	37.966667
2	Pékin	Beijing	116.400002	39.900002

id	id city	year	opened Ceremony Date	closed Ceremony Date	id athletes Representant	athletes Number	motto
1	1	1896	1896/04/06	1896/04/15	NULL	241	NULL
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2	2	2008	2008/08/08	2008/08/24	1	11028	One World one Dream

idathlete	label-fr	label-en	label-cn	birthdate	deathdate
1	Zhang Yining	Zhang Yining	張怡寧	1981/10/05	NULL
2	John Doe	John Doe	NULL	1982/07/08	NULL

If a data is missing for a field in a record, a "NULL" value is added but it is fictional.

A - The relational limits – the data structure

2 – The structure of a database is rigid

idcity	label-fr	label-en	longitude	latitude
1	Athènes	Athens	23.716667	37.966667
2	Pékin	Beijing	116.400002	39.900002

id	id city	year	opened Ceremony Date	closed Ceremony Date	id athletes Representant	Id atheletes repr. 2	athletes Number	motto
1	1	1896	1896/04/06	1896/04/15	NULL	NULL	241	NULL
2	2	2008	2008/08/08	2008/08/24	1	2	11028	One World one Dream

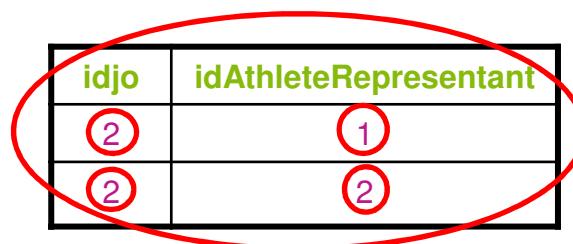
idathlete	label-fr	label-en	label-cn	birthdate	deathdate
1	Zhang Yining	Zhang Yining	張怡寧	1981/10/05	NULL
2	John Doe	John Doe	NULL	1982/07/08	NULL

If a data of the same nature is duplicate in a record, it is necessary to create another field or ...

A - The relational limits – the data structure

2 – The structure of a database is rigid

id	id city	year	opened Ceremony Date	closed Ceremony Date	id athletes Representant	Id atheletes repr. 2	athletes Number	motto
1	1	1896	1896/04/06	1896/04/15	NULL	NULL	241	NULL
2	2	2008	2008/08/08	2008/08/24	1	2	11028	One World one Dream



idjo	idAthleteRepresentant
2	1
2	2

idathlete	label-fr	label-en	label-cn	birthdate	deathdate
1	Zhang Yining	Zhang Yining	張怡寧	1981/10/05	NULL
2	John Doe	John Doe	NULL	1982/07/08	NULL

... or a table

A - The relational limits – the data structure

2 – The structure of a database is rigid

idcity	label-fr	label-en	label-cn	longitude	latitude
1	Athènes	Athens	NULL	23.716667	37.966667
2	Pékin	Beijing	北京市	116.400002	39.900002

id	id city	year	opened Ceremony Date	closed Ceremony Date	id athletes Representant	athletes Number	motto
1	1	1896	1896/04/06	1896/04/15	NULL	241	NULL
2	2	2008	2008/08/08	2008/08/24	1	11028	One World one Dream

idathlete	label-fr	label-en	label-cn	birthdate	deathdate
1	Zhang Yining	Zhang Yining	張怡寧	1981/10/05	NULL
2	John Doe	John Doe	NULL	1982/07/08	NULL

To manage multilingualism, we have to create fields for each language, while the meaning of the field is exactly the same, or create a specific table.

A - The relational limits – the data structure

3 – The relations between two tables are induced

idcity	label-fr	label-en	longitude	latitude
1	Athènes	Athens	23.716667	37.966667
2	Pékin	Beijing	116.400002	39.900002

id	id city	year	opened Ceremony Date	closed Ceremony Date	id athletes Representant	athletes Number	motto
1	1	1896	1896/04/06	1896/04/15	NULL	241	NULL
2	2	2008	2008/08/08	2008/08/24	1	11028	One World one Dream

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1	Zhang Yining	Zhang Yining	張怡寧	1981/10/05	NULL
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2	Pékin	Beijing	116.400002	39.900002

id	id city	year	opened Ceremony Date	closed Ceremony Date	id athletes Representant	athletes Number	motto
1	1	1896	1896/04/06	1896/04/15	NULL	241	NULL
2	2	2008	2008/08/08	2008/08/24	1	11028	One World one Dream

idathlete	label-fr	label-en	label-cn	birthdate	deathdate
1	Zhang Yining	Zhang Yining	張怡寧	1981/10/05	NULL
2	John Doe	John Doe	NULL	1982/07/08	NULL

The relation between data from two tables is induced by the use of common identifiers called foreign keys. The nature of the relationship is not clearly expressed either in the structure or in the data.

A - The relational limits – the data structure

3 – The relations between two tables are induced

idcity	label-fr	label-en	longitude	latitude
1	Athènes	Athens	23.716667	37.966667
2	Pékin	Beijing	116.400002	39.900002

id	id city	year	opened Ceremony Date	closed Ceremony Date	id athletes Representant	athletes Number	motto
1	1	1896	1896/04/06	1896/04/15	NULL	241	NULL
2	2	2008	2008/08/08	2008/08/24	1	11028	One World one Dream

idathlete	label-fr	label-en	label-cn	birthdate	deathdate
1	Zhang Yining	Zhang Yining	張怡寧	1981/10/05	NULL
2	John Doe	John Doe	NULL	1982/07/08	NULL

The extraction of a database does not highlight these relations. We have to retrieve data from different tables to maintain the relationship.

A - The relational limits – the data structure

4 – The identifier of a record is a data like all other

idcity	label-fr	label-en	longitude	latitude
1	Athènes	Athens	23.716667	37.966667
2	Pékin	Beijing	116.400002	39.900002

id	id city	year	opened Ceremony Date	closed Ceremony Date	id athletes Representant	athletes Number	motto
1	1	1896	1896/04/06	1896/04/15	NULL	241	NULL
2	2	2008	2008/08/08	2008/08/24	1	11028	One World one Dream

idathlete	label-fr	label-en	label-cn	birthdate	deathdate
1	Zhang Yining	Zhang Yining	張怡寧	1981/10/05	NULL
2	John Doe	John Doe	NULL	1982/07/08	NULL

A - The relational limits – the data structure

4 – The identifier of a record is a data like all other

idcity	label-fr	label-en	longitude	latitude
1	Athènes	Athens	23.716667	37.966667
2	Pékin	Beijing	116.400002	39.900002

id	id city	year	opened Ceremony Date	closed Ceremony Date	id athletes Representant	athletes Number	motto
1	1	1896	1896/04/06	1896/04/15	NULL	241	NULL
2	2	2008	2008/08/08	2008/08/24	1	11028	One World one Dream

idathlete	label-fr	label-en	label-cn	birthdate	deathdate
1	Zhang Yining	Zhang Yining	張怡寧	1981/10/05	NULL
2	John Doe	John Doe	NULL	1982/07/08	NULL

The identifier of a record does not have a standardized form.
It is independent from the database or from the table (so from the structure).

B - The relational limits - Data interoperability

1 - The identifiers are local and specific to a database

idcity	label-fr	label-en					
1	Athènes	Athens					
2	Pékin	Beijing					
id	id city	year	opened Ceremony Date	closed Ceremony Date	id athletes Representant	athletes Number	motto
1	1	1896	1896/04/06	1896/04/15	NULL	241	NULL
2	2	2008	2008/08/08	2008/08/24	1	11028	One World one Dream
idathlete	label-fr	label-en	label-cn	birthdate	deathdate		
1	Zhang Yining	Zhang Yining	張怡寧	1981/10/05	NULL		
2	John Doe	John Doe	NULL	1982/07/08	NULL		

It is not possible to identify natively two equivalent resources between two different databases.

B - The relational limits - Data interoperability

2 - The names of fields are specific to a database

idcity	label-fr	label-en
1	Athènes	Athens
2	Pékin	Beijing

idcity	Name	longitude	latitude
A	Athènes	23.716667	37.966667
B	Pékin	116.400002	39.900002

id	id city	year	opened Ceremony Date	closed Ceremony Date	id athletes Representant	athletes Number	motto
1	1	1896	1896/04/06	1896/04/15	NULL	241	NULL
2	2	2008	2008/08/08	2008/08/24	1	11028	One World one Dream

idathlete	label-fr	label-en	label-cn	birthdate	deathdate
1	Zhang Yining	Zhang Yining	張怡寧	1981/10/05	NULL
2	John Doe	John Doe	NULL	1982/07/08	NULL

The structure of a database is local. There are no standards for naming properties and to assign them to a normalization of a particular data type.

B - The relational limits - Data interoperability

3 - The database structure is not based on any inheritance mechanism

idcity	label-fr	label-en
1	Athènes	Athens
2	Pékin	Beijing

idcity	Name	longitude	latitude
A	Athènes	23.716667	37.966667
B	Pékin	116.400002	39.900002

id	id city	year	opened Ceremony Date	closed Ceremony Date	id athletes Representant	athletes Number	motto
1	1	1896	1896/04/06	1896/04/15	NULL	241	NULL
2	2	2008	2008/08/08	2008/08/24	1	11028	One World one Dream

idathlete	label-fr	label-en	label-cn	birthdate	deathdate
1	Zhang Yining	Zhang Yining	張怡寧	1981/10/05	NULL
2	John Doe	John Doe	NULL	1982/07/08	NULL

Is that table a specialisation of the description of a person?

Is that table a specialisation of the description of an event?

It is not possible to connect a table to a generic model of local or external description which may inherit the characteristics. It requires to construct a CDM from scratch almost.

B - The relational limits - Data interoperability

4 - There is no standard representation for exchanging RDB on a network.

idcity	label-fr	label-en
1	Athènes	Athens
2	Pékin	Beijing

idcity	Name	longitude	latitude
A	Athènes	23.716667	37.966667
B	Pékin	116.400002	39.900002

id	id city	year	opened Ceremony Date	closed Ceremony Date	id athletes Representant	athletes Number	motto
1	1	1896	1896/04/06	1896/04/15	NULL	241	NULL
2	2	2008	2008/08/08	2008/08/24	1	11028	One World one Dream

idathlete	label-fr	label-en	label-cn	birthdate	deathdate
1	Zhang Yining	Zhang Yining	張怡寧	1981/10/05	NULL
2	John Doe	John Doe	NULL	1982/07/08	NULL

The extraction of a database is specific for each database, there is no standard syntax for exchanging data in a database and merge with another base.

B - The relational limits - Data interoperability

5 - There is no standard way to query a BDR directly on the Web.

idcity	label-fr	label-en
1	Athènes	Athens
2	Pékin	Beijing

idcity	Name	longitude	latitude
A	Athènes	23.716667	37.966667
B	Pékin	116.400002	39.900002

id	id city	year	opened Ceremony Date	closed Ceremony Date	id athletes Representant	athletes Number	motto
1	1	1896	1896/04/06	1896/04/15	NULL	241	NULL
2	2	2008	2008/08/08	2008/08/24	1	11028	One World one Dream

idathlete	label-fr	label-en	label-cn	birthdate	deathdate
1	Zhang Yining	Zhang Yining	張怡寧	1981/10/05	NULL
2	John Doe	John Doe	NULL	1982/07/08	NULL

The extraction of a database is specific for each database, there is no standard syntax for exchanging data in a database and merge with another base.

The RDB is no more the good technology

- We have to handle more and more data
- We have to handle more and more heterogeneous data
- We have to handle more and more faster
- We have to handle more and more in a non intrusive manner
- How do we best handle data in such a context?

The official semantic web activity page says: "The Semantic Web provides a common framework that allows data to be shared and reused across application, enterprise, and community boundaries".

Triple Stores

Three architecture principles

In-Memory
Stores

Temporary storage of triples in memory

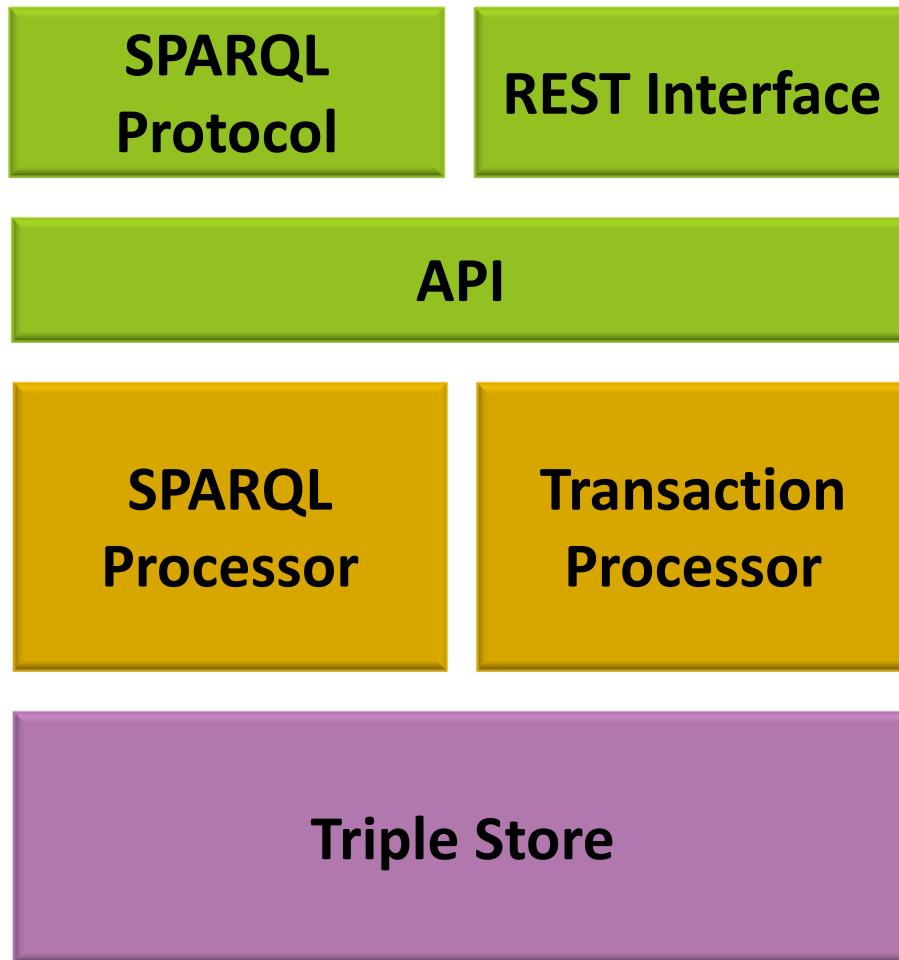
Native
Stores

Persistent storage of triples in native stores with their own storage implementation

Non-Native
Stores

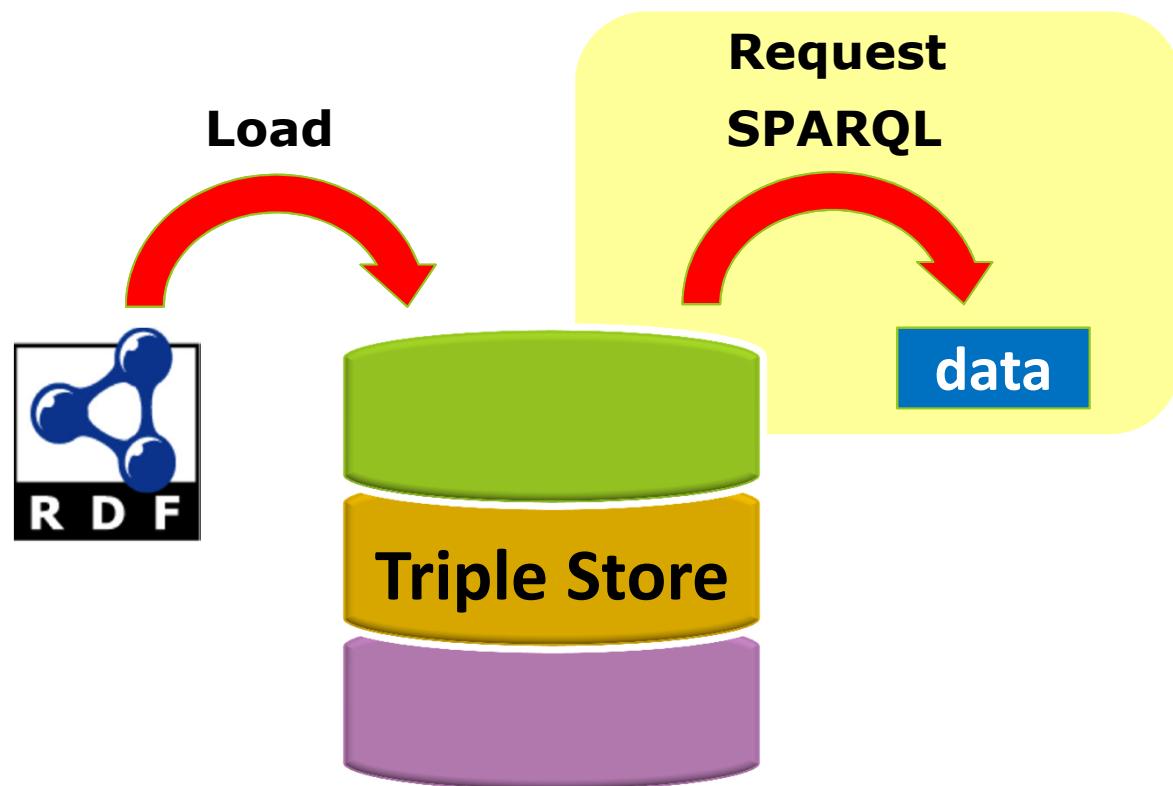
Persistent storage of triples on top of third-party databases

Triple Stores



- select, ask : query
- construct
- describe : export graph
- Insert, delete

SPARQL endpoint



SPARQL Endpoints Status

<http://sparqles.okfn.org/>



The screenshot shows a dashboard with four main sections: Availability (Up / Down), Performance (Cold / Warm), Interoperability (SPARQL 1.0 / 1.1), and Discoverability (VoID, Server). A search bar indicates 469 endpoints.

AVAILABILITY	PERFORMANCE	INTEROPERABILITY	DISCOVERABILITY
Up / Down	Cold / Warm	SPARQL 1.0 / 1.1	VoID, Server

SEARCH: 469 endpoints

- SPARQL Endpoints Status tool monitors
 - the availability,
 - the performance,
 - the interoperability and
 - the discoverability
- of SPARQL Endpoints registered in Datahub.

```
SELECT ?o WHERE { ?s dbpedia2:blackboard ?o. }  
SELECT ?o WHERE { ?s dbpedia2:blackboard ?o. }
```



SPARQL



SPARQL Protocol and RDF Query Language

Query, Access, Transformation

Data in RDF

A query language

An access protocol

An XML format for the results



Introduction *requesting the RDF data*

- **RDF** (Resource Description Framework)
 - a flexible and extensible way to represent the information of the Web resources
 - a graph model intended to describe formally Web resources and their metadata, and to allow automatic processing of such descriptions.
- A document is structured as a set of RDF triples.
- An RDF triple is an association (subject, predicate, object)
 - The **subject** is the resource being described
 - The **predicate** is a type of property applicable to this resource
 - The **object** represents a data or another resource : it is the value of the property.

Pattern Matching



- SPARQL (SPARQL Protocol and RDF Query Language)

- Only 4 reading possibilities

- **SELECT** query
 - Extracts the data from a SPARQL endpoint, the results are in an **array**.
- **ASK** query
 - Requests a SPARQL endpoint to obtain a **boolean answer**.
- **CONSTRUCT** query
 - Extracts the data from a SPARQL endpoint in order to transform them in **triples** (RDF format).
- **DESCRIBE** query
 - Use to extract a part of the RDF **graph** in a silo. Returns all triples containing all URIs that meet the requirements of pattern matching specified in the query. URIs can be present as subject or object of the triples.



W3C : 2008,
SPARQL 1.0

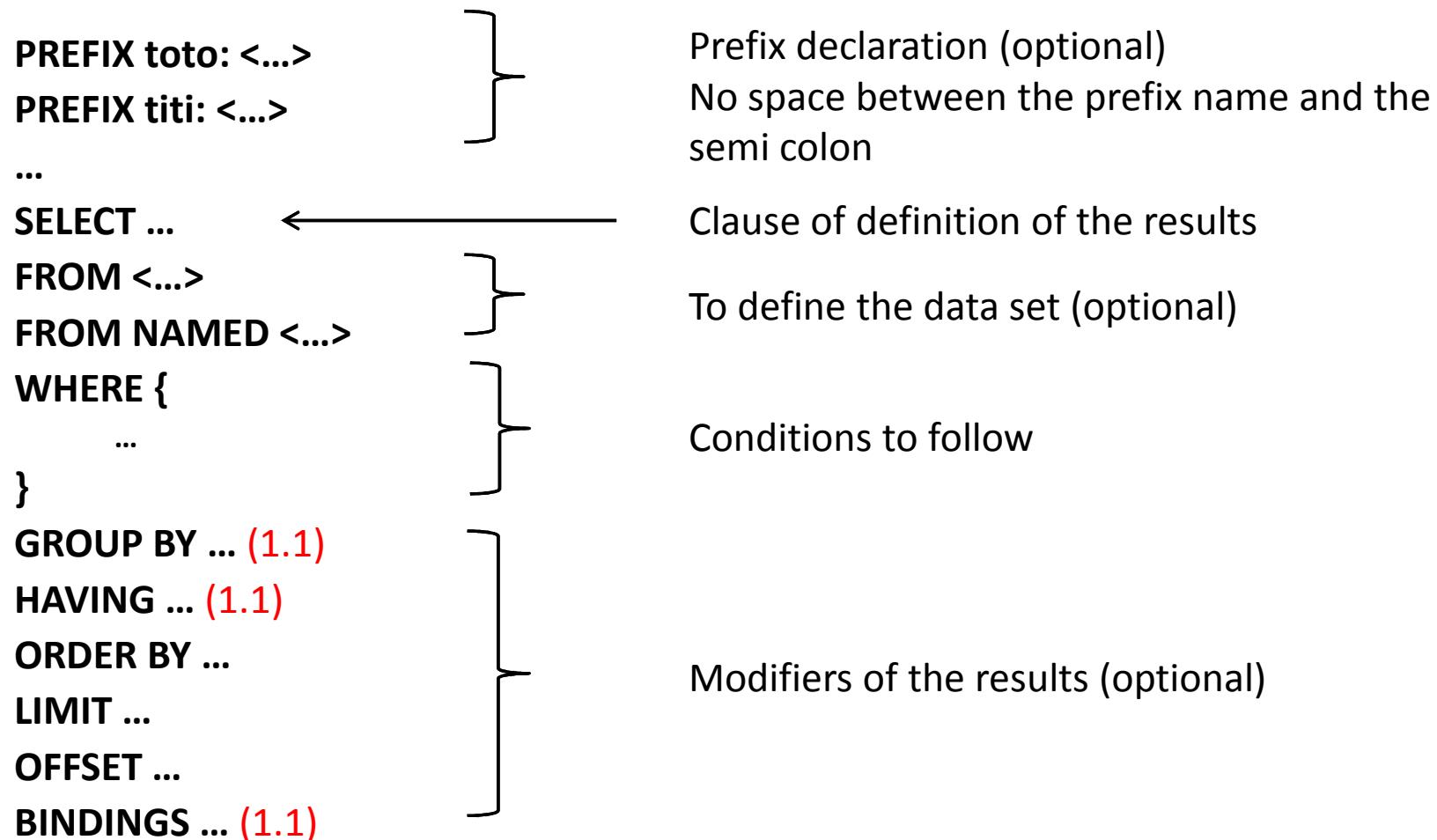


- SPARQL 1.1 looks more and more SQL
- It allows read operations and write data without the need to use or know specific languages: RDF or RDFS, OWL ...

- Two new functions:
 - insert Data
 - Insertion of new triplets
 - Delete Data
 - Remove selected triplets

W3C : 2013,
SPARQL 1.1

SELECT, anatomy of a query



I want to know the URI and the name of the persons who know Jean Dupond, and I search in his FOAF profile.

Jean Dupond knows persons.
These persons are named ???.

```
PREFIX foaf: <http://xmlns.com/foaf/0.1/>

SELECT ?persons ?name

FROM <http://www.jean_dupond.com/foaf.rdf>

WHERE {

    <http://www.jean_dupond.com/foaf.rdf/i> foaf:knows ?persons.
    ?persons foaf:name ?name.

}
```

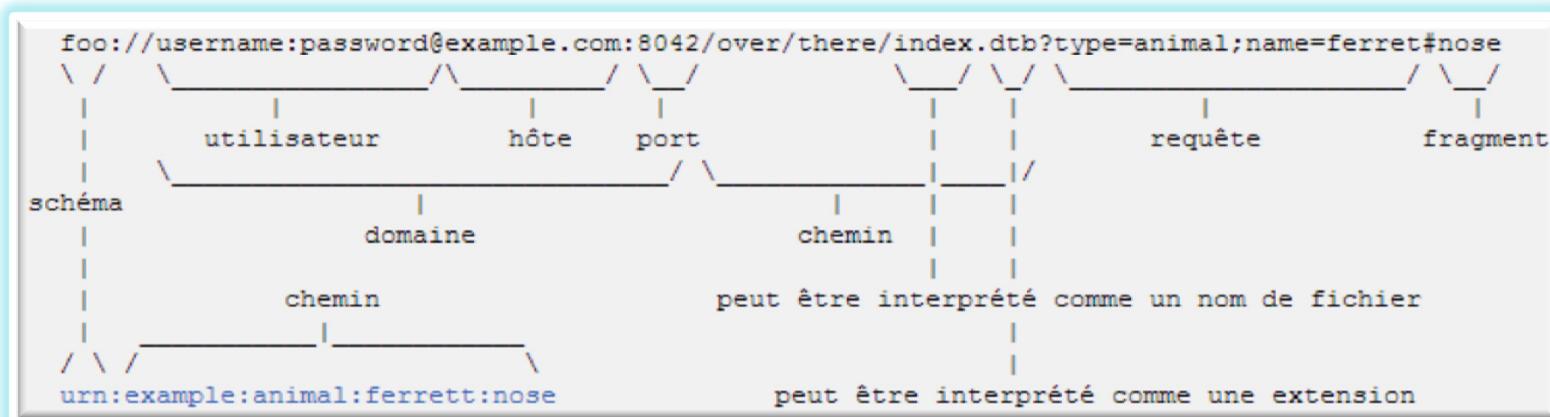
URI

RFC 4935

When we need to name the things

- ftp://ftp.is.co.za/rfc/rfc1808.txt
- <http://www.dila.premier-ministre.gouv.fr/qui-sommes-nous/dila.html>
- ldap://[2001:db8::7]/c=GB?objectClass?one
- <mailto:thomas.saint-aubin@dila.gouv.fr>
- news:comp.infosystems.www.servers.unix
- <tel:+1-816-555-1212>
- telnet://192.0.2.16:80/
- urn:oasis:names:specification:docbook:dtd:xml:4.1.2

zoom
zoom



PREFIX

- **PREFIX... they rewrite**

```
PREFIX rdf: <http://xmlns.com/foaf/0.1/>
PREFIX rdfs: <http://www.w3.org/2000/01/rdf-schema#>
PREFIX owl: <http://www.w3.org/2002/07/owl#>
PREFIX xsd: <http://www.w3.org/2001/XMLSchema#>
PREFIX dc: <http://purl.org/dc/elements/1.1/>
PREFIX foaf: <http://xmlns.com/foaf/0.1/>
```

- **Where can we find others?**

LITERALS

- *Plain literals*
 - “un exemple de littéral”
- *Plain literals with the mentions of a country*
 - “bonjour les amis”@fr
- *Typed literals*
 - “123”^^xsd:integer
- *Shortcuts*
 - **true** → “true”^^xsd:boolean
 - **3** → “3”^^xsd:integer
 - **4.2** → “4.2”^^xsd:decimal
 - **30/06/2014** → “2014-06-30”^^xsd:date
 - **30/06/2014 12:00** → “2014-06-30T12:00:00+02:00”^^xsd:datetime

VARIABLES

- *Variables:*
 - **?var1 ?encoreUneVar \$etUneAutre**
- Never start with a digit
- Case sensitive
- Without space
- Meaningful because they are used to name columns
(SPARQL 1.1 introduces AS to give the possibility to change the column name)

FILTER

- To complete the way to find triples according to a given search, SPARQL offers a mechanism of filters.
- A filter is expressed in a line placed in the mask of the graph.
- Example:

```
?s1 ?p1 ?o1 .  
?s2 ?p2 ?o2 .  
FILTER ( expression )
```

- The filter removes the solutions where the expression is false.

FROM

FROM <url> is used to identify the content in the default graph.

```
SELECT (COUNT(*) as ?n)
FROM <http://localhost:9091/project/kiosques/source/kiosques-ouverts-csv-rdf-1>
WHERE { ?s ?p ?o }
```

It is equivalent to the following request in a named graph.

```
PREFIX g: <http://localhost:9091/project/kiosques/source/kiosques-ouverts-csv-rdf-1>
SELECT (COUNT(*) as ?n)
WHERE { GRAPH g: { ?s ?p ?o } }
```

FROM NAMED

```
<http://grapha.com> = { <a1> <p> <a2> . }
<http://graphb.com> = { <b1> <p> <b2> . }
<http://graphc.com> = { <c1> <p> <c2> . }
<http://graphd.com> = { <d1> <p> <d2> . }
```

```
SELECT ?s WHERE { ?s <p> ?o }
→ <a1>, <b1> <c1>, <d1>
```

```
FROM <http://grapha.com>
SELECT ?s WHERE { ?s <p> ?o }
→ <a1>.
```

```
FROM NAMED <http://grapha.com>
SELECT ?s WHERE { ?s <p> ?o }
→ rien
```

```
FROM <http://grapha.com>
FROM <http://graphb.com>
FROM NAMED <http://graphc.com>
FROM NAMED <http://graphd.com>
SELECT ?s WHERE { ?s <p> ?o }
→ <a1>, <b1>.
```

```
FROM <http://grapha.com>
FROM <http://graphb.com>
FROM NAMED <http://graphc.com>
FROM NAMED <http://graphd.com>
SELECT ?s WHERE { GRAPH ?g { ?s <p> ?o } }
→ <c1>, <d1>.
```

THE COMPONENTS OF A QUERY

■ Prolog

- PREFIX Definition

```
PREFIX g: <http://localhost:9091/project/kiosques/source/kiosques-ouverts-csv-rdf-1>
SELECT *
WHERE { GRAPH g: { ?s ?p ?o } }
```

■ Result

- SELECT, DESCRIBE, CONSTRUCT or ASK

```
PREFIX g: <http://localhost:9091/project/kiosques/source/kiosques-ouverts-csv-rdf-1>
SELECT *
WHERE { GRAPH g: { ?s ?p ?o } }
```

■ Specification of the data set

```
SELECT *
FROM <http://localhost:9091/project/kiosques/source/kiosques-ouverts-csv-rdf-1>
WHERE { ?s ?p ?o }
```

THE COMPONENTS OF A QUERY

■ Pattern

- The things we are looking for

```
PREFIX g: <http://localhost:9091/project/kiosques/source/kiosques-ouverts-csv-rdf-1>
SELECT ?s
WHERE { GRAPH g: { ?s ?p ?o } }
```

■ Modifiers

- ORDER BY, LIMIT, OFFSET

```
PREFIX g: <http://localhost:9091/project/kiosques/source/kiosques-ouverts-csv-rdf-1>
SELECT *
WHERE { GRAPH g: { ?s ?p ?o } }
ORDER BY ?s
OFFSET 20
LIMIT 50
```

BIND

Data

```
@prefix dc:  
  <http://purl.org/dc/elements/1.1/> .  
@prefix :      <http://example.org/book/>  
. .  
@prefix ns:    <http://example.org/ns#> .  
  
:book1  dc:title      "SPARQL Tutorial" .  
:book1  ns:price      42 .  
:book1  ns:discount   0.2 .  
  
:book2  dc:title      "The Semantic Web" .  
:book2  ns:price      23 .  
:book2  ns:discount   0.25 .
```

Query

```
PREFIX dc:  
  <http://purl.org/dc/elements/1.1/>  
PREFIX ns:  <http://example.org/ns#>  
  
SELECT ?title ?price  
{ ?x ns:price ?p .  
  ?x ns:discount ?discount  
  BIND (?p*(1-?discount) AS ?price)  
  FILTER(?price < 20)  
  ?x dc:title ?title .  
}
```

Another query

```
PREFIX dc:  
  <http://purl.org/dc/elements/1.1/>  
PREFIX ns:  <http://example.org/ns#>  
  
SELECT ?title ?price  
{ { ?x ns:price ?p .  
    ?x ns:discount ?discount  
    BIND (?p*(1-?discount) AS ?price)  
  }  
  { ?x dc:title ?title . }  
  FILTER(?price < 20)  
}
```

title	price
"The Semantic Web"	17.25

AGGREGATION

Data

```
@prefix : <http://books.example/> .  
  
:org1 :affiliates :auth1, :auth2 .  
:auth1 :writesBook :book1, :book2 .  
:book1 :price 9 .  
:book2 :price 5 .  
:auth2 :writesBook :book3 .  
:book3 :price 7 .  
:org2 :affiliates :auth3 .  
:auth3 :writesBook :book4 .  
:book4 :price 7 .
```

Query

```
PREFIX : <http://books.example/>  
SELECT (SUM(?lprice) AS ?totalPrice)  
WHERE {  
    ?org :affiliates ?auth .  
    ?auth :writesBook ?book .  
    ?book :price ?lprice .  
}  
GROUP BY ?org  
HAVING (SUM(?lprice) > 10)
```

totalPrice

21

GROUP BY

- **GROUP BY** : In order to calculate total values for a solution, the solution is first divided into one or more groups, and the total value is calculated for each group.

Query

```
SELECT (AVG(?y) AS ?avg)
WHERE {
    ?a :x ?x ;
    :y ?y .
}
GROUP BY ?x
```

- ▶ **NB** : AVG, the average function.

SUB REQUEST

Data

```
@prefix : <http://people.example/> .  
:alice :name "Alice", "Alice Foo", "A. Foo" .  
:alice :knows :bob, :carol .  
:bob :name "Bob", "Bob Bar", "B. Bar" .  
:carol :name "Carol", "Carol Baz", "C. Baz" .
```

Query

```
PREFIX : <http://people.example/>  
PREFIX : <http://people.example/>  
SELECT ?y ?minName  
WHERE { :alice :knows ?y .  
       { SELECT ?y (MIN(?name) AS ?minName)  
         WHERE { ?y :name ?name .  
                  } GROUP BY ?y  
       }  
}
```

Y	minName
:bob	"B. Bar"
:carol	"C. Baz"

This is achieved by first evaluating the internal request

```
SELECT ?y (MIN(?name) AS ?minName)  
WHERE {  
       ?y :name ?name .  
     } GROUP BY ?y
```

We obtain the following solutions

Y	minName
:alice	"A. Foo"
:bob	"B. Bar"
:carol	"C. Baz"

that are then binded with the results of the external query

Y
:bob
:carol

Return a name (the one with the lowest sort order) for all the people that know Alice and have a name.

DUPLICATE RESULTS

Data

```
@prefix foaf: <http://xmlns.com/foaf/0.1/> .  
  
_:x    foaf:name    "Alice" .  
_:x    foaf:mbox    <mailto:alice@example.com> .  
  
_:y    foaf:name    "Alice" .  
_:y    foaf:mbox    <mailto:asmith@example.com> .  
  
_:z    foaf:name    "Alice" .  
_:z    foaf:mbox    <mailto:alice.smith@example.com>  
.
```

Query

```
PREFIX foaf: <http://xmlns.com/foaf/0.1/>  
SELECT ?name WHERE { ?x foaf:name ?name }
```

DUPLICATE RESULTS

Data

```
@prefix foaf: <http://xmlns.com/foaf/0.1/> .  
  
_:x    foaf:name    "Alice" .  
_:x    foaf:mbox    <mailto:alice@example.com> .  
  
_:y    foaf:name    "Alice" .  
_:y    foaf:mbox    <mailto:asmith@example.com> .  
  
_:z    foaf:name    "Alice" .  
_:z    foaf:mbox    <mailto:alice.smith@example.com>  
.
```

Query

```
PREFIX foaf: <http://xmlns.com/foaf/0.1/>  
SELECT ?name WHERE { ?x foaf:name ?name }
```

name
"Alice"
"Alice"
"Alice"

DUPLICATE RESULTS

Data

```
@prefix foaf: <http://xmlns.com/foaf/0.1/> .  
  
_:x    foaf:name    "Alice" .  
_:x    foaf:mbox     <mailto:alice@example.com> .  
  
_:y    foaf:name    "Alice" .  
_:y    foaf:mbox     <mailto:asmith@example.com> .  
  
_:z    foaf:name    "Alice" .  
_:z    foaf:mbox     <mailto:alice.smith@example.com>  
. 
```

```
PREFIX foaf: <http://xmlns.com/foaf/0.1/>  
SELECT DISTINCT ?name WHERE { ?x foaf:name ?name }
```

Query

```
PREFIX foaf: <http://xmlns.com/foaf/0.1/>  
SELECT ?name WHERE { ?x foaf:name ?name }
```

name
"Alice"
"Alice"
"Alice"

DUPLICATE RESULTS

Data

```
@prefix foaf: <http://xmlns.com/foaf/0.1/> .  
  
_:x    foaf:name    "Alice" .  
_:x    foaf:mbox    <mailto:alice@example.com> .  
  
_:y    foaf:name    "Alice" .  
_:y    foaf:mbox    <mailto:asmith@example.com> .  
  
_:z    foaf:name    "Alice" .  
_:z    foaf:mbox    <mailto:alice.smith@example.com>  
.
```

```
PREFIX foaf: <http://xmlns.com/foaf/0.1/>  
SELECT DISTINCT ?name WHERE { ?x foaf:name ?name }
```

Query

```
PREFIX foaf: <http://xmlns.com/foaf/0.1/>  
SELECT ?name WHERE { ?x foaf:name ?name }
```

name
"Alice"
"Alice"
"Alice"

name
"Alice"

TYPES

xsd:integer	
xsd:decimal	
xsd:float	
xsd:double	xsd:nonPositiveInteger
xsd:string	xsd:negativeInteger
xsd:boolean	xsd:long
xsd:dateTime	xsd:int
	xsd:short
	xsd:byte
	xsd:nonNegativeInteger
	xsd:unsignedLong
	xsd:unsignedInt
	xsd:unsignedShort
	xsd:unsignedByte
	xsd:positiveInteger

TYPES & CASTS

```
xsd:integer  
xsd:decimal  
xsd:float  
xsd:double  
xsd:string  
xsd:boolean  
xsd:dateTime  
  
SELECT ?n  
WHERE  
{  
  BIND (xsd:integer("34") AS ?n)  
}
```

```
xsd:nonPositiveInteger  
xsd:negativeInteger  
xsd:long  
xsd:int  
xsd:short  
xsd:byte  
xsd:nonNegativeInteger  
xsd:unsignedLong  
xsd:unsignedInt  
xsd:unsignedShort  
xsd:unsignedByte  
xsd:positiveInteger
```

IF (exp1, exp2, exp3)

?x = 2, ?z = 0,?y unbounded.

IF(?x = 2, "yes", "no") → ???

IF(bound(?y), "yes", "no") → ???

IF(?x=2, "yes", 1/?z) → ???

IF(?x=1, "yes", 1/?z) → ???

IF("2" > 1, "yes", "no") → ???

IF (exp1, exp2, exp3)

?x = 2, ?z = 0,?y unbounded.

IF(?x = 2, "yes", "no")	→ "yes"
IF(bound(?y), "yes", "no")	→ "no"
IF(?x=2, "yes", 1/?z)	→ "yes", and the expression 1/?z is not evaluated
IF(?x=1, "yes", 1/?z)	→ error
IF("2" > 1, "yes", "no")	→ error

IF (exp1, exp2, exp3)

?x = 2, ?z = 0,?y unbounded.

IF(?x = 2, "yes", "no")	→ "yes"
IF(bound(?y), "yes", "no")	→ "no"
IF(?x=2, "yes", 1/?z)	→ "yes", and the expression 1/?z is not evaluated
IF(?x=1, "yes", 1/?z)	→ error
IF("2" > 1, "yes", "no")	→ error

```
SELECT ?result
WHERE
{
    BIND ( IF (3>4, « True», « False») AS ?result)
}
```

Data

```
@prefix foaf: <http://xmlns.com/foaf/0.1/> .  
  
_:a foaf:name "Alice".  
_:a foaf:mbox <mailto:alice@work.example> .  
  
_:b foaf:name "Ms A.".  
_:b foaf:mbox <mailto:alice@work.example> .
```

name1	name2
"Alice"	"Ms A."
"Ms A."	"Alice"

Query 1

```
PREFIX foaf: <http://xmlns.com/foaf/0.1/>  
SELECT ?name1 ?name2  
WHERE { ?x foaf:name ?name1 ;  
        foaf:mbox ?mbox1 .  
        ?y foaf:name ?name2 ;  
        foaf:mbox ?mbox2 .  
        FILTER (sameTerm(?mbox1, ?mbox2) && !sameTerm(?name1, ?name2))  
    }
```

Query 2

```
PREFIX foaf: <http://xmlns.com/foaf/0.1/>  
SELECT ?name1 ?name2  
WHERE { ?x foaf:name ?name1 ;  
        foaf:mbox ?mbox1 .  
        ?y foaf:name ?name2 ;  
        foaf:mbox ?mbox2 .  
        FILTER (?mbox1 = ?mbox2 && ?name1 != ?name2)  
    }
```

Different writing of the same IRI

<http://example.org/book/book1>

BASE <http://example.org/book/>

<book1>

PREFIX book: <http://example.org/book/>

book:book1

IRI Syntax

Different writing of the same query

```
PREFIX dc: <http://purl.org/dc/elements/1.1/>
SELECT ?title
WHERE { <http://example.org/book/book1> dc:title ?title }
```

```
PREFIX dc: <http://purl.org/dc/elements/1.1/>
PREFIX : <http://example.org/book/>
SELECT $title
WHERE { :book1 dc:title $title }
```

```
BASE <http://example.org/book/>
PREFIX dc: <http://purl.org/dc/elements/1.1/>
SELECT $title
WHERE { <book1> dc:title ?title }
```

Triple Syntax

Triple Syntax

```
?x foaf:name ?name ;  
    foaf:mbox ?mbox .  
=   
?x foaf:name ?name .  
?x foaf:mbox ?mbox .
```

```
?x foaf:nick "Alice" , "Alice_" .  
=   
?x foaf:nick "Alice" .  
?x foaf:nick "Alice_" .
```

```
?x foaf:name ?name ; foaf:nick "Alice" , "Alice_" .  
=   
?x foaf:name ?name .  
?x foaf:nick "Alice" .  
?x foaf:nick "Alice_" .
```

```
PREFIX foaf: <http://xmlns.com/foaf/0.1/>
SELECT ?name ?mbox
WHERE {
    ?x foaf:name ?name .
    ?x foaf:mbox ?mbox .
}
```

```
PREFIX foaf: <http://xmlns.com/foaf/0.1/>
SELECT ?name ?mbox
WHERE {
    { ?x foaf:name ?name . }
    { ?x foaf:mbox ?mbox . }
}
```

Groups of patterns

```
{ ?x foaf:name ?name .  
?x foaf:mbox ?mbox .  
FILTER regex(?name, "Smith")  
}
```

FILTER

```
{ FILTER regex(?name, "Smith")  
?x foaf:name ?name .  
?x foaf:mbox ?mbox .  
}
```

```
{ ?x foaf:name ?name .  
FILTER regex(?name, "Smith")  
?x foaf:mbox ?mbox .  
}
```

Optional Pattern

Data

```
@prefix foaf: <http://xmlns.com/foaf/0.1/> .  
@prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#> .  
  
_:a rdf:type foaf:Person .  
_:a foaf:name "Alice" .  
_:a foaf:mbox <mailto:alice@example.com> .  
_:a foaf:mbox <mailto:alice@work.example> .  
_:b rdf:type foaf:Person .  
_:b foaf:name "Bob" .
```

Query

```
PREFIX foaf: <http://xmlns.com/foaf/0.1/>  
SELECT ?name ?mbox  
WHERE { ?x foaf:name ?name .  
       OPTIONAL { ?x foaf:mbox ?mbox }  
 }
```

name	mbox
"Alice"	<mailto:alice@example.com>
"Alice"	<mailto:alice@work.example>
"Bob"	

```

@prefix dc: <http://purl.org/dc/elements/1.1/> .
@prefix : <http://example.org/book/> .
@prefix ns: <http://example.org/ns#> .

:book1 dc:title "SPARQL Tutorial" .
:book1 ns:price 42 .
:book2 dc:title "The Semantic Web" .
:book2 ns:price 23 .

```

title	price
"SPARQL Tutorial"	
"The Semantic Web"	23

```

PREFIX dc: <http://purl.org/dc/elements/1.1/>
PREFIX ns: <http://example.org/ns#>
SELECT ?title ?price
WHERE { ?x dc:title ?title .
        OPTIONAL { ?x ns:price ?price . FILTER (?price < 30) }
}

```

Test the absence of a pattern

Data

```
@prefix : <http://example/> .  
@prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#> .  
@prefix foaf: <http://xmlns.com/foaf/0.1/> .  
  
:alice rdf:type foaf:Person .  
:alice foaf:name "Alice" .  
:bob rdf:type foaf:Person .
```

Query

```
PREFIX rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#>  
PREFIX foaf: <http://xmlns.com/foaf/0.1/>  
SELECT ?person  
WHERE {  
  ?person rdf:type foaf:Person .  
  FILTER NOT EXISTS { ?person foaf:name ?name }  
}
```

person
[<http://example/bob>](http://example/bob)

Test the presence of a pattern

Data

```
@prefix : <http://example/> .  
@prefix rdf: <http://www.w3.org/1999/02/22-rdf-syntax-  
ns#> .  
@prefix foaf: <http://xmlns.com/foaf/0.1/> .  
  
:alice rdf:type foaf:Person .  
:alice foaf:name "Alice" .  
:bob rdf:type foaf:Person .
```

Query

```
PREFIX rdf: <http://www.w3.org/1999/02/22-rdf-syntax-ns#>  
PREFIX foaf: <http://xmlns.com/foaf/0.1/>
```

```
SELECT ?person  
WHERE  
{  
    ?person rdf:type foaf:Person .  
    FILTER EXISTS { ?person foaf:name ?name }  
}
```

person
<http://example/alice>

Minus

Data

```
@prefix : <http://example/> .  
@prefix foaf: <http://xmlns.com/foaf/0.1/> .  
  
:alice foaf:givenName "Alice" ;  
       foaf:familyName "Smith" .  
  
:bob   foaf:givenName "Bob" ;  
       foaf:familyName "Jones" .  
  
:carol  foaf:givenName "Carol" ;  
       foaf:familyName "Smith" .
```

Query

```
PREFIX : <http://example/>  
PREFIX foaf: <http://xmlns.com/foaf/0.1/>  
  
SELECT DISTINCT ?s  
WHERE {  
  ?s ?p ?o .  
  MINUS {  
    ?s foaf:givenName "Bob" .  
  }  
}
```

s
`<http://example/carol>`
`<http://example/alice>`

LISTS IN & NOT IN

2 IN (1, 2, 3)	???
2 IN ()	???
2 IN (< http://example/iri >, "str", 2.0)	???
2 IN (1/0, 2)	???
2 IN (2, 1/0)	???
2 IN (3, 1/0)	???

2 NOT IN (1, 2, 3)	???
2 NOT IN ()	???
2 NOT IN (< http://example/iri >, "str", 2.0)	???
2 NOT IN (1/0, 2)	???
2 NOT IN (2, 1/0)	???
2 NOT IN (3, 1/0)	???

LISTS IN & NOT IN

2 IN (1, 2, 3)	true
2 IN ()	false
2 IN (<http://example/iri>, "str", 2.0)	true
2 IN (1/0, 2)	true
2 IN (2, 1/0)	true
2 IN (3, 1/0)	error

2 NOT IN (1, 2, 3)	false
2 NOT IN ()	true
2 NOT IN (<http://example/iri>, "str", 2.0)	false
2 NOT IN (1/0, 2)	false
2 NOT IN (2, 1/0)	false
2 NOT IN (3, 1/0)	error

STRINGS

■ STRSTARTS, STRENDNS, CONTAINS, STRBEFORE & STRAFTER

Argument1	Argument2	Compatible?		
"abc"	"b"	Oui	strStarts("foobar", "foo")	true
"abc"	"b"^^xsd:string	Oui	strStarts("foobar"@en, "foo"@en)	True
"abc"^^xsd:string	"b"	Oui	strEnds("foobar"^^xsd:string, "bar"^^xsd:string)	true
"abc"^^xsd:string	"b"^^xsd:string	Oui	strEnds("foobar"^^xsd:string, "bar")	True
"abc"@en	"b"	Oui	contains("foobar", "bar"^^xsd:string)	true
"abc"@en	"b"@en	Oui	contains("foobar"@en, "foo")	true
"abc"@en	"b"^^xsd:string	Oui	strbefore("abc", "b")	"a"
"abc"@en	"b"@en	Oui	strbefore("abc"@en, "bc")	"a"@en
"abc"@fr	"b"@ja	Non	strbefore("abc"@en, "b"@cy)	error
"abc"	"b"@ja	Non	strbefore("abc"^^xsd:string, "")	""^^xsd:string
"abc"	"b"@en	Non	strbefore("abc", "xyz")	" "
"abc"^^xsd:string	"b"@en	Non	strafter("abc"@en, "z")	" "
			strafter("abc"@en, ""@en)	"abc"@en
			strafter("abc"@en, "")	"abc"@en

STRINGS

► CONCAT

concat("foo", "bar")	"foobar"
concat("foo"@en, "bar"@en)	"foobar"@en
concat("foo"^^xsd:string, "bar"^^xsd:string)	"foobar"^^xsd:string
concat("foo", "bar"^^xsd:string)	"foobar"
concat("foo"@en, "bar")	"foobar"
concat("foo"@en, "bar"^^xsd:string)	"foobar"

strlen("chat")	4
strlen("chat"@en)	4
strlen("chat"^^xsd:string)	4

► SUBSTR

substr("foobar", 4)	"bar"
substr("foobar"@en, 4)	"bar"@en
substr("foobar"^^xsd:string, 4)	"bar"^^xsd:string
substr("foobar", 4, 1)	"b"
substr("foobar"@en, 4, 1)	"b"@en
substr("foobar"^^xsd:string, 4, 1)	"b"^^xsd:string

Categories	Functions / Operators	Examples
Logic & Comparisons	!, &&, , =, !=, <, <=, >, >=, IN, NOT IN	?hasPermit ?age < 25
Conditions (SPARQL 1.1)	EXISTS, NOT EXISTS, IF, COALESCE	NOT EXISTS { ?p foaf:mbox ?email }
Math	+, -, *, /, abs, round, ceil, floor, RAND	?decimal * 10 > ?minPercent
Strings (SPARQL 1.1)	STRLEN, SUBSTR, UCASE, LCASE, STRSTARTS, CONCAT, STREND, CONTAINS, STRBEFORE, STRAFTER	STRLEN(?description) < 255
Date/time (SPARQL 1.1)	now, year, month, day, hours, minutes, seconds, timezone, tz	month(now()) < 4
Tests SPARQL	isURI, isBlank, isLiteral, isNumeric, bound	isURI(?person) !bound(?person)
Constructors (SPARQL 1.1)	URI, BNODE, STRDT, STRLANG, UUID, STRUUID	STRLANG(?text, "en") = "hello"@en
Accessors	str, lang, datatype	lang(?title) = "en"
Hashing (1.1)	MD5, SHA1, SHA256, SHA512	BIND(SHA256(?email) AS ?hash)
Miscellaneous	sameTerm, langMatches, regex, REPLACE	regex(?ssn, "\d{3}-\d{2}-\d{4}")

Construct



CONSTRUCT

```
CONSTRUCT {  
    #the form of triples to be constructed  
}  
WHERE {  
    #the selected constituents  
}
```

```
CONSTRUCT {  
    ?x a foaf:Person .  
    ?y a foaf:Person .  
}  
WHERE {  
    ?x foaf:knows ?y .  
}
```

CONSTRUCT with optional data

```
CONSTRUCT {  
    ?song dcterms:title ?title.  
    ?song music:hasComposer ?composer.  
    ?composer foaf:name ?cname.  
    ?composer music:birthYear ?y.  
}  
WHERE {  
    ?song dcterms:title ?title.  
    OPTIONAL {  
        ?song music:hasComposer ?composer.  
        ?composer foaf:name ?cname.  
        ?composer music:birthYear ?y.  
    }  
}
```

To learn more ...



SPARQL 1.1 Query Language

W3C Recommendation 21 March 2013

<http://www.w3.org/TR/sparql11-query/>



4

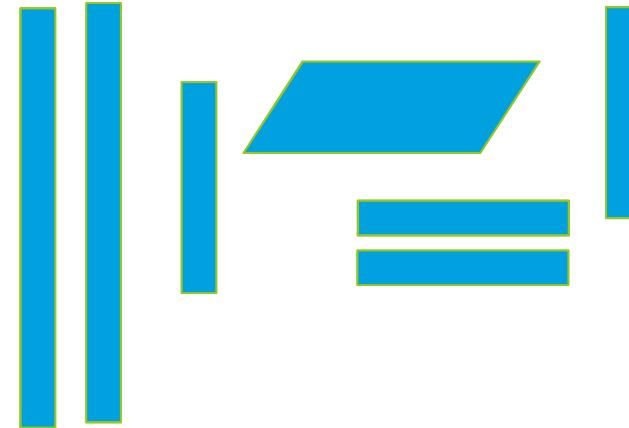
Understanding the basics of knowledge modeling (ontology, OWL)

1. The basic principles of the Web of Data: from raw and heterogeneous Data to the lingua franca of the Web of Data
2. The RDF semantic and its main syntaxes
3. Handling and querying Data (triple store, end point, SPARQL)
4. Understanding the basics of knowledge modeling (ontology, OWL)
5. The lifting and linking of Data in order to publish 5 stars Data (DataLift).

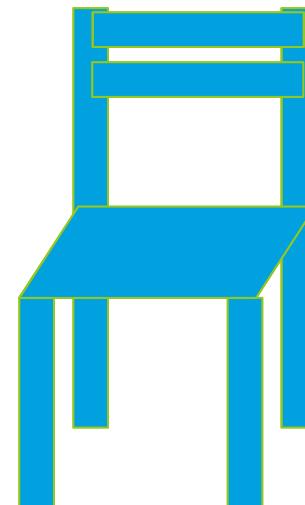
What can we do with an ontology?

- Understand
- Capture
- Model
- Transform
- Integrate contextualized knowledge
- Do
- Share
- Reuse

- Ontologies provide a common vocabulary
 - The meaning of the classes and properties is well defined, it allows us to "understand" the data that we exchange.
- Ontologies contain general knowledge, valid for all instances of data
 - they allow us to deduce new facts from the available data



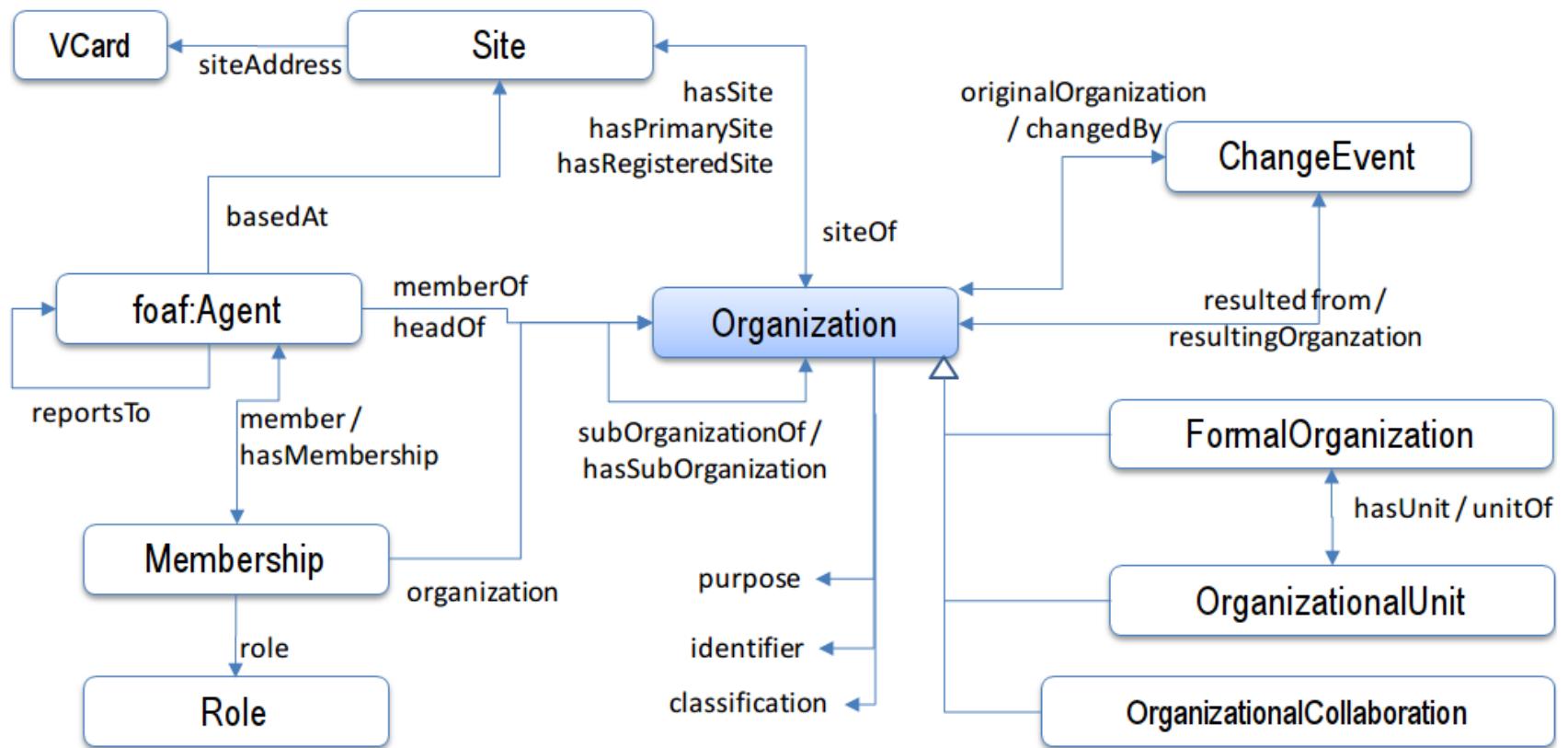
When is it a chair?

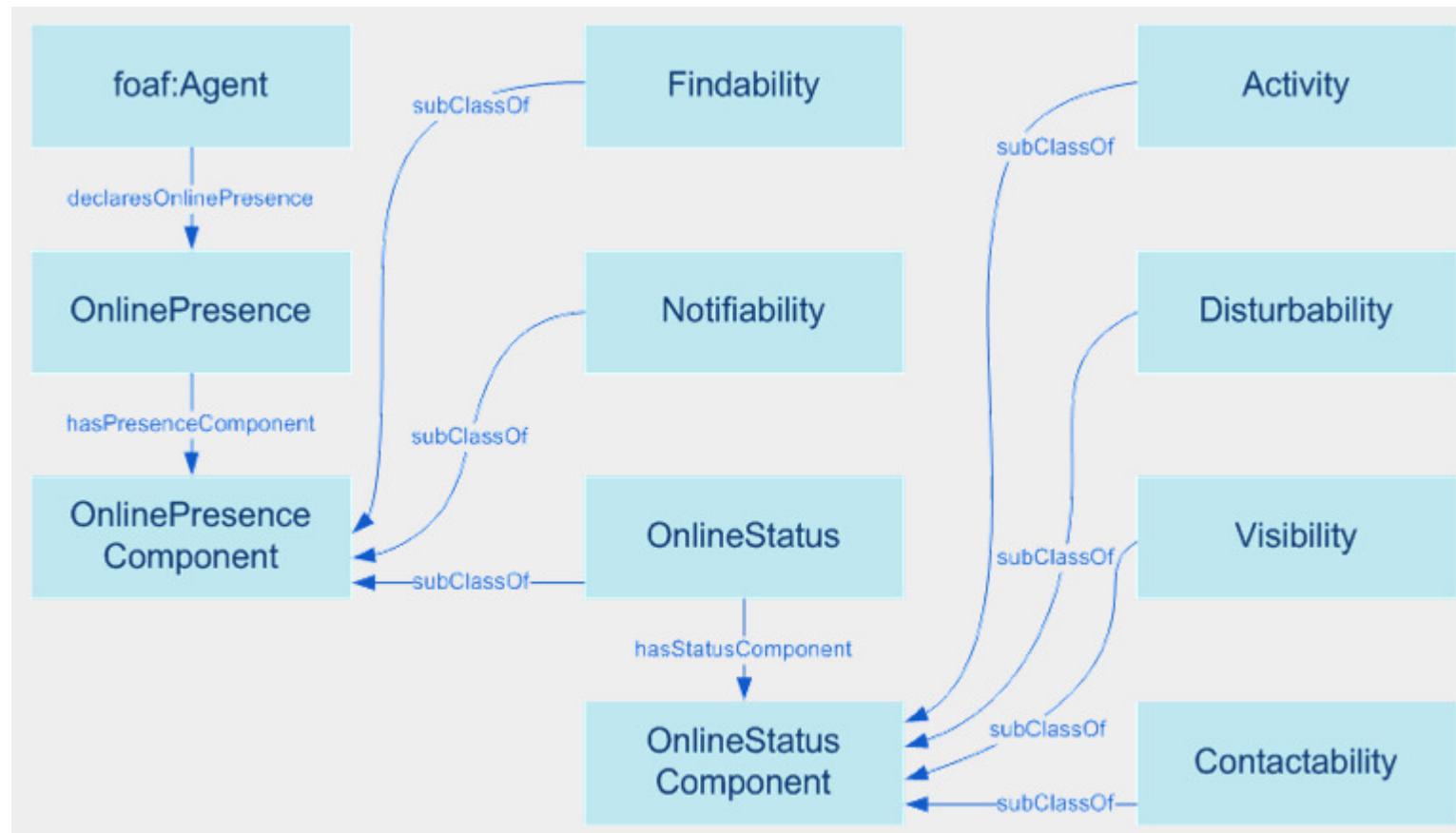


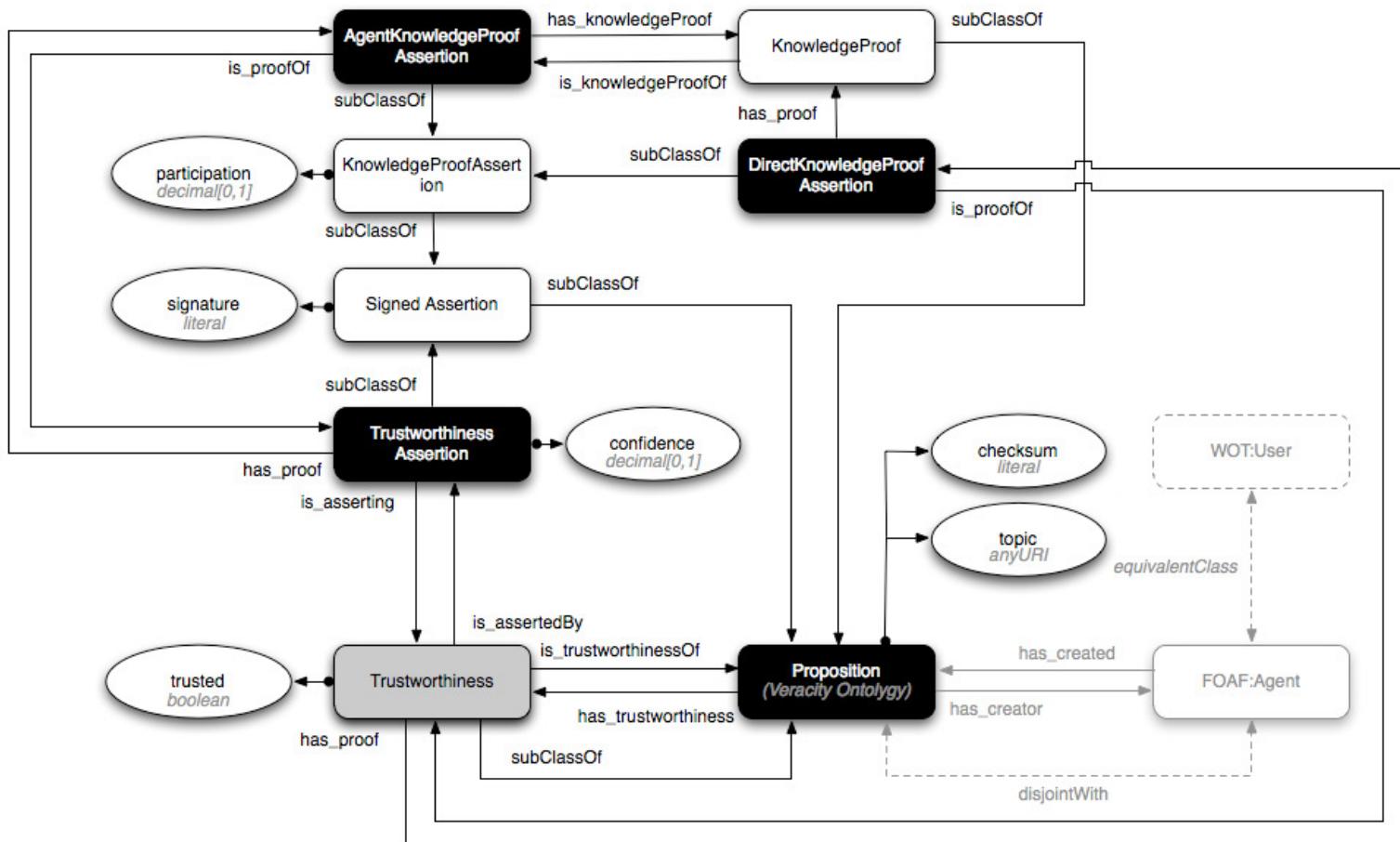
Ontology

Ontology samples







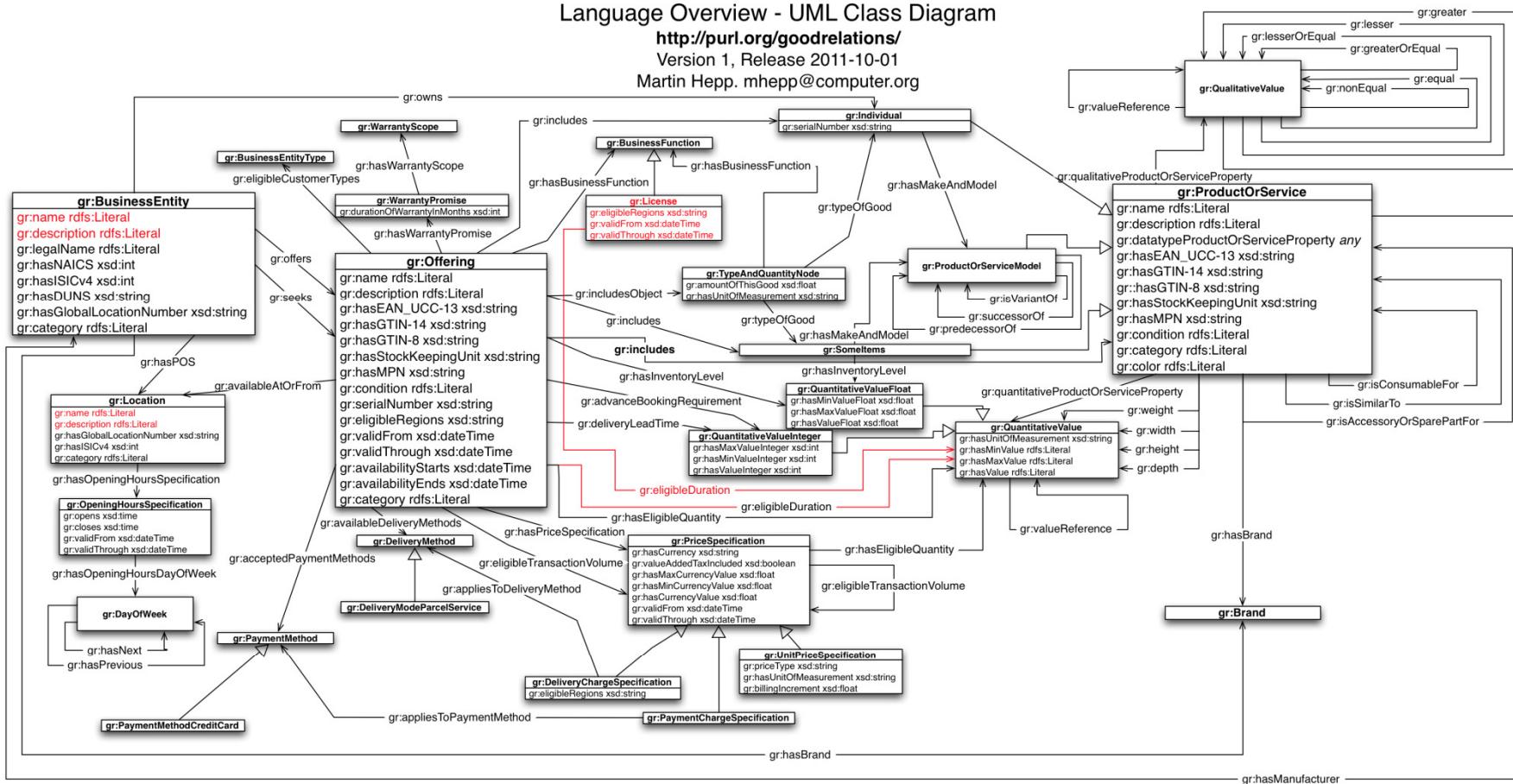


The GoodRelations Ontology for E-Commerce Language Overview - UML Class Diagram

<http://purl.org/goodrelations/>

Version 1, Release 2011-10-01

Martin Hepp. mhepp@computer.org



Notes:

1. The following GoodRelations elements are only shortcuts for simpler annotation or querying. See the documentation at <http://purl.org/goodrelations/> for details:
`gr:hasValue` (shortcut for setting both `hasMinValue` and `hasMaxValue` properties to the same value in one turn)
`gr:hasValueFloat` (shortcut for setting both `hasMinValueFloat` and `hasMaxValueFloat` properties to the same value in one turn)
`gr:hasValueInteger` (shortcut for setting both `hasMinValueInteger` and `hasMaxValueInteger` properties to the same value in one turn)
2. The following elements are now deprecated, but you can still use them, e.g. for staying compatible with older data consumers (e.g. Yahoo SearchMonkey):
`gr:ActualProductOrServiceInstance` (now `gr:Individual`)
`gr:ProductOrServicesSomeInstancesPlaceholder` (now `gr:SomeItems`)
`gr:LocationOfSalesOrServiceProvisioning` (now `gr:Location`)
3. For the recommended cardinality of attributes, see the GoodRelations Language Reference at <http://purl.org/goodrelations/v1.html>.
4. `gr:valueReference` links can also exist between a `gr:QualitativeValue` and a `gr:QuantitativeValue` and vice versa, but this rare case is not shown for readability.
5. `gr:name` and `gr:description` can now be attached to any GoodRelations type, but this is not shown here for readability.

Red highlighting indicates elements added or changed in this release.



OWL

Fragments from the
Wine Ontology
example in the OWL
Guide...

Namespaces

```
<rdf:RDF xmlns =http://www.example.org/wine#
           xmlns:vin =http://www.example.org/wine#
           xmlns:food=http://www.example.org/food#
           xmlns:owl ="http://www.w3.org/2002/07/owl#"
           xmlns:rdf =http://www.w3.org/1999/02/22-rdf-syntax-ns#
           xmlns:rdfs=http://www.w3.org/2000/01/rdf-schema#
           xmlns:xsd ="http://www.w3.org/2000/10/XMLSchema#">
```

Ontology header

```
<owl:Ontology rdf:about="http://www.example.org/wine">
  <rdfs:comment>An example OWL ontology</rdfs:comment>
  <owl:priorVersion
    rdf:resource="http://www.example.org/wine-2102.owl"/>
  <owl:imports rdf:resource="http://www.example.org/food.owl"/>
  <rdfs:label>Wine Ontology</rdfs:label>
  ...
</owl:Ontology>
```

Simple classes

```
<owl:Class rdf:ID="Winery"/>
```

```
<owl:Class rdf:ID="Region"/>
```

```
<owl:Class rdf:ID="ConsumableThing"/>
```

```
<owl:Class rdf:ID="Wine">
```

```
  <rdfs:subClassOf rdf:resource="&food;PotableLiquid"/>
```

```
  <rdfs:label xml:lang="en">wine</rdfs:label>
```

```
  <rdfs:label xml:lang="fr">vin</rdfs:label>
```

```
  ...
```

```
</owl:Class>
```

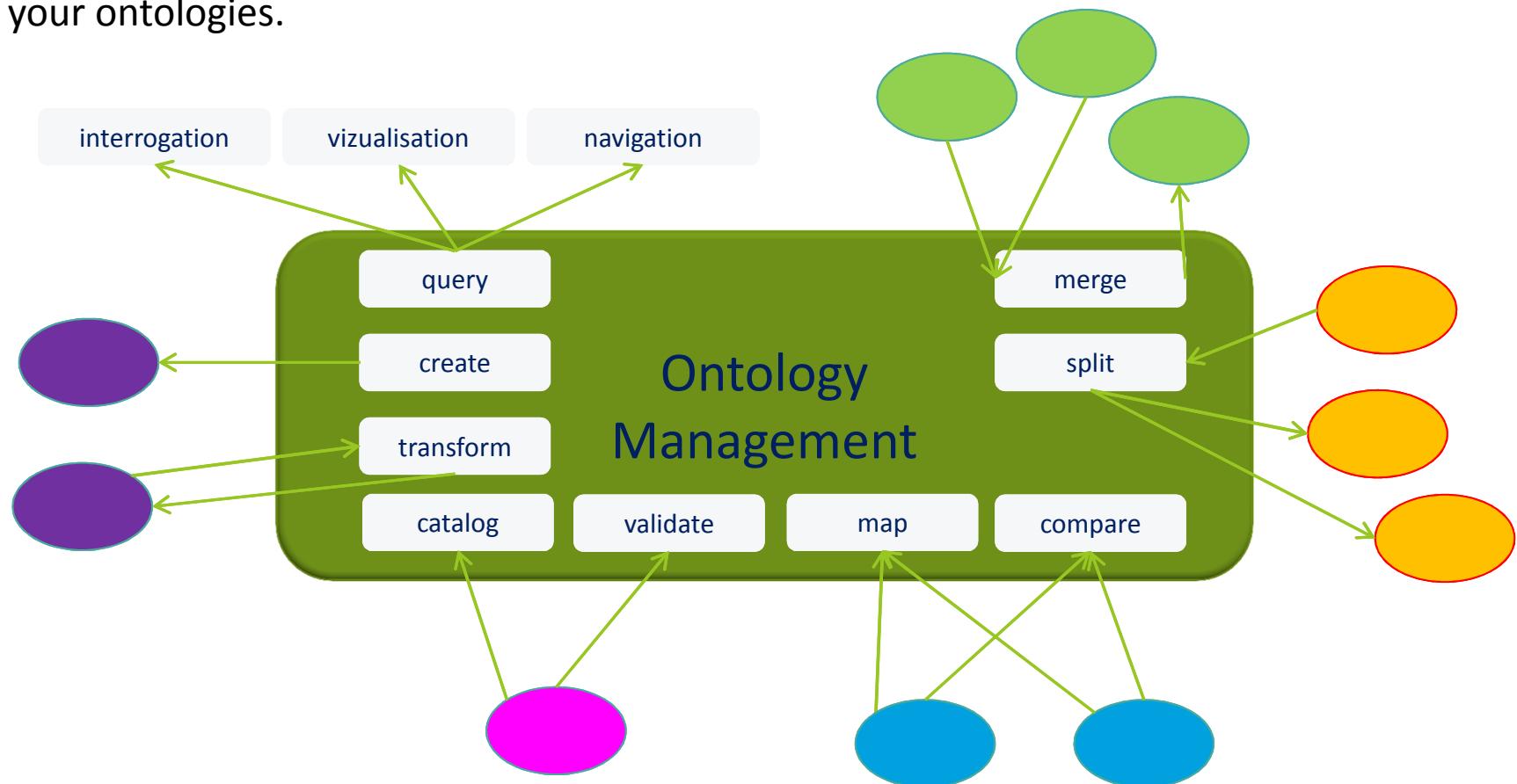
Some tools

A small panorama



Ontology Management

- Supporting the ontology data model are a number of processing technologies to manage your ontologies.



Ontology creation

properties

A person **knows** other persons

logic

« is a brother of » is a symetrical property

classes (a value of a property)

A person is a **living been**

Ontology creation is not a uniforme question...

- Tow approches
 - Bottom – Up
 - Top – Down
- Different worlds
 - Unstructured data
 - Semi structured data
 - Structured data
- Knowledges at different levels
 - Task, sensor, ...
 - Application, system, ...
 - Domain, ecosystem, ...
 - World, univers, ...



Ontology Creation

- **From texts**

- **Terminae (2002)**
 - Manual Conceptualisation
- **Text2Onto (2005)**
 - Automatic Conceptualisation
- **OntoGen (2006)**
 - semi-automatic Conceptualisation
- **Etc.**

- **From data**

- Generalisation from data analyse

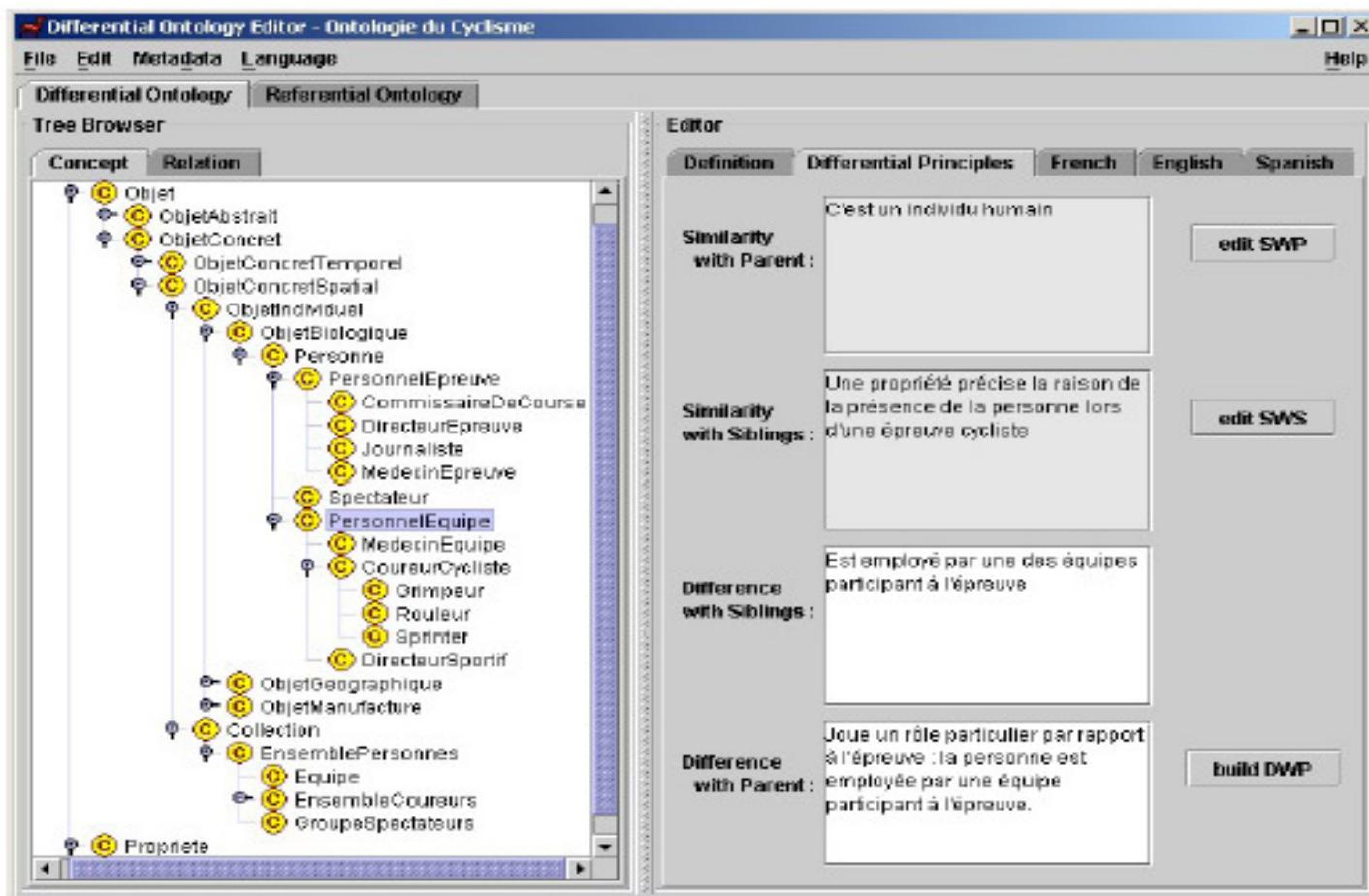
- **From structures**

- Transformation

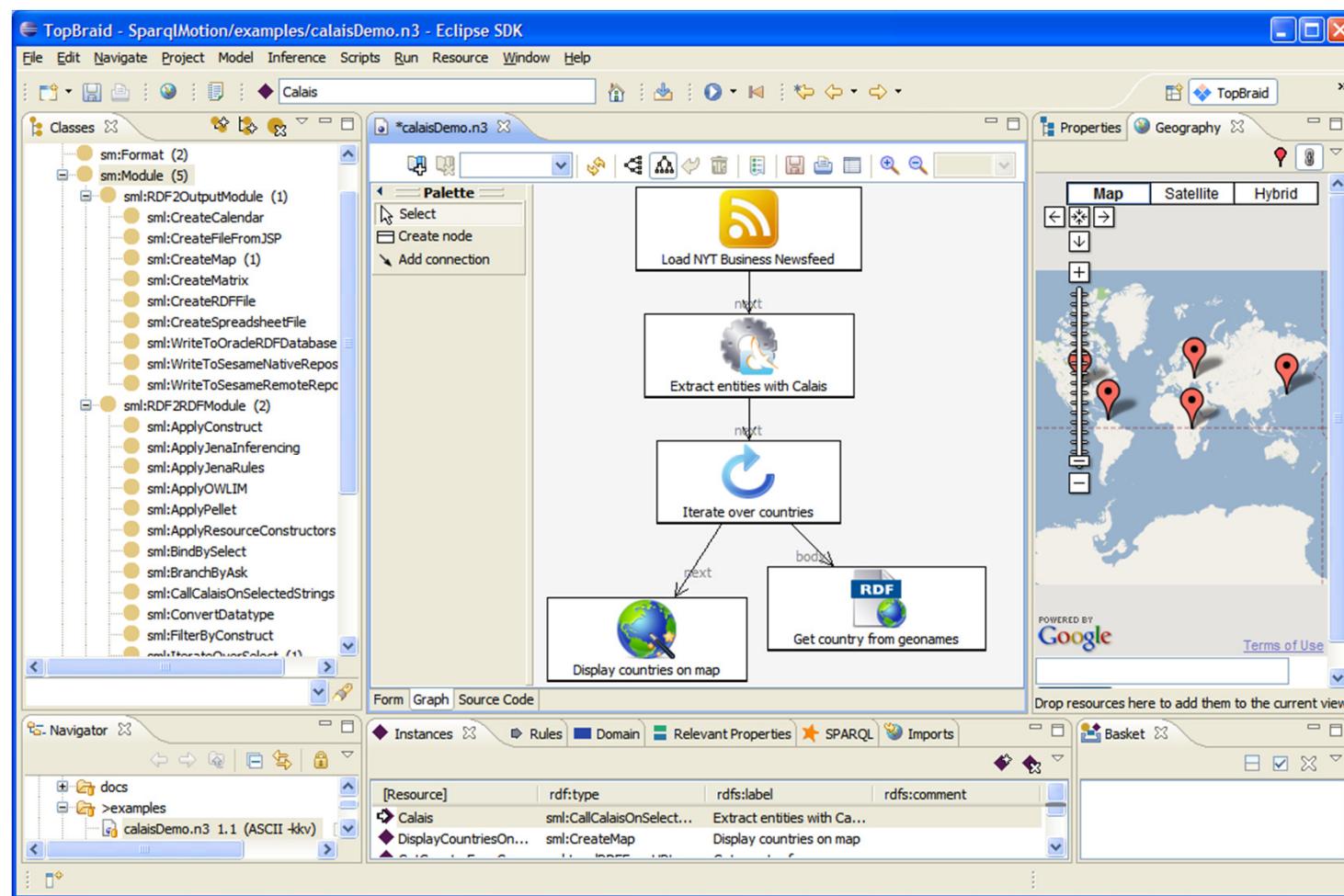
DOE : Differential Ontology Editor

- **DOE** is a simple ontology editor which allows the user to build ontologies according to the methodology proposed by *Bruno Bachimont*.
- The specification process is divided in 3 steps.
 1. The user is invited to build taxonomies of concepts and relations, explicitly justifying the position of each item (notion) in the hierarchy. For each notion, the user builds a definition following 4 principles which come from the *Differential Semantics* theory. Hence, the user has to explicit why a notion is similar but more specific than its parent (2 principles), and why this notion is similar but different from its siblings (2 others principles). The user can also add synonyms and encyclopedic definition in a few languages for all notions.
 2. The two taxonomies are considered from an extensional semantics point of view. The user can augment them with new entities (defined) or add constraints onto the domains of the relations.
 3. The ontology can be translated into a knowledge representation language, which allows to use it in an appropriate ontology-based system or to import it into another ontology-building tool to specify it further:

DOE : Differential Ontology Editor

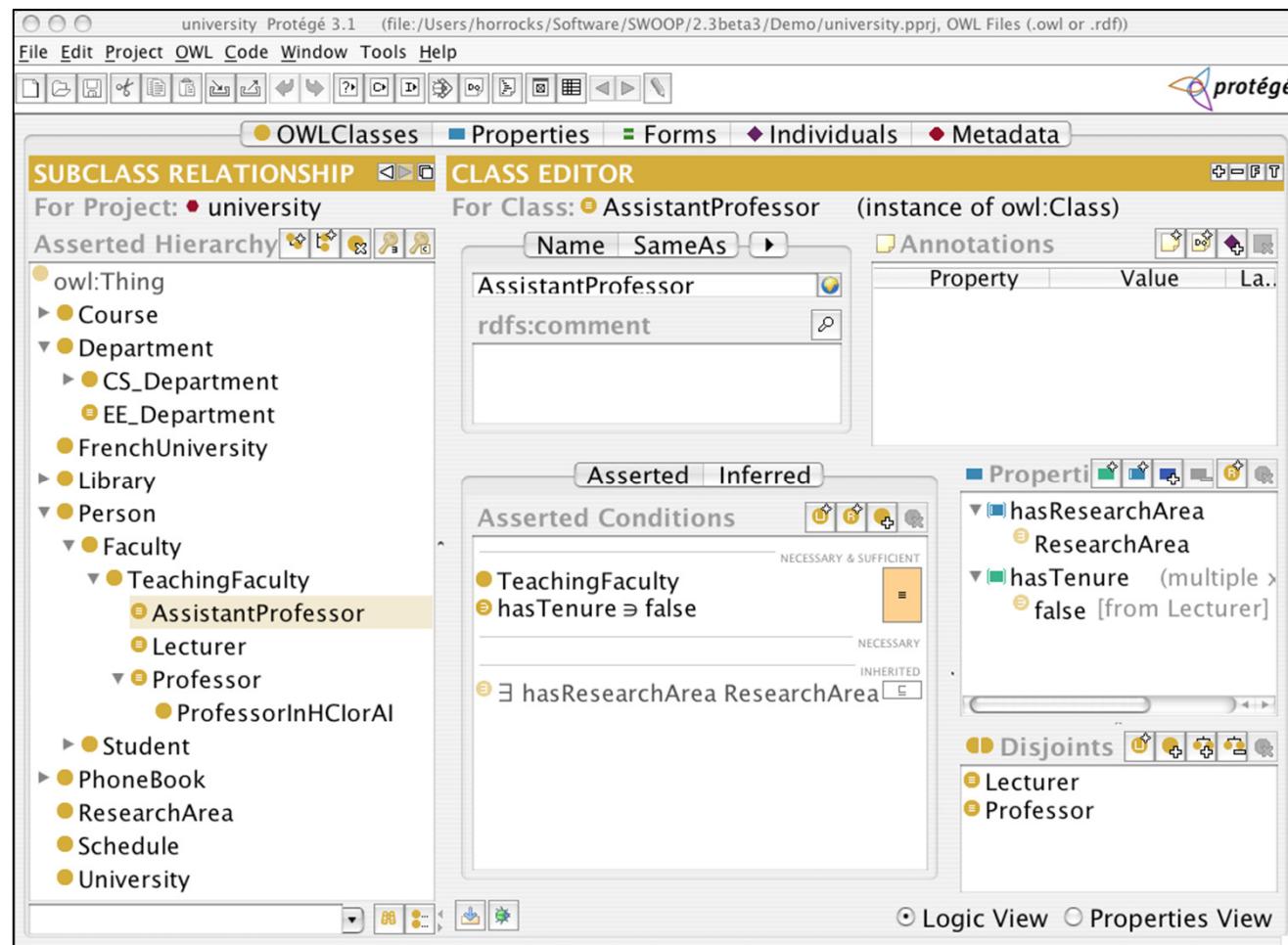


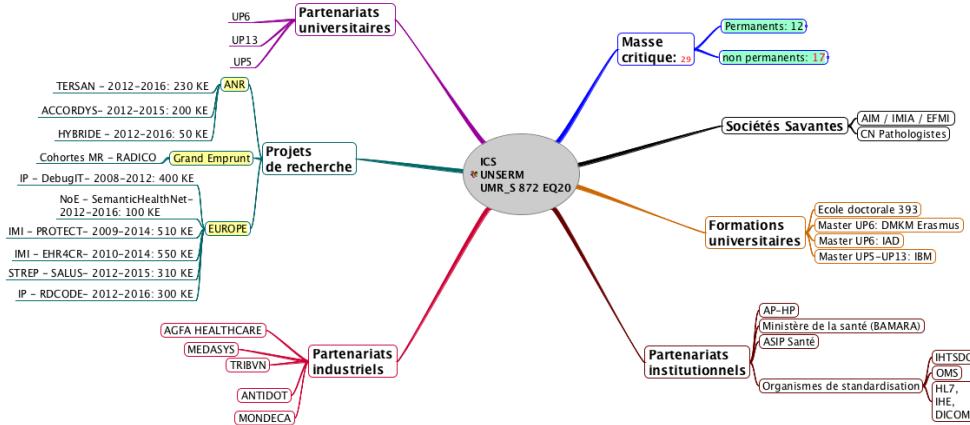
TopBraid composer



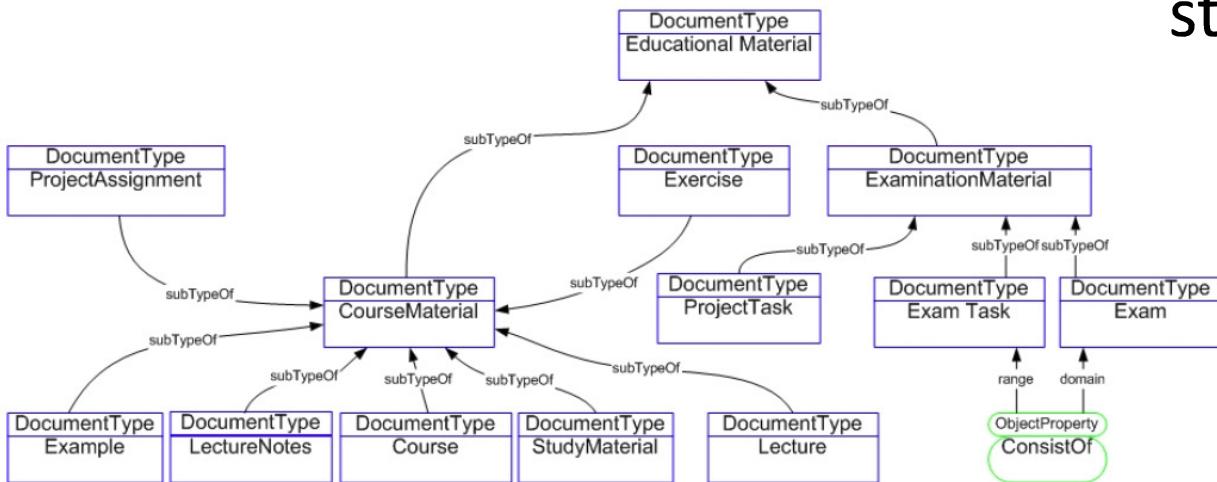
Protégé

An ontology editor





It helps to
structure thinking



Oyster - Distributed Ontology Registry

Oyster is a distributed registry that exploits semantic web techniques in order to provide a solution for the management of ontology related metadata in distributed environments.



[ontology](#) [document](#) [term](#) [across ontologies](#)

[Swoogle Search](#) [?](#)

Searching over 10,000 ontologies

[manual](#) o [news](#) o [faq](#) o [web-service](#) o [submit-url](#) o [sw-archive](#) o [feedback](#) o [swoogle2005](#)

Swoogle © 2004-2007, [ebiquity group](#) at UMBC
This work is licensed under a [Creative Commons Attribution-NonCommercial-ShareAlike 2.5 License](#).



Catalog : LOV

Vocabulary spaces

Classifying vocabularies is a brand new challenge for Library Science, and no existing classification scheme seems to fit the need so far. Unless we missed something, and in that case we are open to suggestions from the librarian community.



This feature is highly experimental and its structure is unstable, likely to be often challenged by the growth of LOV, but gives a high-level view of the ecosystem, complementary to other ways of exploration such as [SPARQL endpoints](#), [full-text search](#), and [navigation through links](#). Vocabulary Spaces are also used as one of the facets in the [full-text search](#).



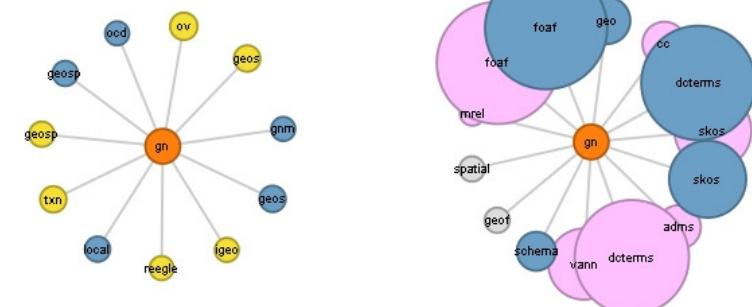


Metadata:

Property	Value
is part of vocabulary space	All > Where & When > Geography
Vocabulary URI	http://www.geonames.org/ontology
Prefix	gn
Namespace URI	http://www.geonames.org/ontology#
Description	The Geonames ontology provides elements of description for geographical features defined in the geonames.org data base @en
Date issued	2006-10-14
Last modified	2012-10-29
Language	Russian, Swedish, English, Bulgarian, Norwegian
Creator	Bernard Vatant
Publisher	Geonames.org
Class number	7
Property number	26
Homepage	http://www.geonames.org/ontology/documentation.html
See also	http://stats.lod2.eu/vocabularies/95
Represented by	format-geonames

Vocabulary links:

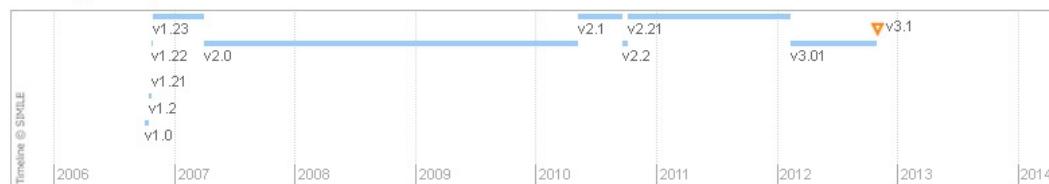
Vocabularies referencing "gn" (11)



Vocabularies referenced by "gn" (14)



Vocabulary history:



The LOV dataset is licensed under Creative Commons CC BY 3.0. It is developed in the framework of the Datalift project and supported by the Open Knowledge Foundation (OKFN).

Catalog : LOV

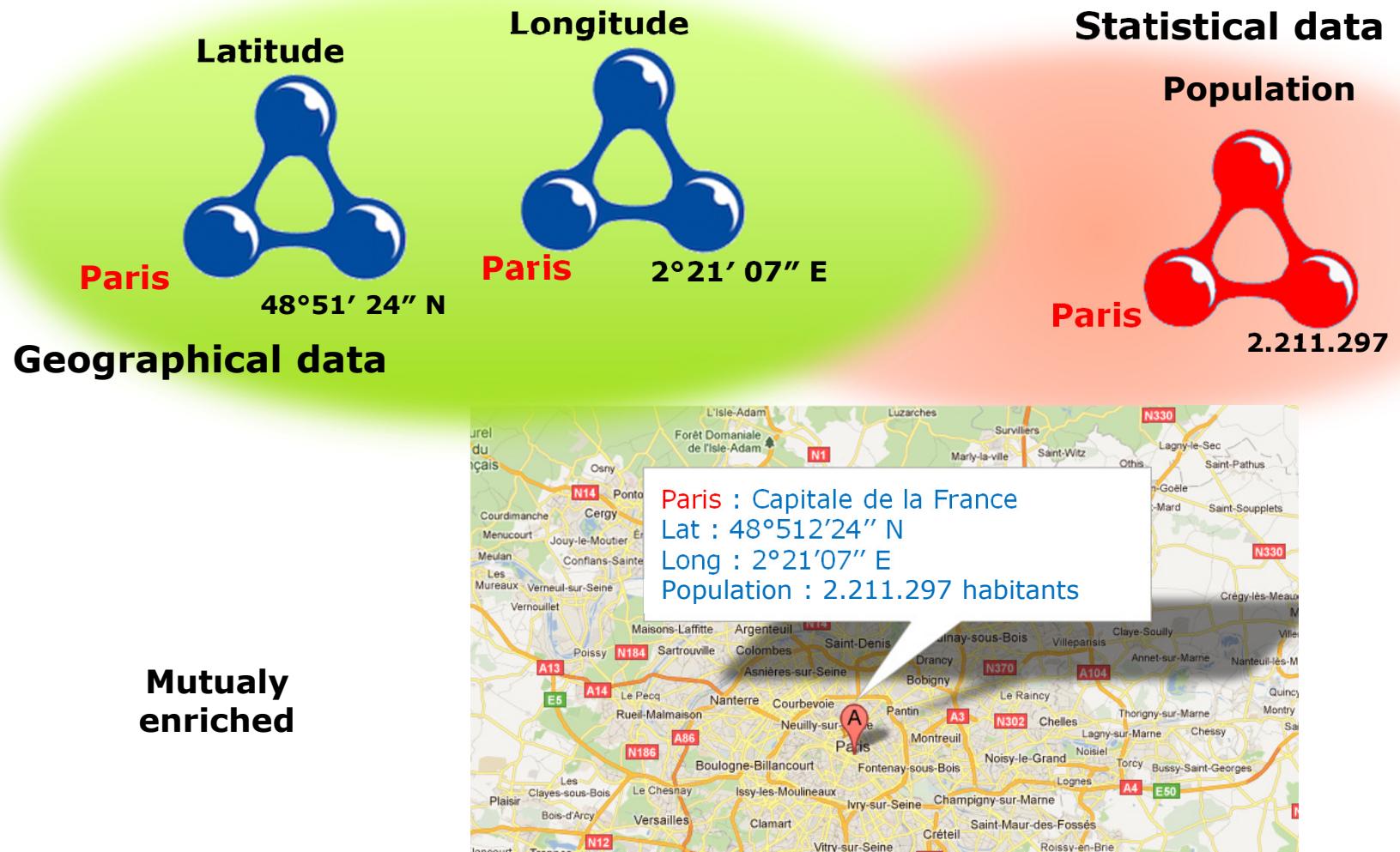


Languages to express ontologies

- Two major languages (W3C)
 - **RDFS**
 - extensible language for knowledge representation
 - provides basic elements for the definition of ontologies or vocabularies intended to structure RDF resources
 - V1 in March 1999
 - Final recommendation in February 2004
 - the main components of RDFS are included in OWL , a more expressive language
 - **OWL**
 - knowledge representation built on the RDF model
 - based on the description logic
 - designed as an extension of RDFS
 - V1 in feb. 2004
 - V2 in oct. 2009

Ontology alignment

What is interconnection?



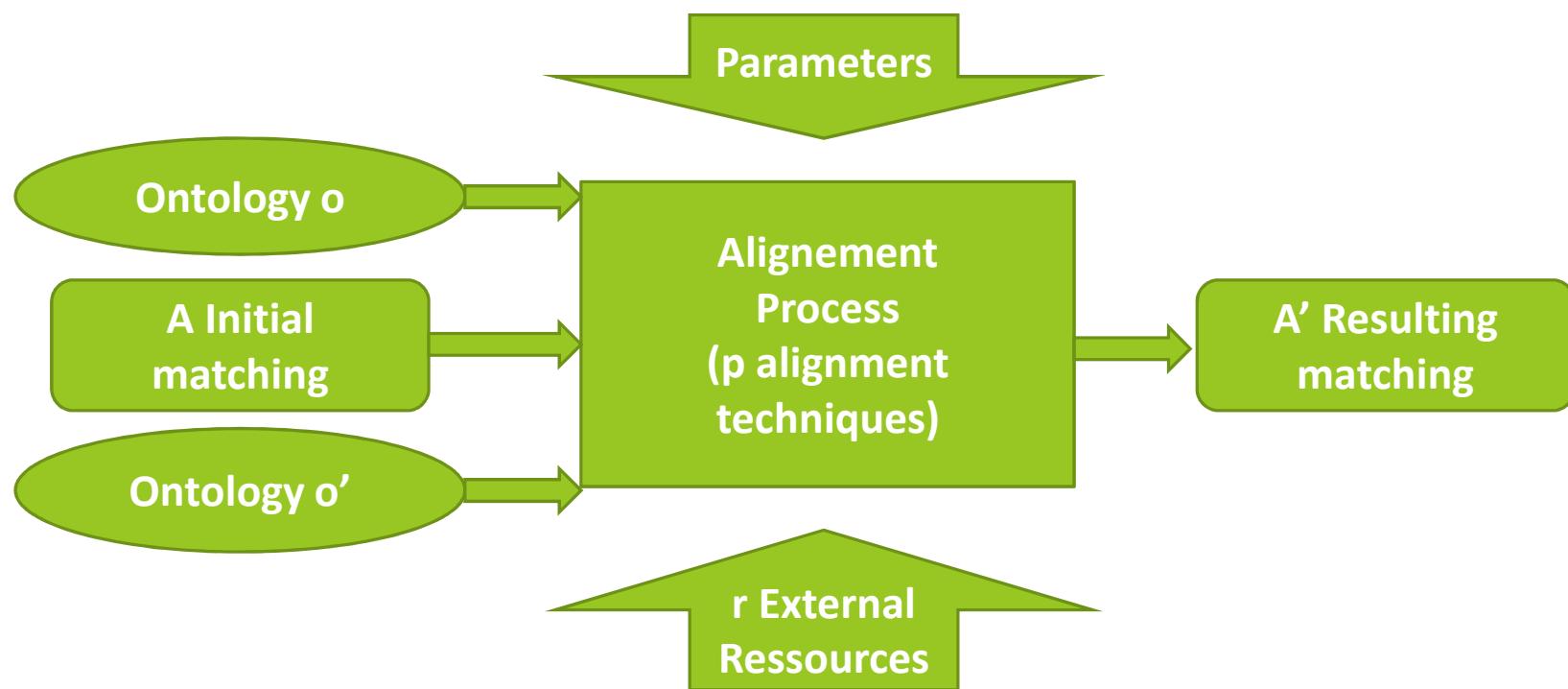
Ontology matching is a function f applied:

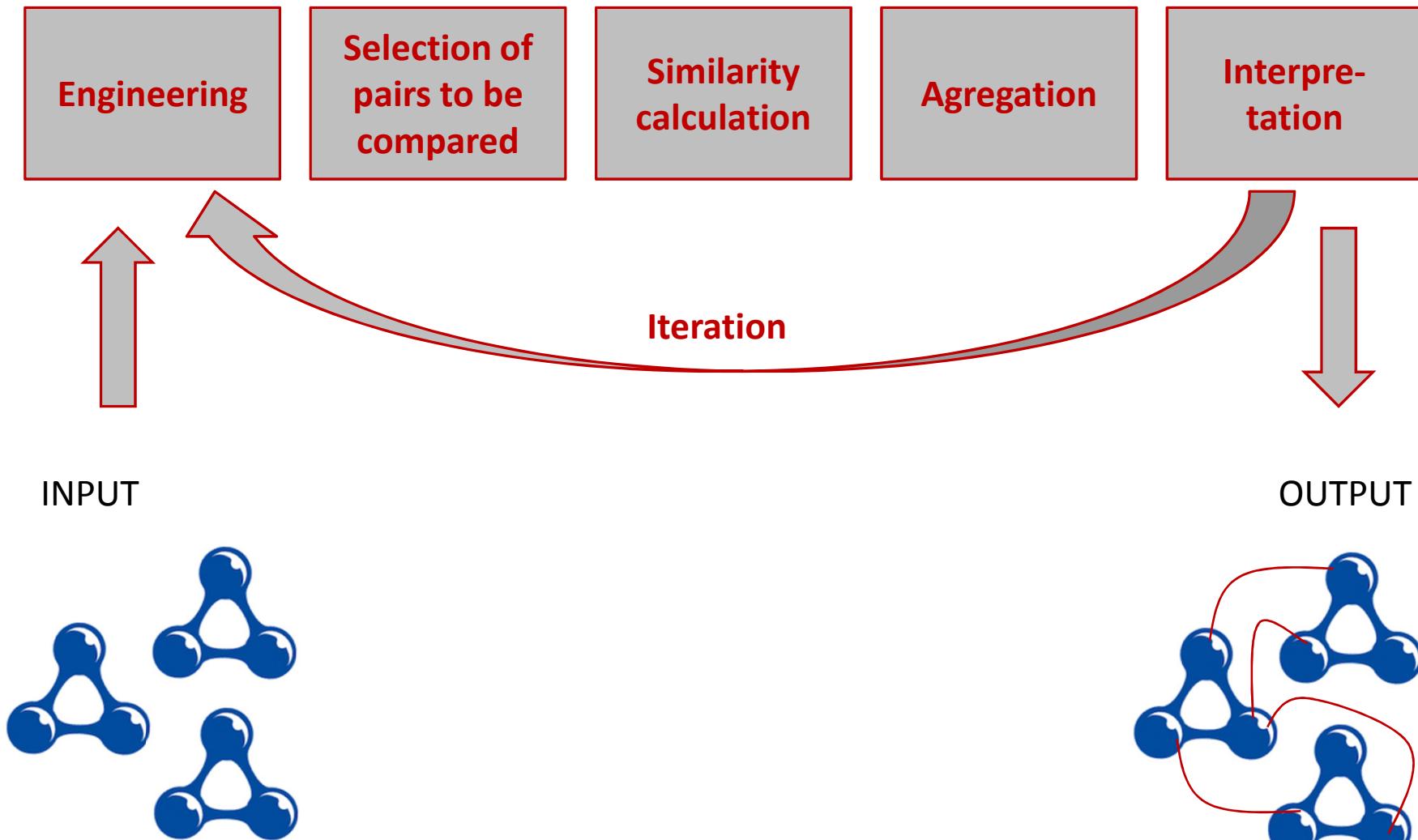
- to two ontologies o and o' ,
- to a set of original matching A ,
- from a set of parameters applicable to p alignment techniques implemented in the process
- and a set of external resources r ,

and that produces a set of matching A' between the two ontologies.

[Euzenat and Shvaiko, 2007]

$$A' = f(o; o'; A; p; r)$$





- Jaccard distance [Jaccard, 1901]
- n-grammes technique [Shannon, 1948]
- Levenshtein distance [Levenshtein, 1966]
 - minimal cost of edition operations necessary to transforme
- Lin mesure [Lin 1998]
- etc.

$$Sim_{Lin}(x, y) = 2 \frac{\sum_{t \in (tri(x) \cap tri(y))} logP(t)}{\sum_{t \in tri(x)} logP(t) + \sum_{t \in tri(y)} logP(t)}$$

For a long time, tools were created.
SILK is an interesting implementation.



The lifting and linking of Data in order to publish 5 stars Data (DataLift)

1. [The basic principles of the Web of Data: from raw and heterogeneous Data to the lingua franca of the Web of Data](#)
2. [The RDF semantic and its main syntaxes](#)
3. [Handling and querying Data \(triple store, endpoint, SPARQL\)](#)
4. [Understanding the basics of knowledge modeling \(ontology, OWL\)](#)
5. [The lifting and linking of Data in order to publish 5 stars Data \(DataLift\).](#)



The 5 stars of TBL



2010: At the "Gov 2.0 Expo" Tim Berners-Lee presented a wide opening became famous data that refers

5 stars



Available on the web (whatever format) but with an open licence, to be **Open Data**



Available as **machine-readable structured data** (e.g. excel instead of image scan of a table)



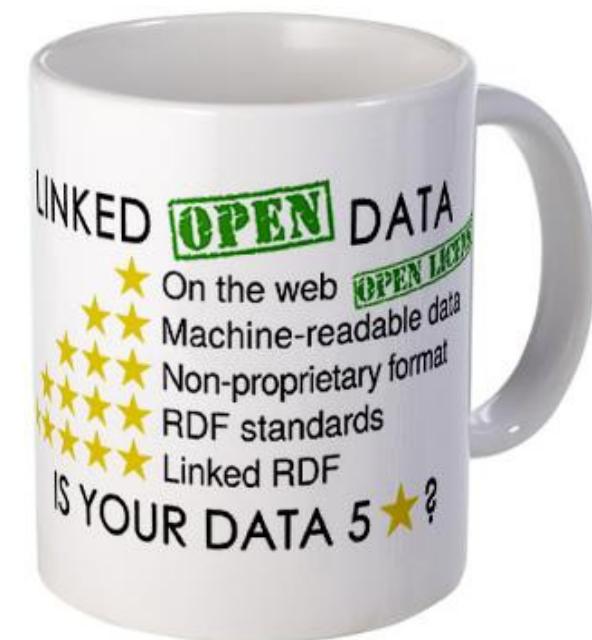
as (2) plus **non-proprietary format** (e.g. CSV instead of excel)



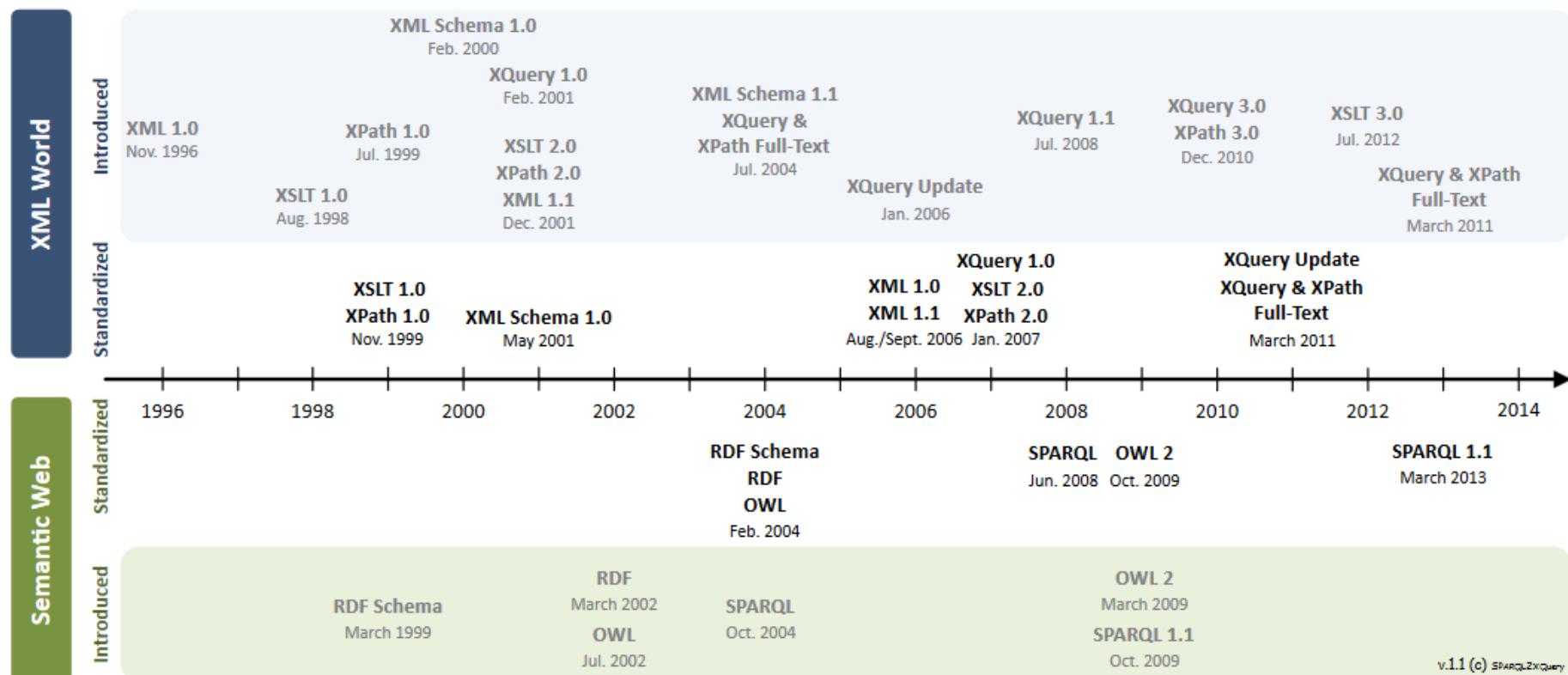
All the above plus, Use **open standards from W3C** (RDF and SPARQL) to identify things, so that people can point at your stuff



All the above, plus: Link your data to other people's data to **provide context**



4 stars – Standards & norms



4 stars – shared vocabularies

Subject	predicate	object
ASIN:B00005014I	dc:type	enregistrement musical
ASIN:B00005014I	dc:format	DVD
ASIN:B00005014I	dc:title	'Variations Goldberg' @fr
ASIN:B00005014I	bibo:performer	Glenn Gould
ASIN:B00005014I	mo:composer	Johannn-Sebastian Bach

The predicates have a **declared semantic** in **shared vocabularies** accessible via their URI

dc:	Dublin Core Terms
bibo:	Bibliographic Ontology
mo:	Music Ontology

5 stars – shared referentials

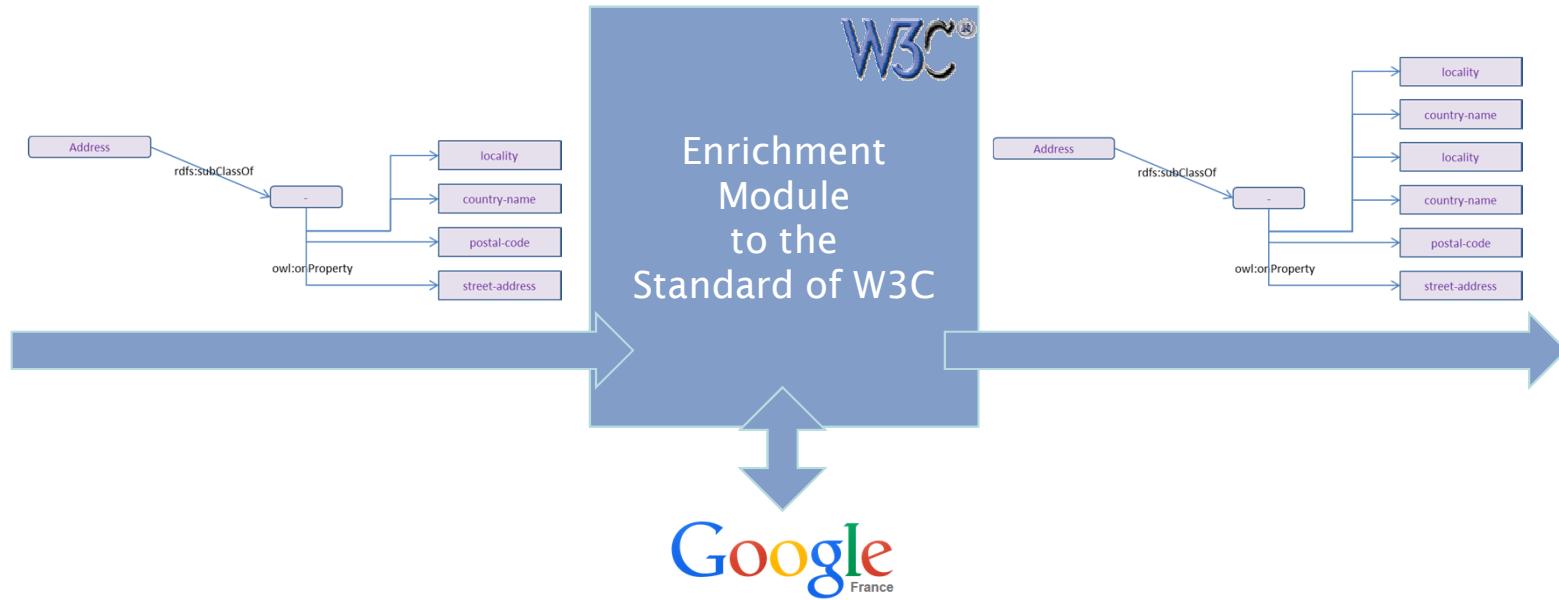
The description uses **entities of reference defined by URI**

subject	predicate	object
ASIN:B00005014I	dc:type	http://schema.org/MusicRecording
ASIN:B00005014I	dc:format	http://purl.org/spar/fabio/dvd
ASIN:B00005014I	bibo:performer	http://viaf.org/viaf/54148399/
ASIN:B00005014I	mo:composer	http://viaf.org/viaf/12304462/

The data structure does not change with respect to the original description

A demonstration





Fichier Accueil Insertion Mise en page Formules Données Révision Affichage

Calibri 11 A A Standard

Police Alignement Nombre

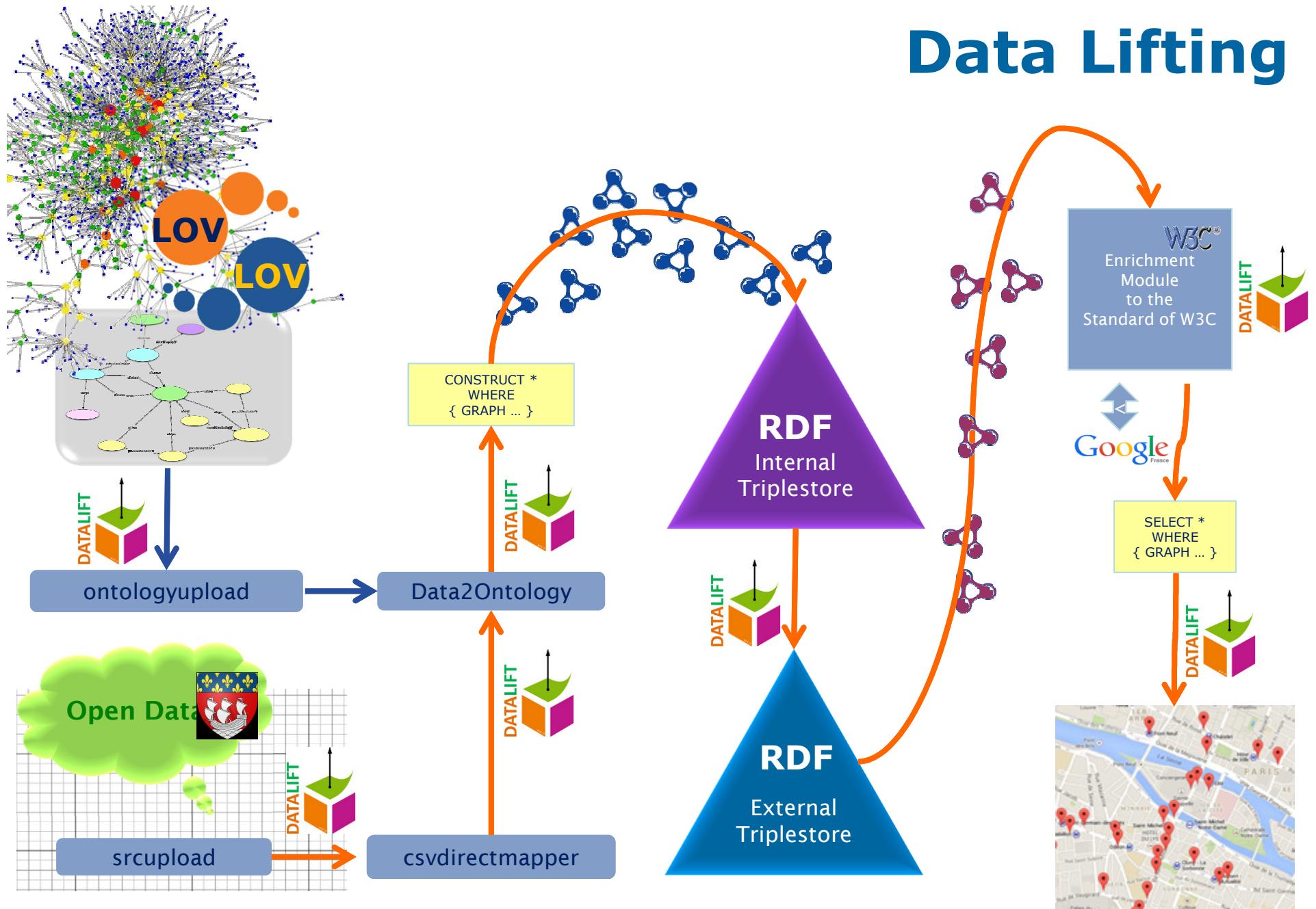
A6 f PCE DU PALAIS ROYAL / ST HONORE

A	B	C	D	E	F
1 ADRESSE	ARRDT	OUV.	FERM.	OUV. DIM	
2 PCE DU CHATELET	75001	08H30	20H00	OUI	
3 8 PCE DU PALAIS ROYAL / RIVOLI	75001	07H30	19H00	OUI	
4 8 RUE DE L'AMIRAL DE COLIGNY	75001	06H30	19H00	OUI	
5 2 BLD DU PALAIS	75001	08H00	19H30	MAT	
6 PCE DU PALAIS ROYAL / ST HONORE	75001	07H00	21H00	OUI	
7 2 QUAI DU LOUVRE	75001	08H00	19H30	MAT	
8 2 PCE STE OPPORTUNE	75001	06H30	20H00	MAT	
9 1 BLD DE LA MADELEINE	75001	06H45	19H30	MAT	
10 104 RUE DE RIVOLI	75001	08H00	20H00		
11 39 RUE CROIX DES PETITS CHAMPS	75001	07H00	17H30	MAT	
12 284 RUE ST HONORE	75001	07H15	19H45		

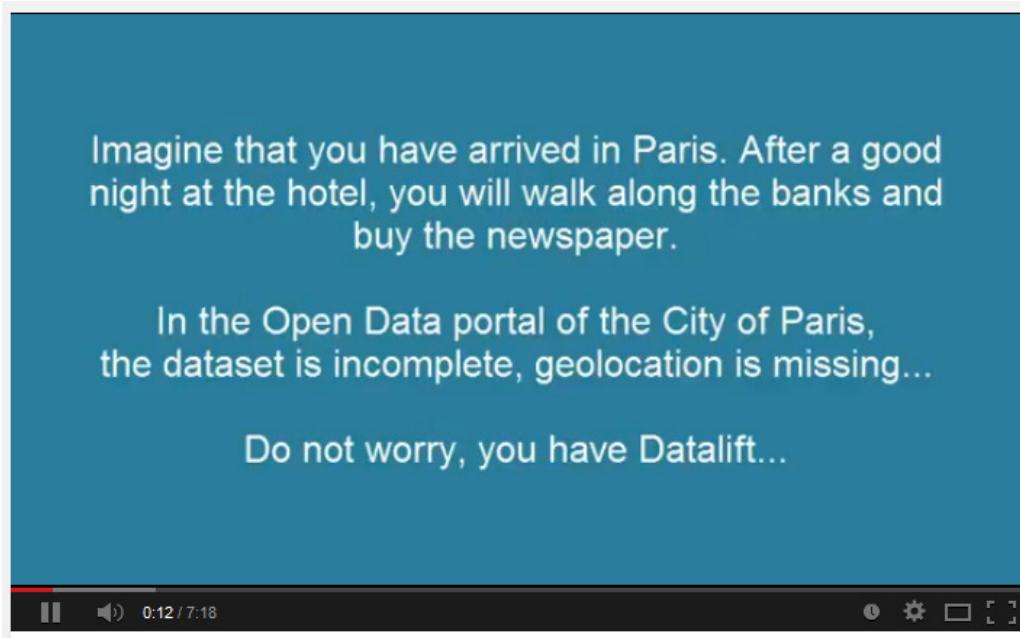
The data are « flat »

And incomplete for geolocalisation

Data Lifting



Using **Datalift**, the all-in-one linked data production platform



<http://www.youtube.com/watch?v=dMvEPpUZJqE>

Demo on youtube

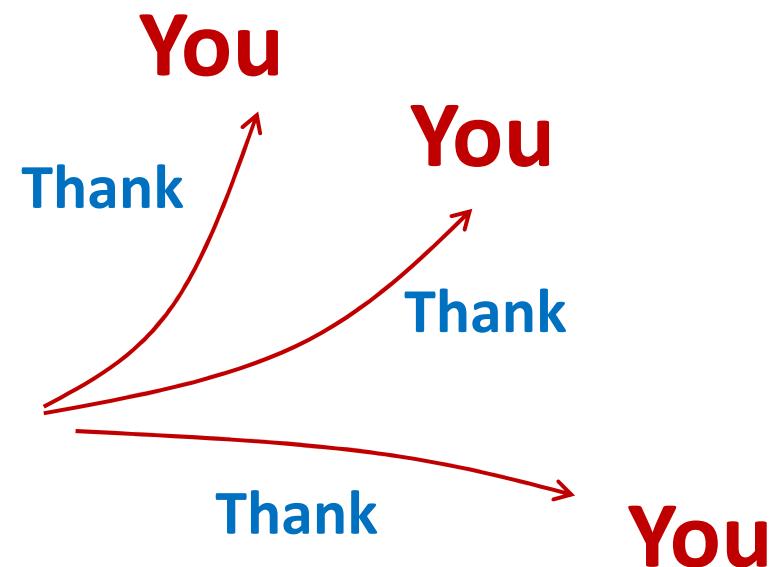
Datalift is an open source plateform
www.datalift.org

**DataLift is also an association,
which is freely open to all students.**



Fourth European Business Intelligence Summer School (eBISS 2014)

July 6 - July 11, 2014 Berlin, Germany



Gabriel KEPEKLIAN
R&D director
Atos Integration France

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Atos