

Introduction to Graph Databases

NoSQL and Graph Databases

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Agenda

- 10.10.22. Introduction – Graph data models
- 13.10.22. Graph DB internals. Introduction to Neo4j
- 17.10.22. Querying Neo4j databases
- 20.10.22. Assignment 1. Graphs in relational databases
- 24.10.22. Assignment 2. Basic Cypher queries
- 27.10.22. Assignment 3. Advanced Cypher queries

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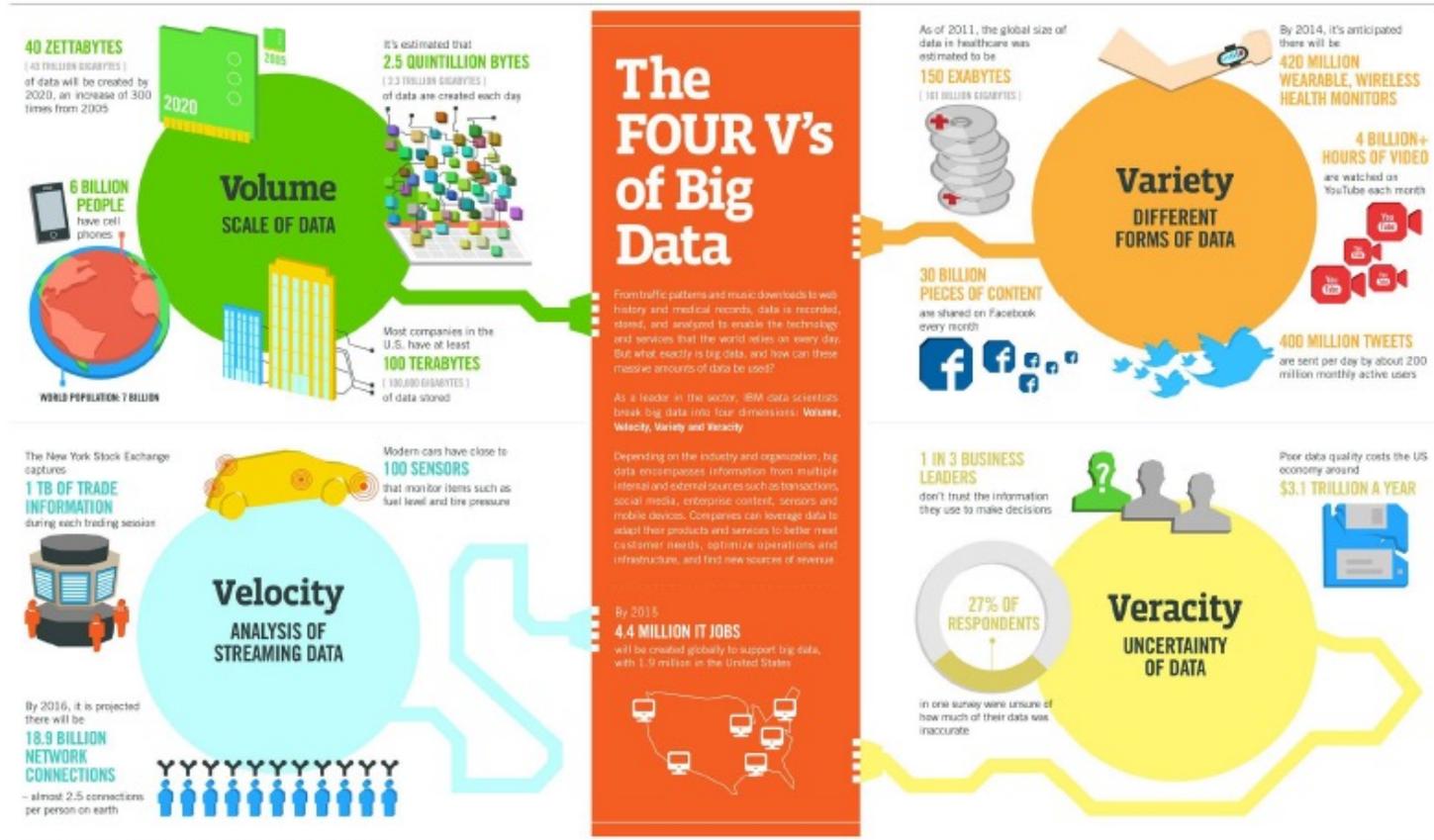
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Motivation

Typical BI scenario years ago ...

- **Early 90's**: Data Warehousing + Data Mining
- Big data: GB...TB!!
- Structured data
 - Mostly relational
 - Spreadsheets
 - (Some) Text
 - Web still in its infancy
- Problem: Data integration
- **Today**: Data deluge on the Web
- Daily PB of data of different kinds
 - Geographic
 - Text
 - Video, image
 - Audio

Volume, Velocity, Variety, Veracity



Main characteristic of these data

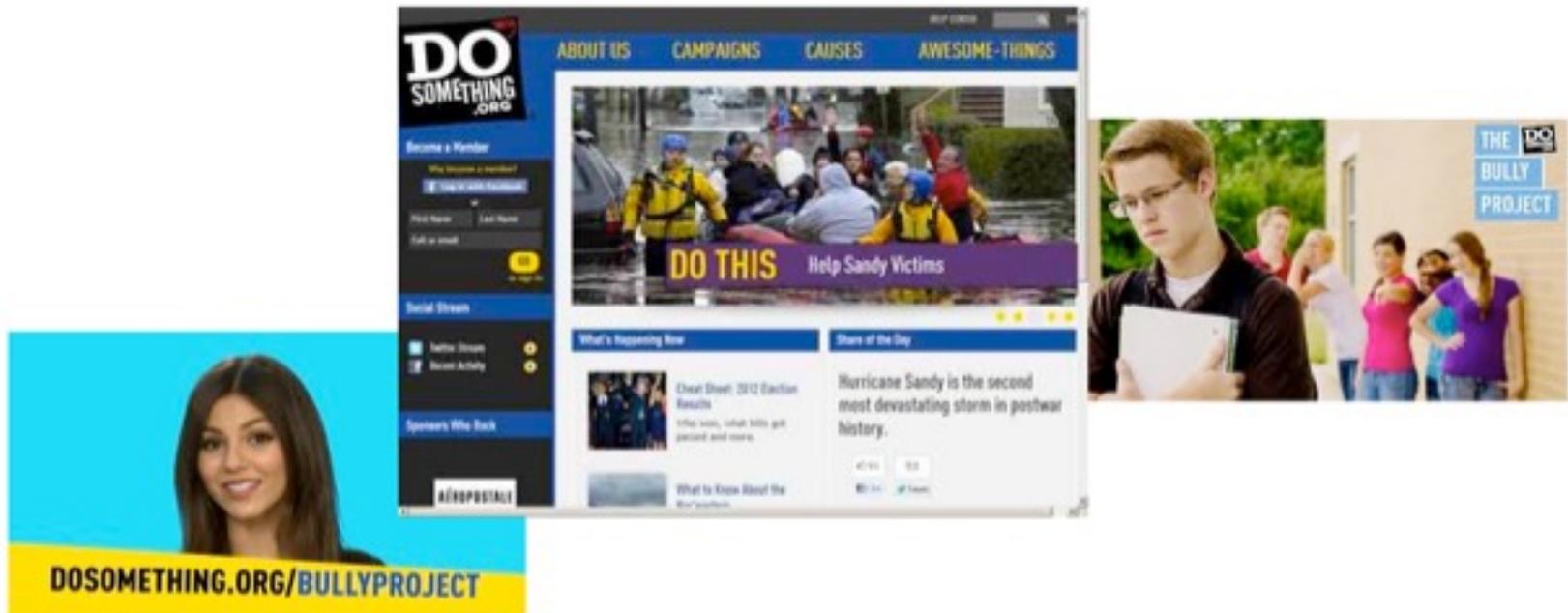


Social media monitoring



Social media monitoring

“Bullying” analysis on social networks



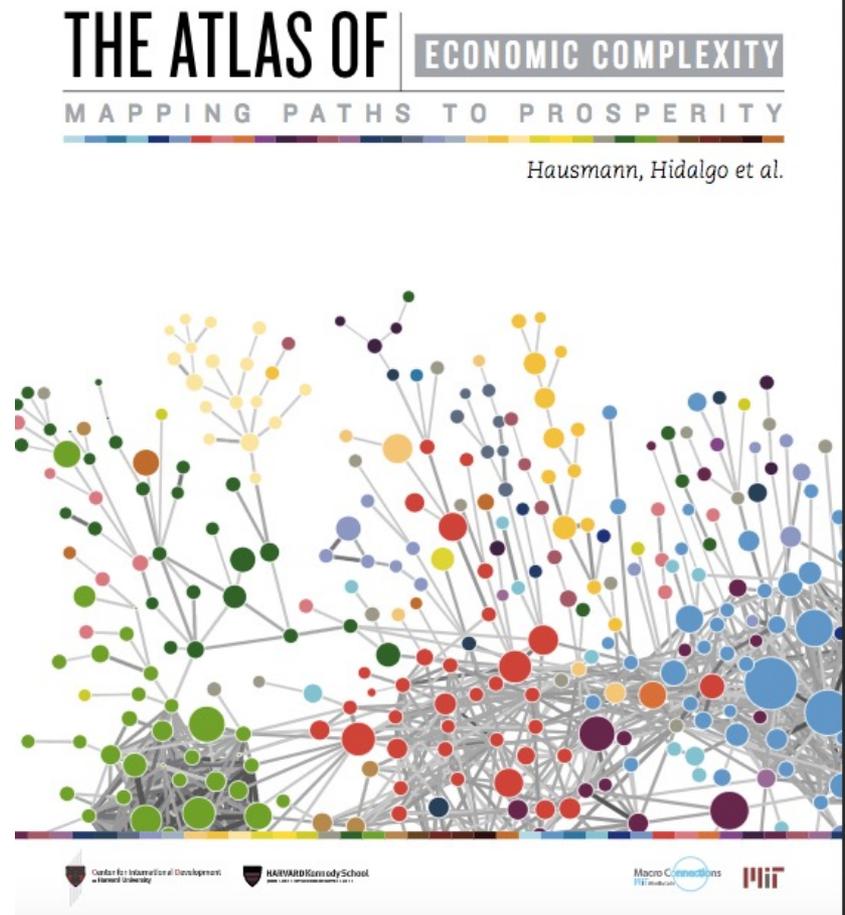
A world of interrelated information

Data of 200 countries, 750 products, 50 years.
Visualizations produced by MIT media lab

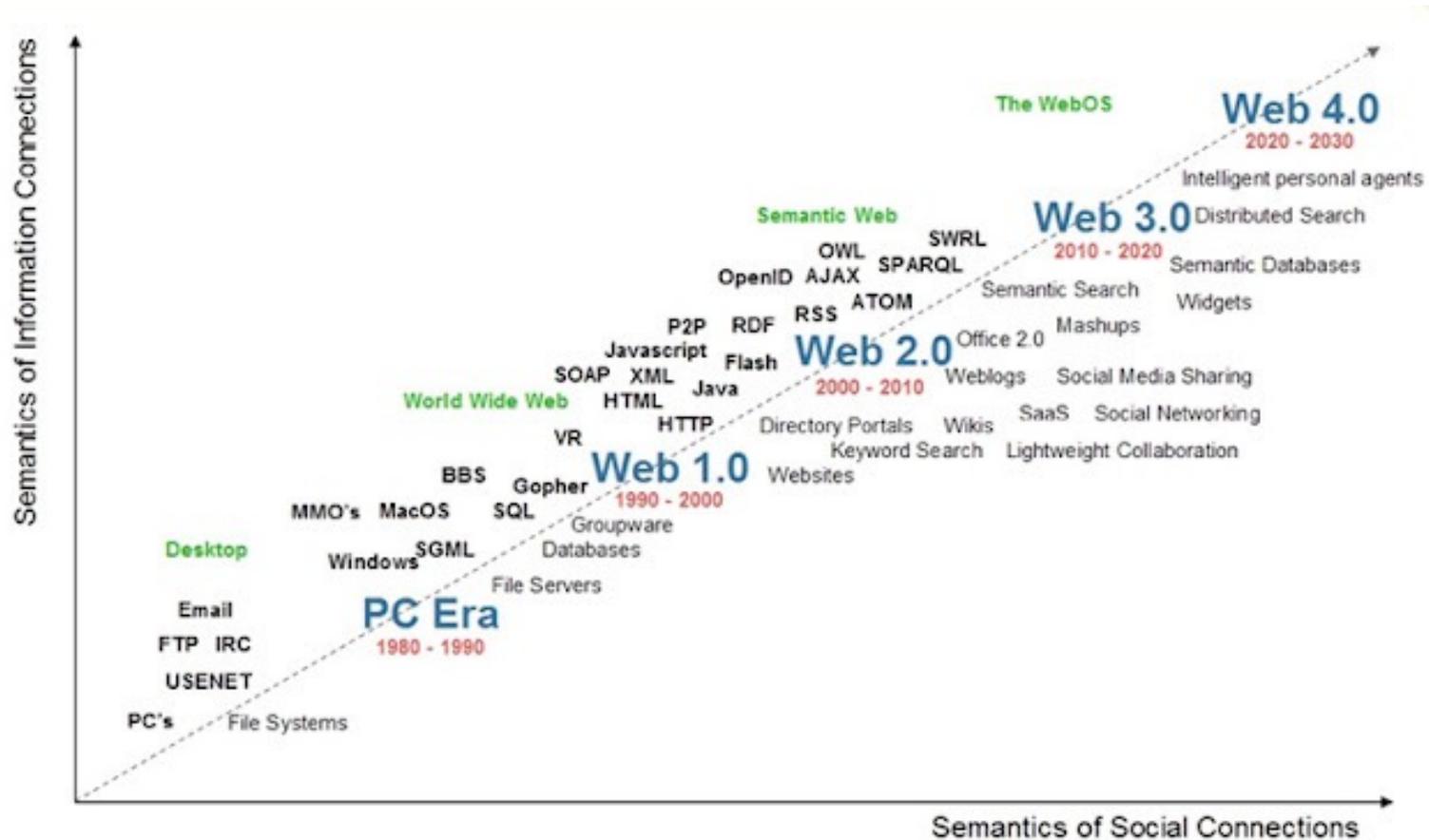
Indicators:

- Capability Distance (countries start making products “close” to the ones they already make, the “chicken and egg” problem)
- Complexity outlook index (how many new products could a country produce)
- Diversity (different products produced by a country)
- Economic complexity (measures how much of the knowledge of the society is transferred to the products it produces)

<http://atlas.cid.harvard.edu/>



Trend : Connectedness



How do we deal with this?

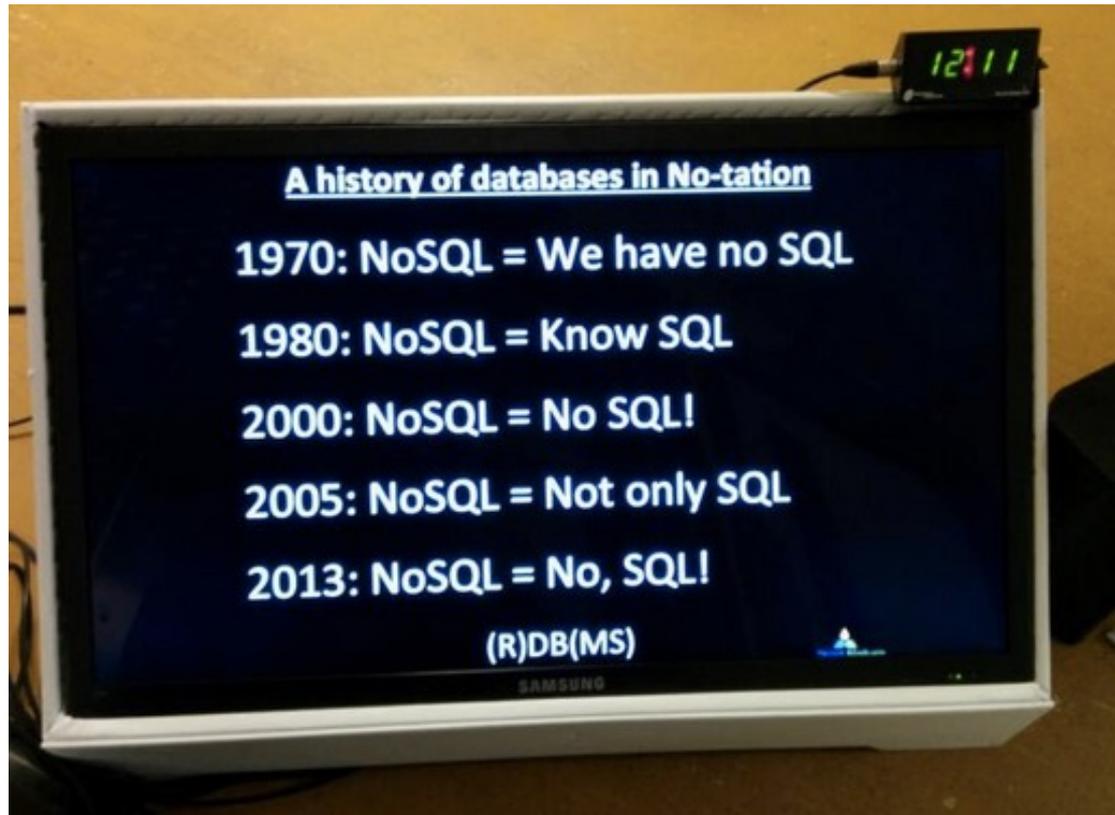
Is traditional DB technology enough?

We must address:

- Connectedness
- Unstructured data
- High Volumes
- Real-time

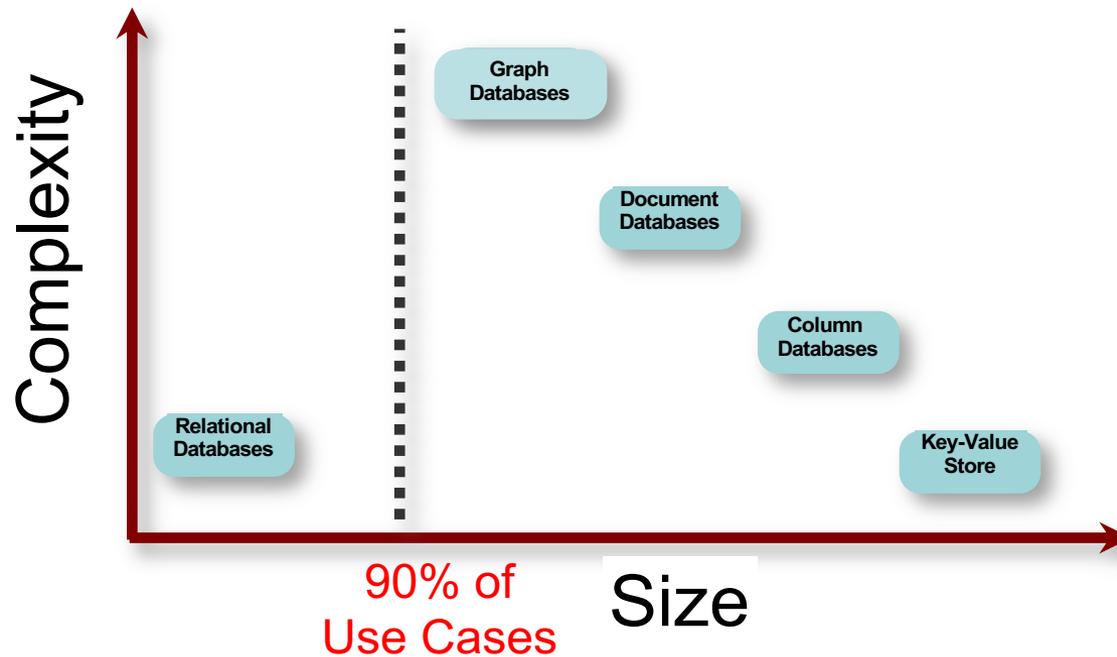
NoSQL technologies

The NoSQL paradigm



- Is SQL the future of NoSQL?

Living in a NoSQL world



NoSQL Motivation

- RDBMS **too rigid** for Big Data scenarios
- Not the best to store **unstructured data**
- **One-size-fits-all approach** no longer valid in many scenarios
- RDBMS **hard to scale** for billions of rows
- Data structures used in RDBMS optimized for systems with **small amounts of memory**

NoSQL characteristics

NoSQL characteristics

- **Not** using the **relational model** for storing data

NoSQL characteristics

- **Not** using the **relational model** for storing data
- **Not** using **SQL** for retrieving data

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- **Not** using **SQL** for retrieving data
- **No schema**, allowing fields to be added to any record, without control

NoSQL characteristics

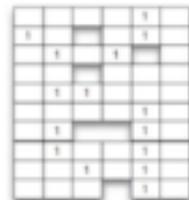
- **Not** using the **relational model** for storing data
- **Not** using **SQL** for retrieving data
- **No schema**, allowing fields to be added to any record, without control
- Ability to run on **clusters** of commodity hardware
- Ability to web-scale, with **horizontal scalability** in mind
- Trade-off traditional consistency for other useful properties (e.g., **no ACID** support most of the time)

Types of NoSQL Stores

Key-Value



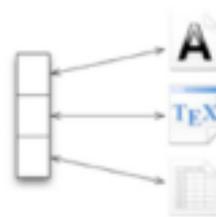
Column



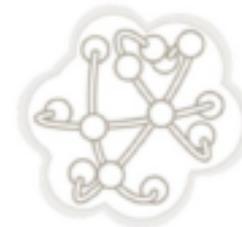
APACHE
HBASE



Document



Graph

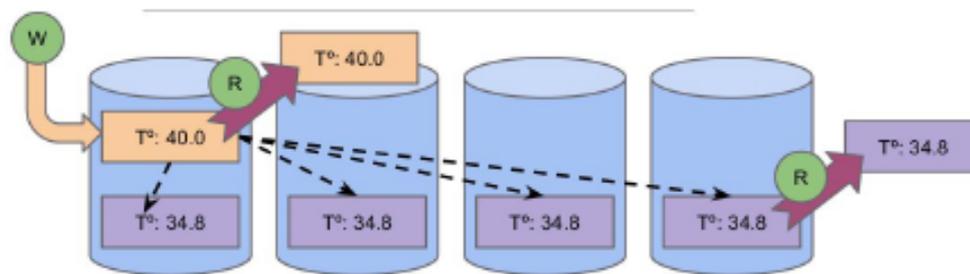


Recall the CAP theorem

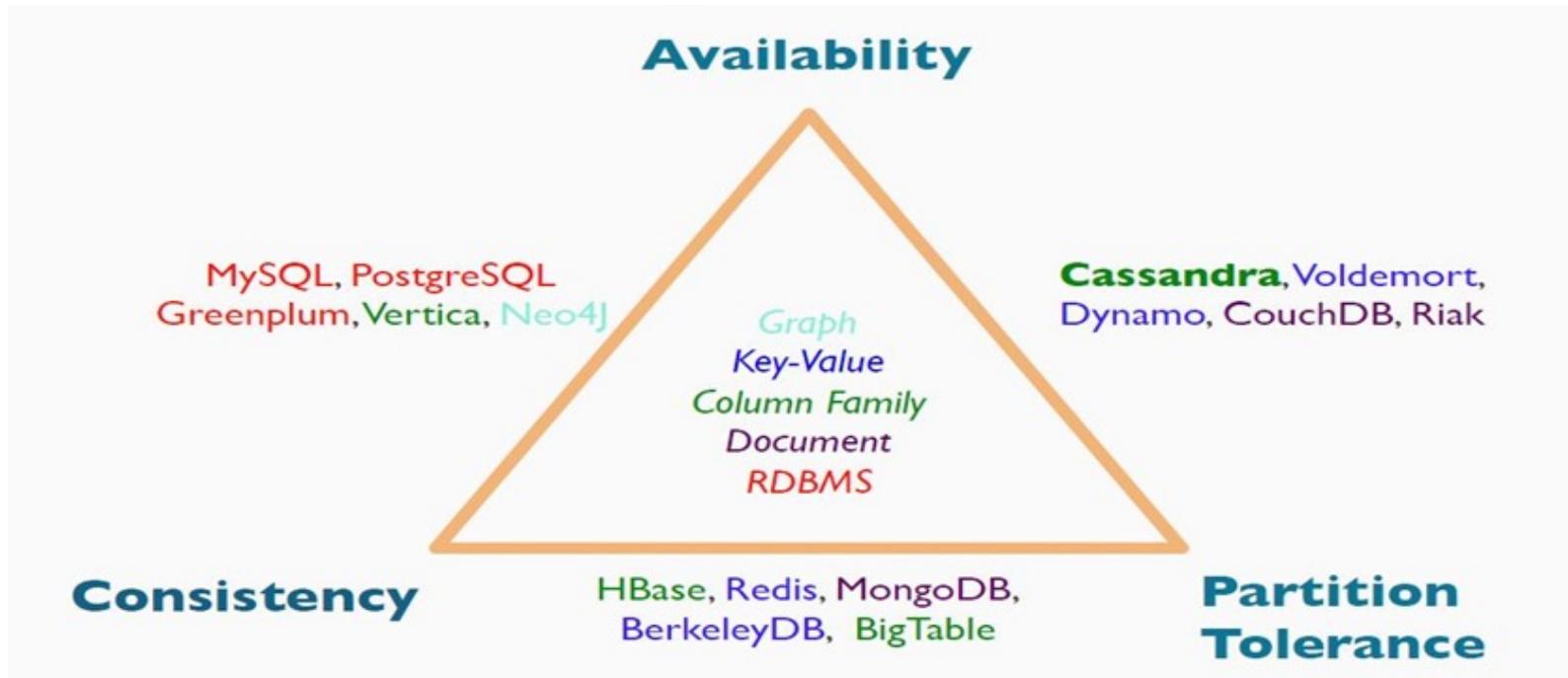
- **Consistency:** Strong consistency of updates
- **Availability:** Guarantees that every request receives a response whether it succeeded or failed when retrieving data
- **Partition tolerance:** The system continues to operate despite arbitrary message loss or failure of part of the system
- **CAP Theorem:** A distributed data system cannot guarantee the three properties above, but only any combination of two of them.
- In other words, the CAP theorem states that in the presence of a network partition, one has to choose between consistency and availability.
- For example, a RDBMS can only guarantee **CA**

Eventual consistency

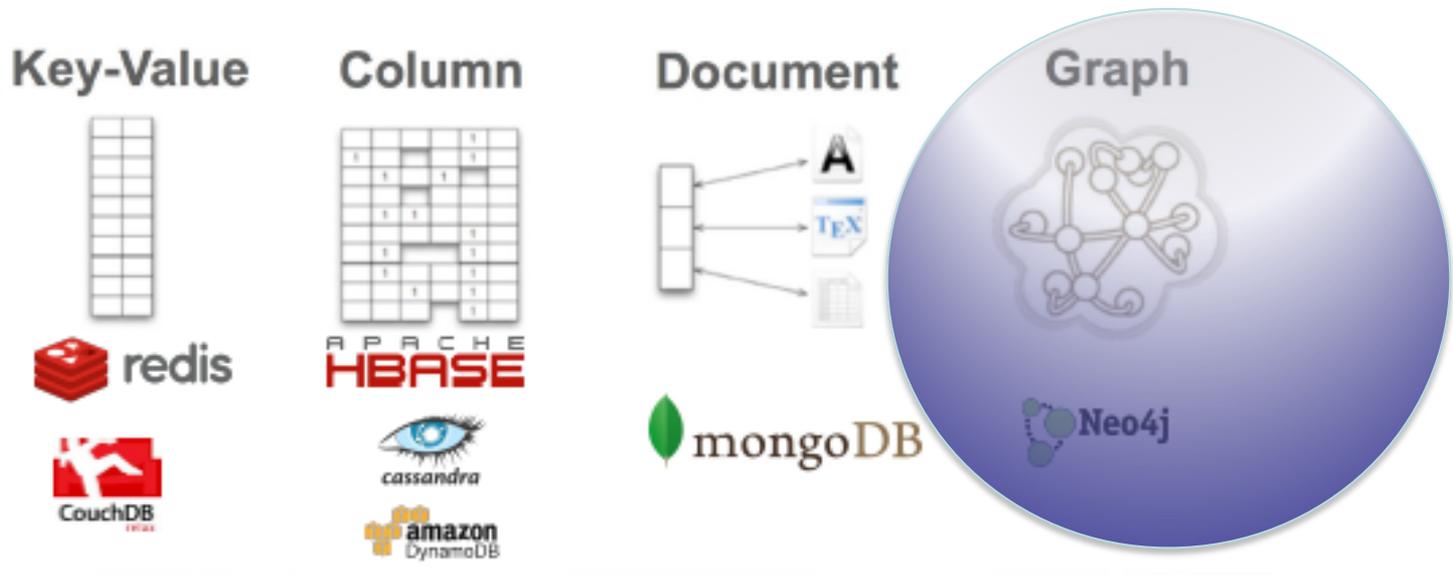
- **Assumption:** in the absence of new writes, **eventually** consistency will be achieved, and all replicas that are responsible for a data item will agree on the same version and return the last updated value.
- With fewer replicas, R/W operations complete more quickly, lowering latency.



CAP theorem

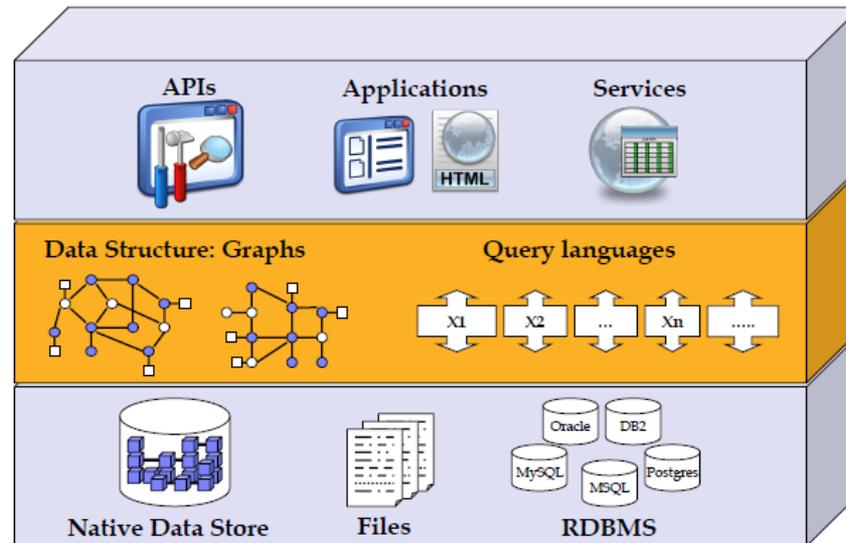


In the remainder....



Introduction to Graph Databases

Architecture



But first something some reminders

- Database models (Codd)

Data structures

Integrity constraints

Query Language

Database models

| Database model | Abstraction level | Data structure | Information focus |
|-----------------------|--------------------------|-----------------------|--------------------------|
| Network | Physical | Pointers, records | Records |
| Relational | Logical | Relations | Data, attributes |
| Semantic | User | Graph | Schema, relations |
| OO | Physical/logical | Objects | Objects, methods |
| Semi-structured | Logical | Tree | Data, components |
| Graph | Logical/user | Graph | Data, relations |

Database models: graphs

Data structure

Data and/or schema are represented by graphs, or by data structures generalizing the notion of graph (hypergraphs or hypernodes)

Database models (Codd)

Integrity constraints

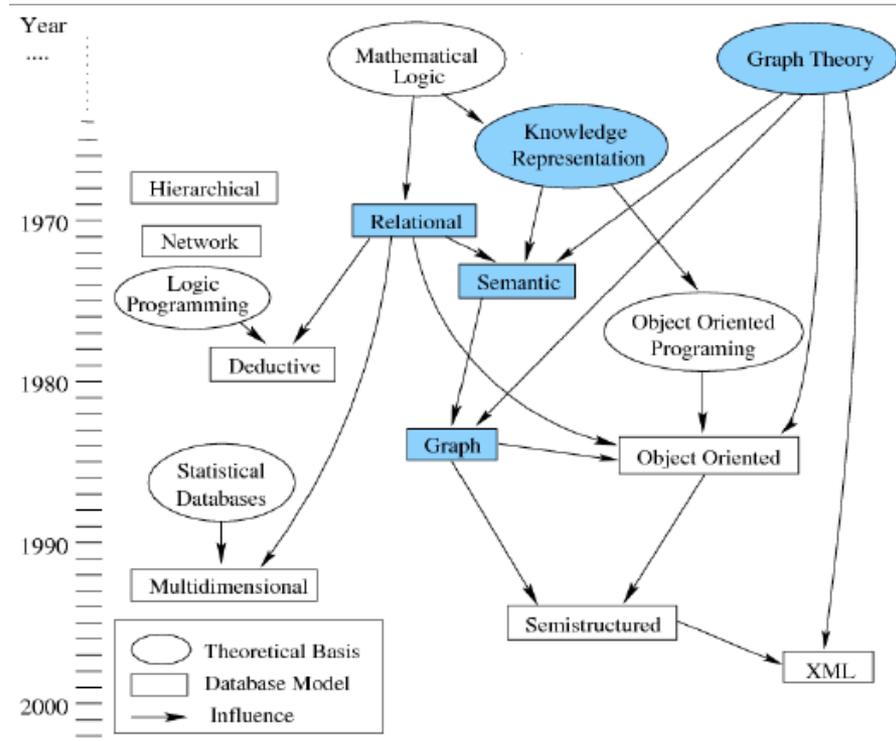
Enforce data consistency. Constraints can be grouped in schema-instance consistency, identity and referential integrity, and functional and inclusion dependencies

Database models: (Codd)

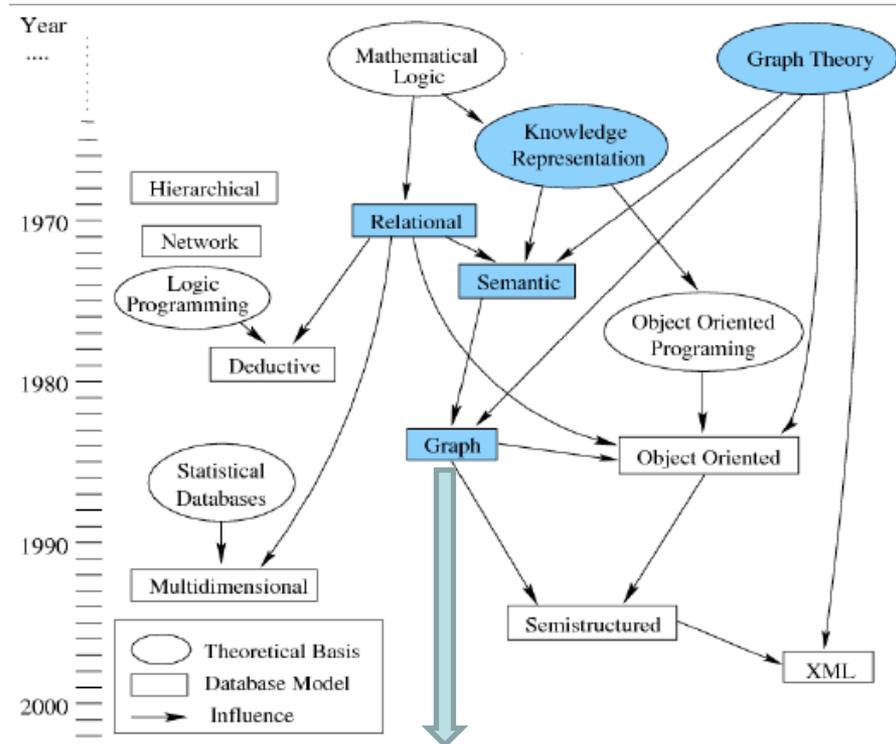
Query language

Data manipulation is expressed by graph transformations, or by operations whose main primitives are on graph features, like paths, neighborhoods, subgraphs, connectivity, and graph statistics.

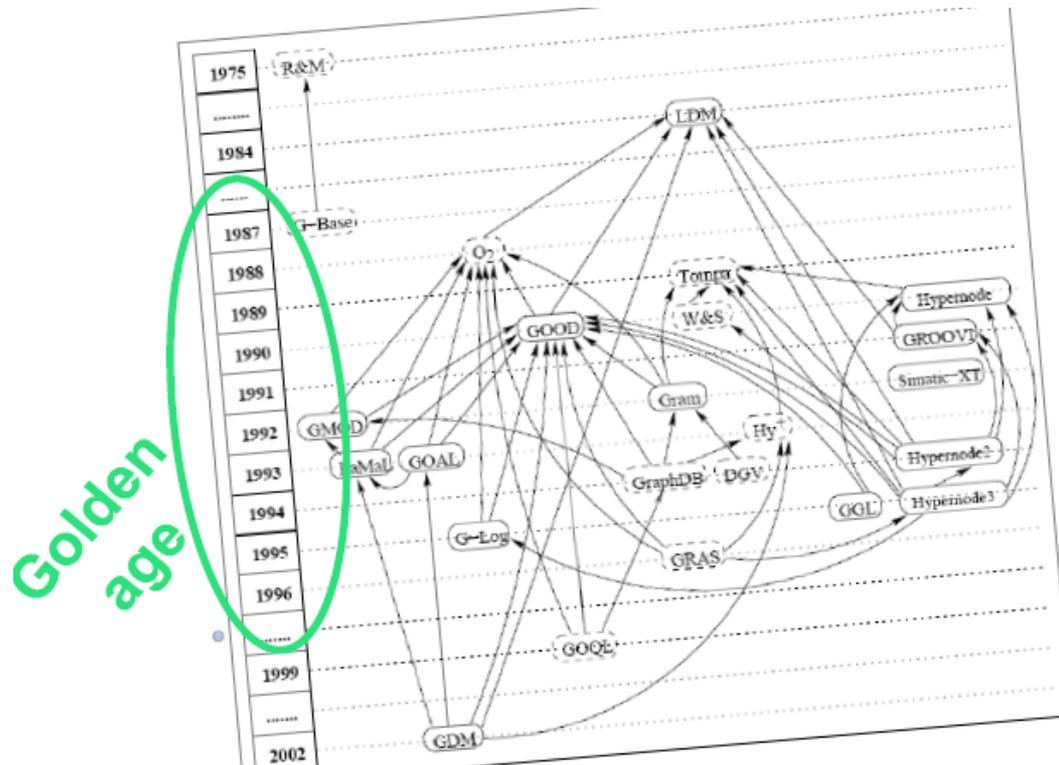
A history of database models (A. Mendelzon)



A history of database models (A. Mendelzon)

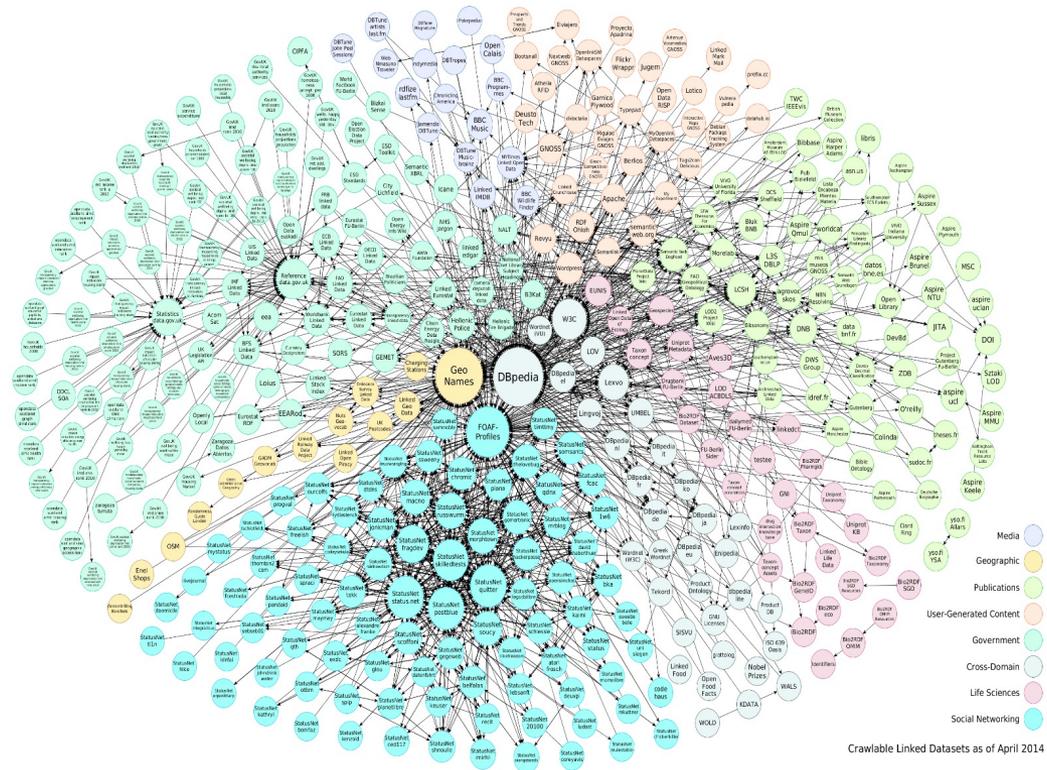


The Golden age of GDB



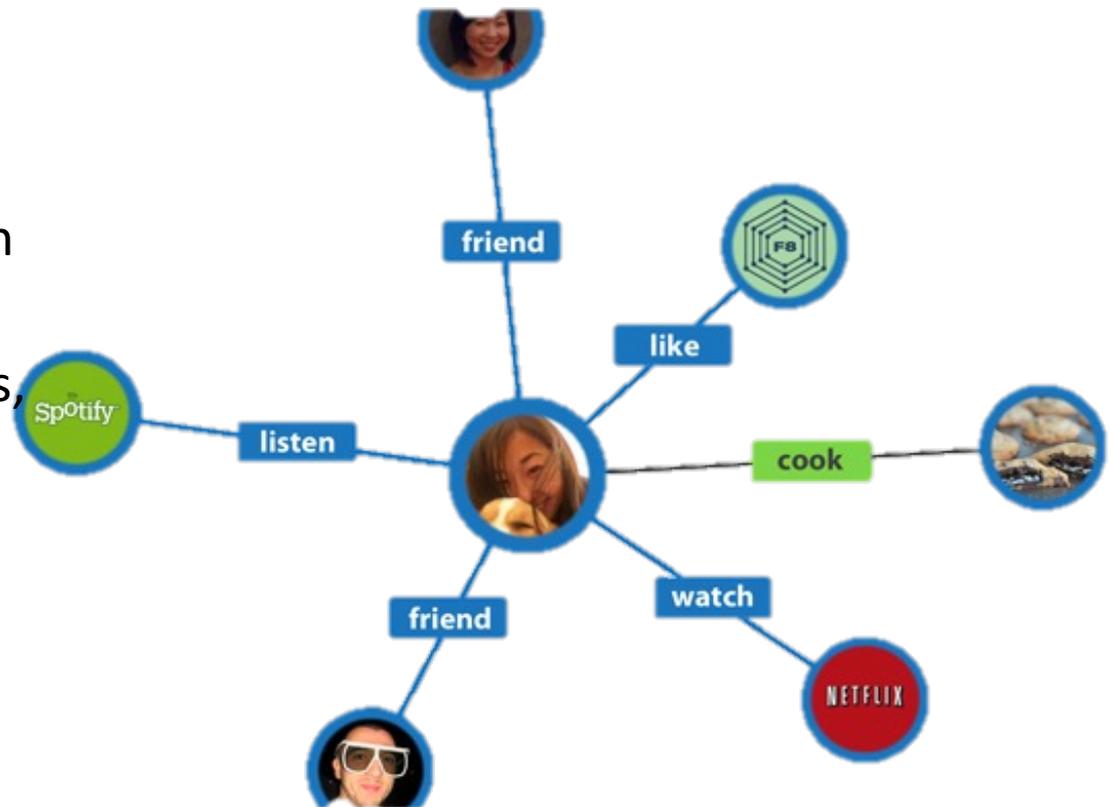
Data connectedness

- Facebook Graph
- LinkedIn Graph
- Linked Data
- Blogs/Tagging



Data connectedness

- Modeling the FB Graph
- Persons, friendships, photos, locations, apps, pages, ads, interests, age range, etc.



Social Network “path” performance

- Experiment:

- ~1k persons
- Average 50 friends per person
- `pathExists(a, b)` limited to depth 4
- Caches warm to eliminate disk IO

| | # persons | query time |
|---------------------|-----------|------------|
| Relational database | 1000 | 2000ms |

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|---------------------|-----------|------------|
| Relational database | 1000 | 2000ms |
| Neo4j | 1000 | 2ms |

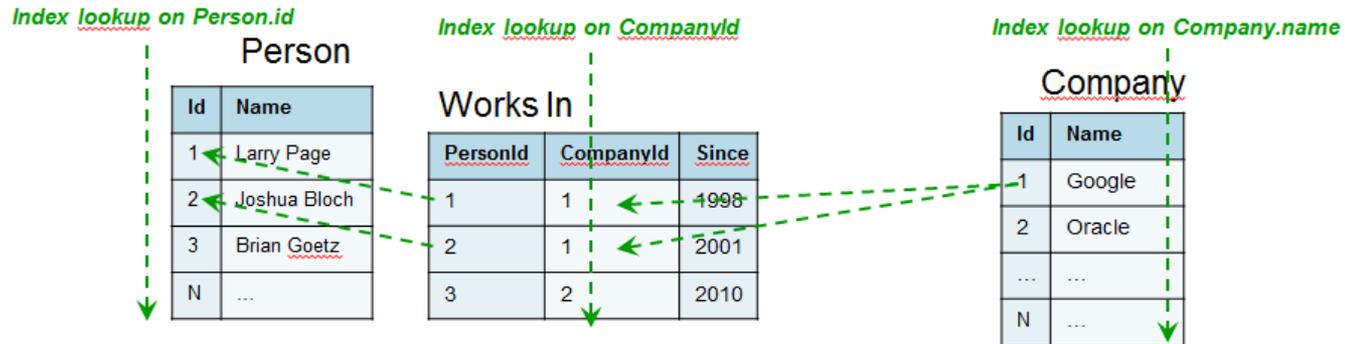
Social Network “path” performance

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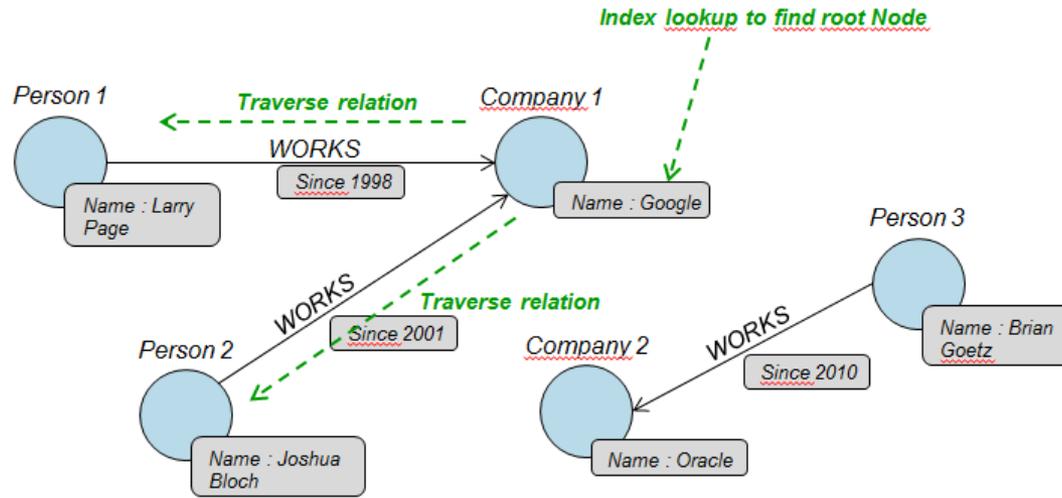
| | # persons | query time |
|---------------------|-----------|------------|
| Relational database | 1000 | 2000ms |
| Neo4j | 1000 | 2ms |
| Neo4j | 1000000 | 2ms |

A typical SQL query



```
Select Person.Name
from Person, Company, WorksIn
where Company.name='Google'
and WorksIn.CompanyId = Company.Id
and WorksIn.PersonId = Person.Id
```

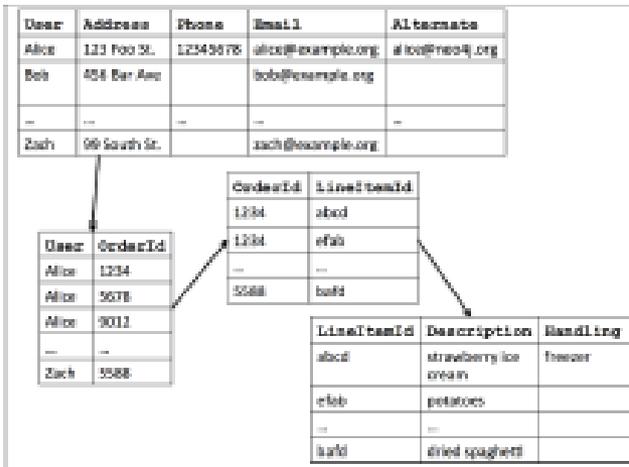
Same query on graphs



The deepest the navigation, the largest the difference with RDBs

Traversing data in a RDBMS

- Based on joining and selecting data



```
SELECT *
FROM user u, user_order uo,
orders o, items i
WHERE u.user = uo.user AND
uo.orderId = o.orderId AND
i.lineItemId = i.LineItemId
AND u.user = 'Alice'
```

Cardinalities:

|User|: 5.000.000
|UserOrder|: 100.000.000
|Orders|: 1.000.000.000
|Item|: 35.000

Query Cost?!

Traversing data in a GDB



Cardinalities:

|User| : 5.000.000
|Orders| : 1.000.000.000
|Item| : 35.000

Query Cost?!
O(N)

| ReportsTo | | | ReportsDirectlyTo | |
|-----------|----------|-------------|-------------------|----------|
| Boss | Employee | | Boss | Employee |
| A | C | Iteration 1 | A | C |
| A | D | | A | D |
| C | E | | C | E |
| C | M | | C | M |
| A | E | Iteration 2 | E | T |
| A | M | | | |
| C | T | Iteration 3 | | |
| A | T | | | |

We want to compute the closure of the relation “ReportsDirectlyTo”, that is, to whom someone “ReportsTo”, either directly or indirectly. SQL supports these kinds of recursive queries. Recursively joining ReportsTo and ReportsDirectlyTo on `RT.Employee=RDT.Boss`.

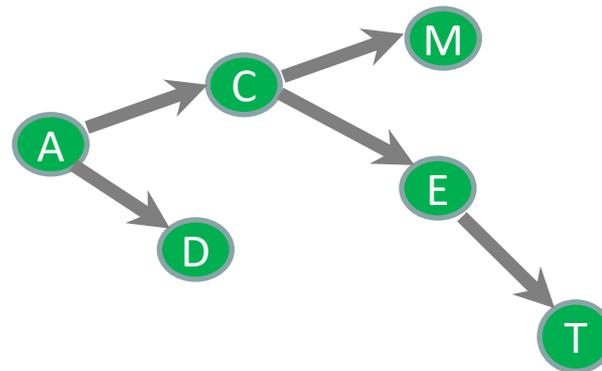
```
WITH recursive ReportsTo(Boss, Employee) AS
  (SELECT Boss, Employee
   FROM ReportsDirectlyTo
   UNION ALL
   SELECT ReportsTo.Boss, ReportsDirectlyTo.Employee
   FROM ReportsTo, ReportsDirectlyTo
   WHERE ReportsTo.Employee = ReportsDirectlyTo.Boss )
SELECT * FROM ReportsTo
```

These queries are normally more expensive in the Relational Model, since they imply **MULTIPLE JOINS**.

Joins are expressed at the schema level rather than at the instance level.

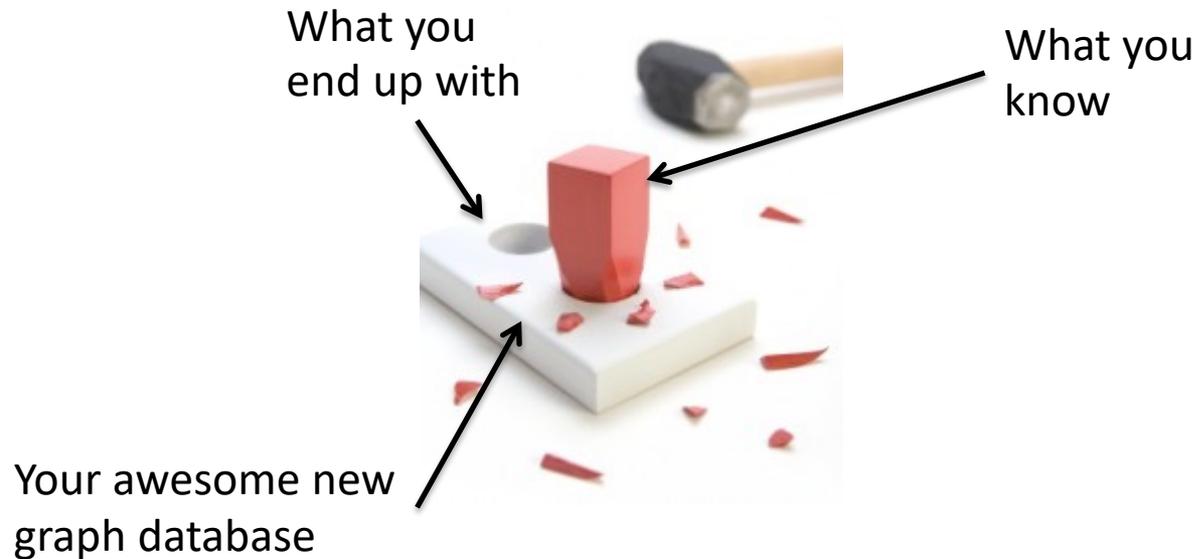
How would we represent this in a Graph Data Model?

| ReportsDirectlyTo | |
|-------------------|----------|
| Boss | Employee |
| A | C |
| A | D |
| C | E |
| C | M |
| E | T |



Paths in a graph are expressed at the instance level (there is no schema). Just check if there is an outgoing edge.

Forcing the relational model in a graph



Think in terms of nodes and edges!

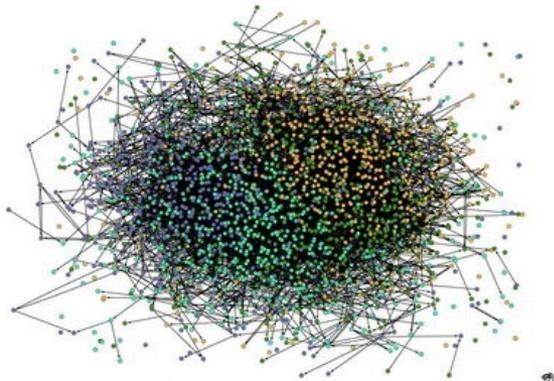
Consequence: graph DB are **hot** again

- The question is: how to process HUGE graphs?
- We now have hardware that can do this
- As usual
 - Software runs behind hardware;
 - Theory and models behind software
 - Current status: problem understanding

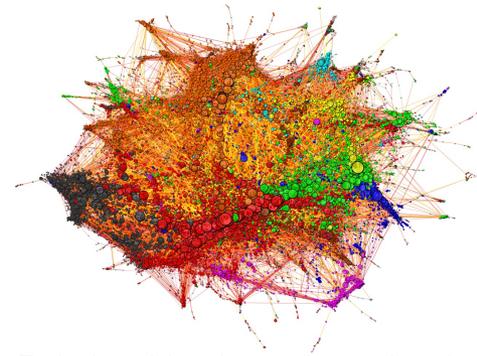
Open questions (almost all of the stuff)

- Use cases
- Data structures
- Query languages and operators
- Benchmarks
- Open world vs. closed world
- Centralized or distributed?
- Dynamics – transactions

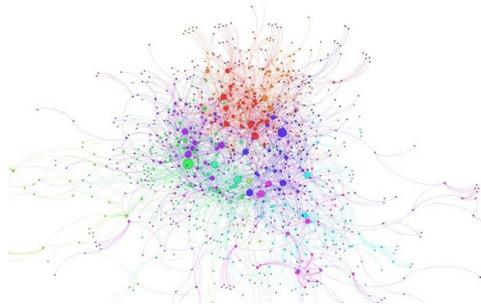
Use cases: social networks



Network of Friends in a High School 

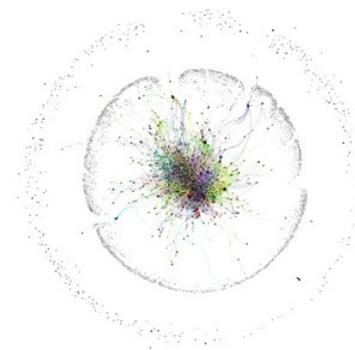


Relationships between artists in Last.fm
<http://sixdegrees.hu/last.fm/>



Network structure of music genres and their stylistic origin

http://www.infosysblogs.com/web2/2013/01/network_structure_of_music_gen.html

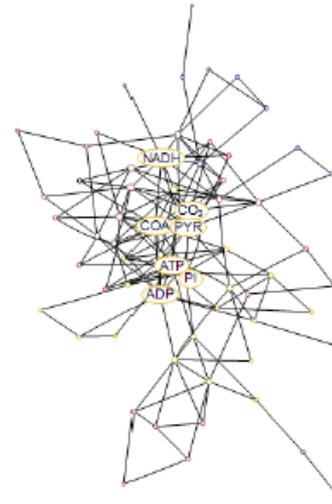


Network structure of Patent Citations

<http://www.infosysblogs.com/web2/2013/07/>

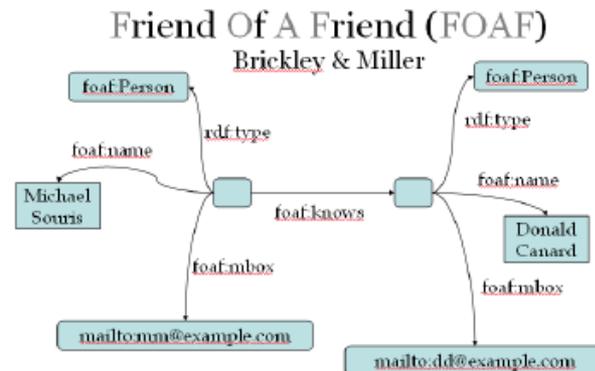
Use cases: biology

| Use Case | Graph Query |
|---|---------------------------------------|
| Chemical structure associated with a node | Node matching |
| Find the difference in metabolisms between two microbes | Graph intersection, union, difference |
| To combine multiple protein interaction graphs | Majority graph query |
| To construct pathways from individual reactions | Graph composition |
| To connect pathways, metabolism of co-existing organisms | Graph composition |
| Identify "important" paths from nutrients to chemical outputs | Shortest path queries |
| Find all products ultimately derived from a particular reaction | Transitive Closure |
| Observe multiple products are co-regulated | Least common ancestor |
| To find biopathways graph motifs | Frequent subgraph recognition |
| Chemical info retrieval | Subgraph isomorphism |
| Kinaze enzyme | Subgraph homomorphism |
| Enzyme taxonomies | Subsumption testing |
| To find biopathways graph motifs | Frequent subgraph recognition |

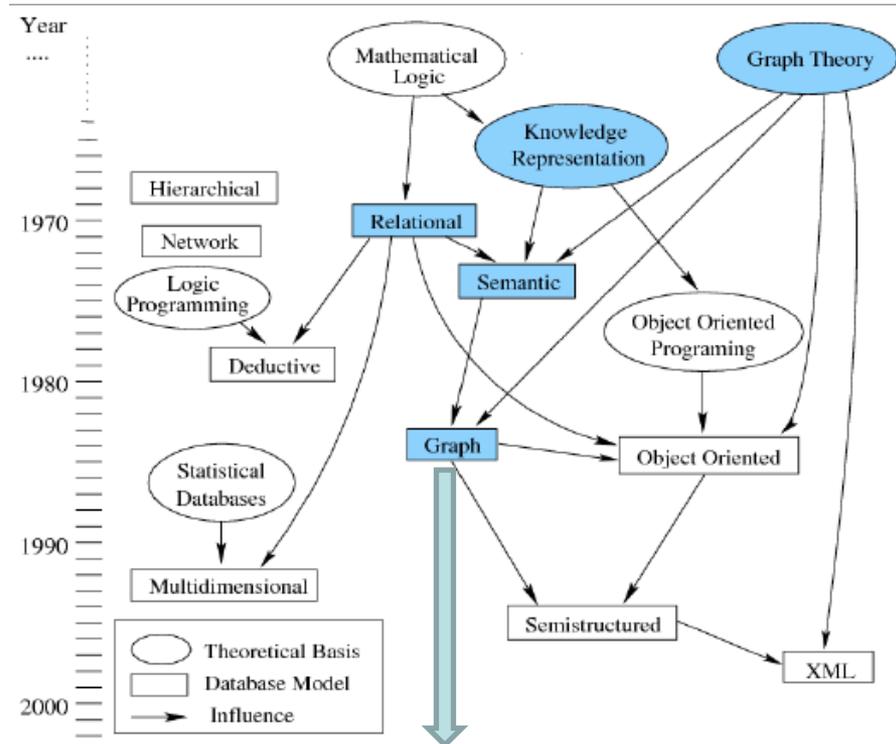


Use cases: the web

| Use Case | Graph Query |
|--|------------------|
| What is/are the most cited paper/s? | Degree of a node |
| What is the influence of article D? | Paths |
| What is the Erdős distance between authors X and author Y? | Distance |
| Are suspects A and B related? | Paths |
| All relatives of degree one of Alice | Adjacency |



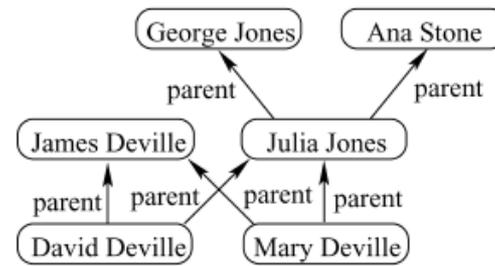
Graph database models



Graph database models

Example: genealogy and data diagram

| NAME | LASTNAME | PERSON | PARENT |
|--------|----------|--------|--------|
| George | Jones | Julia | George |
| Ana | Stone | Julia | Ana |
| Julia | Jones | David | James |
| James | Deville | David | Julia |
| David | Deville | Mary | James |
| Mary | Deville | Mary | Julia |



RDF Model

- RDF allows to express facts: Ana is the mother of Julia
- But we'd like to be able to express more generic knowledge
 - Mothers are female
 - If somebody has a daughter then that person is a parent
- This kind of knowledge is often called *schema* knowledge or *terminological* knowledge.
- RDF Schema allows us to do some schema knowledge modeling, OWL gives even more expressivity

RDF classes & instances

- Classes stand for sets of things. In RDF: Sets of URIs.
- `book:uri` is a member of the class `ex:Textbook`

```
book:uri    rdf:type    ex:Textbook .
```

- An URI can belong to several classes

```
book:uri    rdf:type    ex:Textbook .  
book:uri    rdf:type    ex:WorthReading .
```

- Classes can be arranged in hierarchies: each textbook is a book

```
ex:Textbook  rdfs:subClassOf  ex:Book .
```

RDF pre-defined classes

- Every URI denoting a class is a member of `rdfs:Class`

```
ex:Textbook    rdf:type    rdfs:Class .
```

- This also makes `rdfs:Class` a member of `rdfs:Class`

```
rdfs:Class    rdf:type    rdfs:Class .
```

- `rdfs:Resource` (class of all URIs)
- `rdf:Property` (class of all properties)
- `rdf:XMLLiteral`
- `rdfs:Literal` (each datatype is a subclass)
- `rdf:Bag`, `rdf:Alt`, `rdf:Seq`, `rdfs:Container` , `rdf:List`, `rdf:nil`,
`rdfs:ContainerMembershipProperty`
- `rdfs:Datatype` (contains all datatypes – a class of classes)
- `rdf:Statement`

RDF implicit knowledge

If an RDFS document contains

```
u    rdf:type    ex:Textbook .
```

and

```
ex:Textbook  rdfs:subClassOf  ex:Book .
```

then

```
u    rdf:type    ex:Book .
```

is *implicitly* also the case: it's a *logical consequence*. (We can also say it is *deduced* (deduction) or *inferred* (inference). We do not have to state this explicitly. Which statements are logical consequences is governed by the formal semantics.

RDF implicit knowledge (cont.)

As another example, from

```
ex:Textbook    rdfs:subClassOf    ex:Book .
```

and

```
ex:Book        rdfs:subClassOf    ex:PrintMedia .
```

then

```
ex:Textbook    rdfs:subClassOf    ex:PrintMedia .
```

I.e. `rdfs:subClassOf` is *transitive*.

RDF implicit knowledge (cont.)

Property hierarchies

From

```
ex:isHappilyMarriedTo rdf:subPropertyOf ex:isMarriedTo.
```

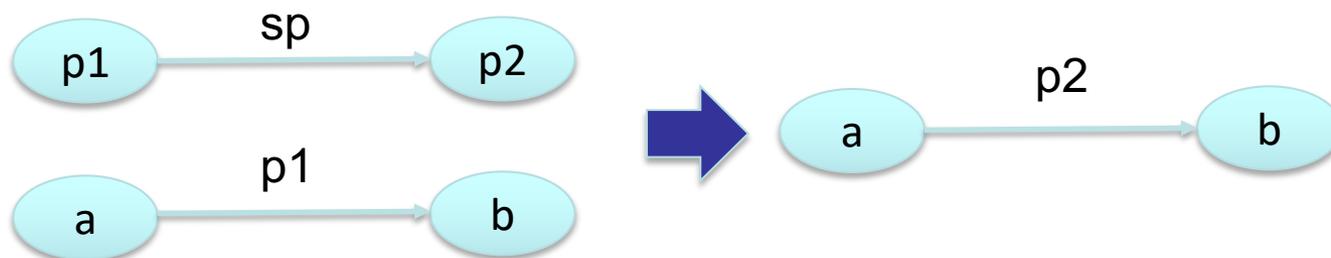
and

```
ex:markus ex:isHappilyMarriedTo ex:anja .
```

then

```
ex:markus ex:isMarriedTo ex:anja .
```

Thus, if p_1 is subproperty of p_2 , and a p_1 b , then a p_2 b



RDF model classes & properties

```

@prefix rdf:
  <http://www.w3.org/1999/02/22-rdf-syntax-ns#>.
@prefix rdfs:
  <http://www.w3.org/2000/01/rdf-schema#>.
@prefix :
  <http://example.org/ns#>

:Post      rdf:type      rdfs:Class.

:Tweet     rdf:type      rdfs:Class;
           rdfs:subClassOf :Post.

:FBPost    rdf:type      rdfs:Class;
           rdfs:subClassOf :Post.

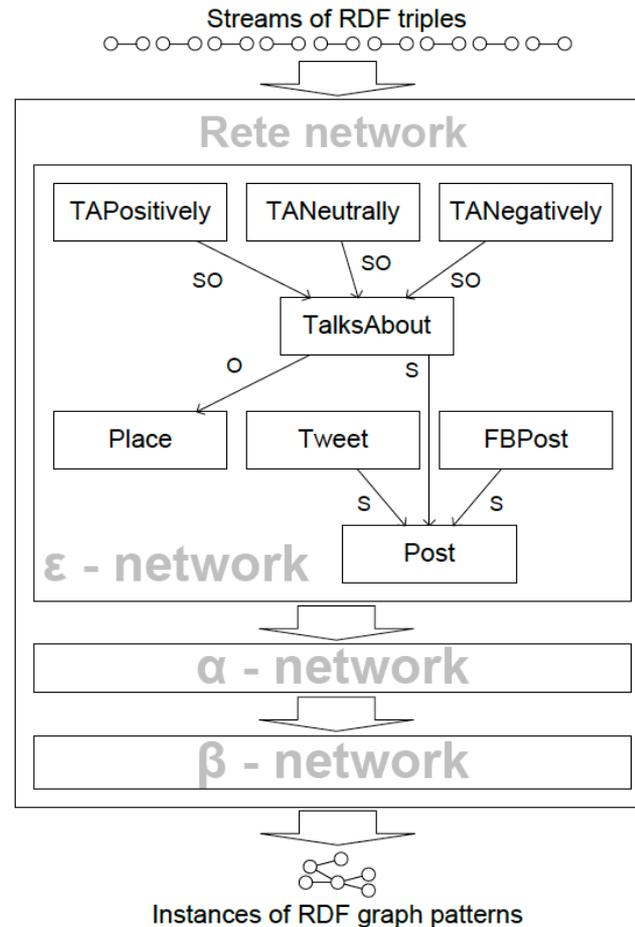
:Place     rdf:type      rdfs:Class.

:TalksAbout rdf:type      rdfs:Property;
           rdfs:domain    :Post;
           rdfs:range     :Place.

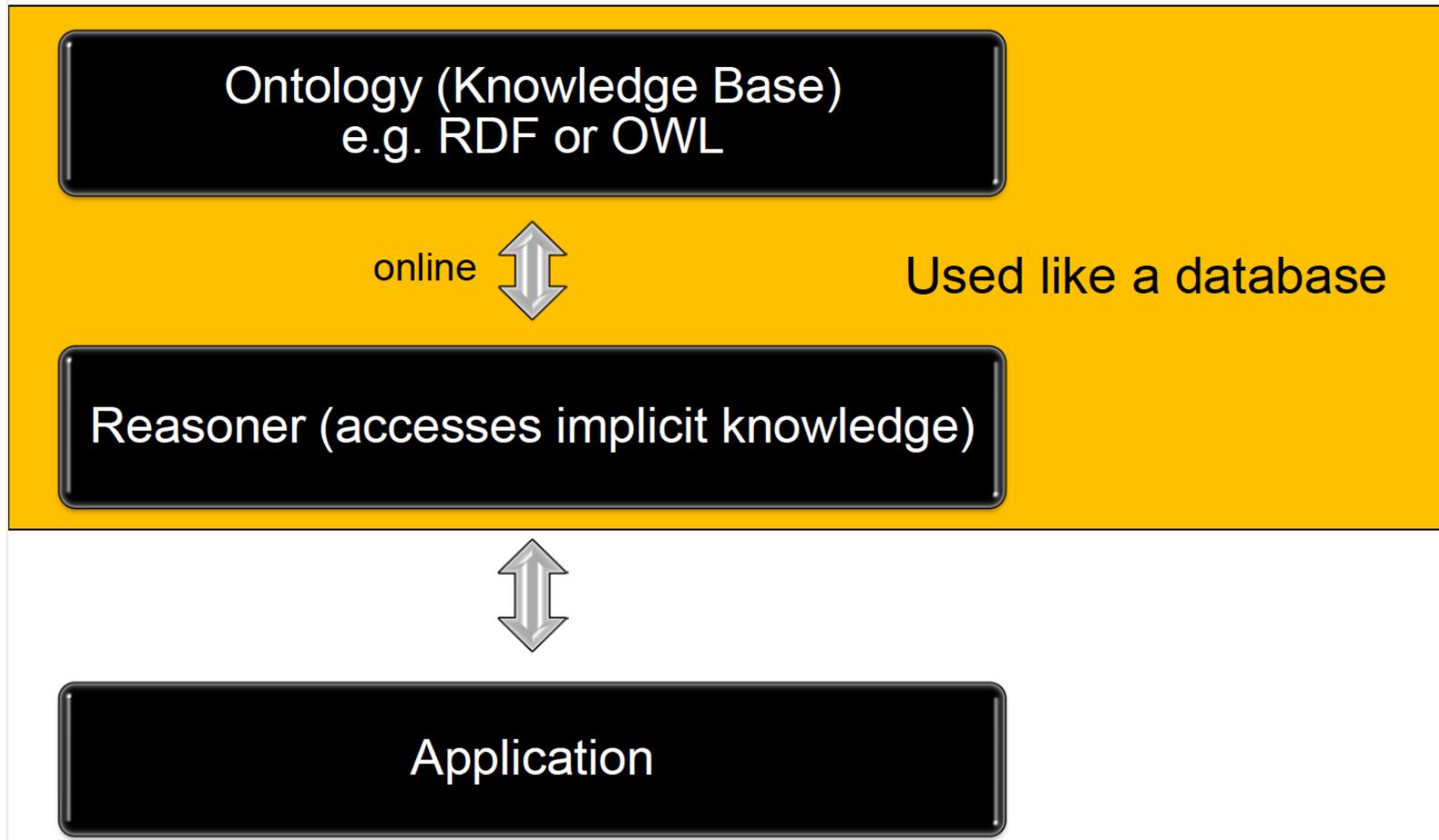
:TAPositively  rdf:type      rdfs:Property;
           rdfs:subPropertyOf :TalksAbout.

:TANeutrally   rdf:type      rdfs:Property;
           rdfs:subPropertyOf :TalksAbout.

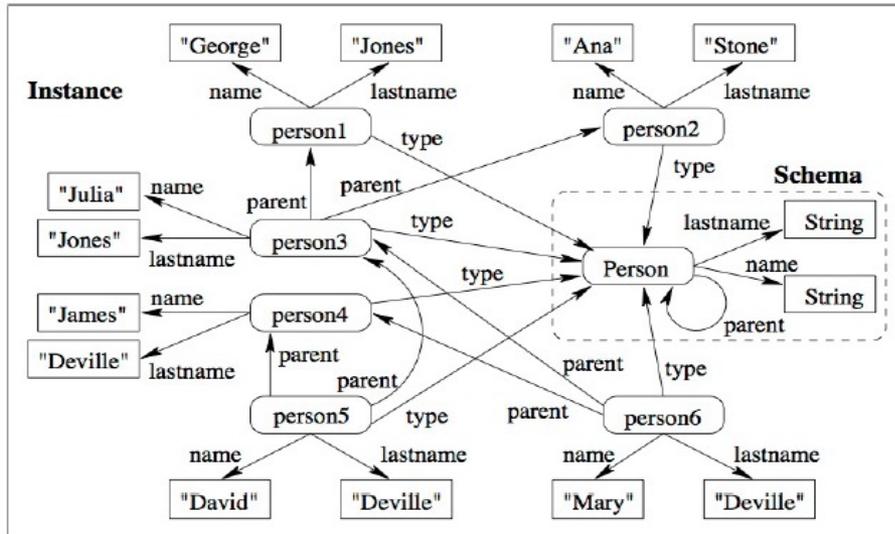
:TANegatively  rdf:type      rdfs:Property;
           rdfs:subPropertyOf :TalksAbout.
    
```



Using implicit knowledge

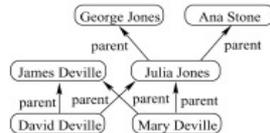


Genealogy – RDF model

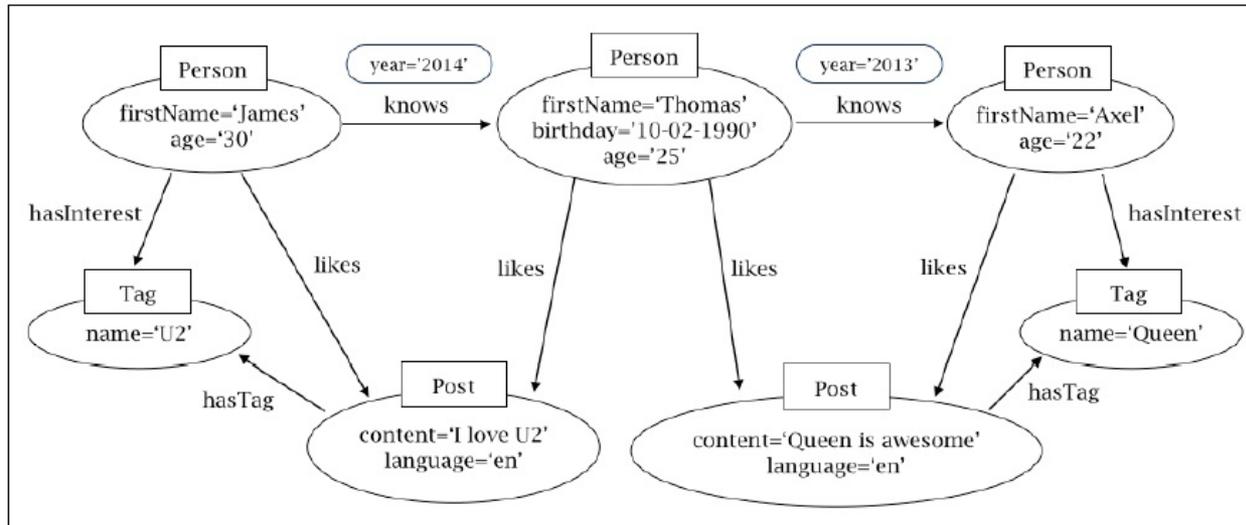


- Not oriented explicitly to model connectivity.
- Originally devised to represent metadata.
- Represents resources and relations between resources.
- No assumption of the application domain.

| NAME | LASTNAME | PERSON | PARENT |
|--------|----------|--------|--------|
| George | Jones | Julia | George |
| Ana | Stone | Julia | Ana |
| Julia | Jones | David | James |
| James | Deville | David | Julia |
| David | Deville | Mary | James |
| Mary | Deville | Mary | Julia |



The property graph data model



Next, we get into GRAPH DB's