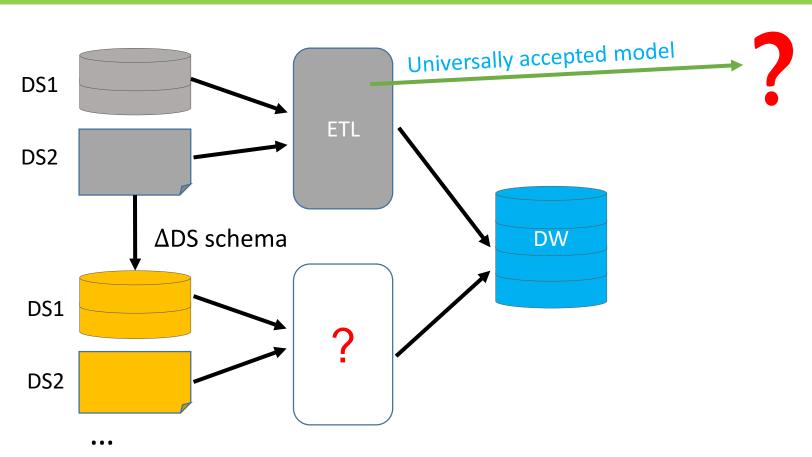
Algorithms and Architecture for Managing Evolving ETL Workflows in a Big Data Environment

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Problem



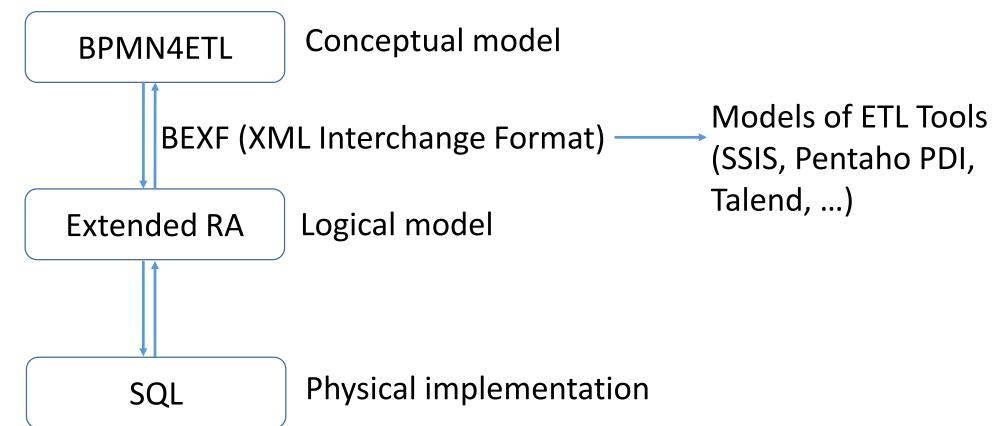
- No agreed-upon model for specifying ETL processes.
- Expensive to manually repair ETL workflows (time, expertise, money).
 Existing ETL tools tacitly assume DSs have static structure not true (Wikipedia had 171 schema versions from Apr 2003 and Nov 2007 [1]).

Objectives

- 1. To propose a methodology for designing ETL processes that will facilitate a smooth transition from gathering user requirements to the actual implementation.
- 2. To develop an Extended Evolving ETL (E3TL) framework to (semi-) automatically repair ETL workflows upon data source changes.

Our approach

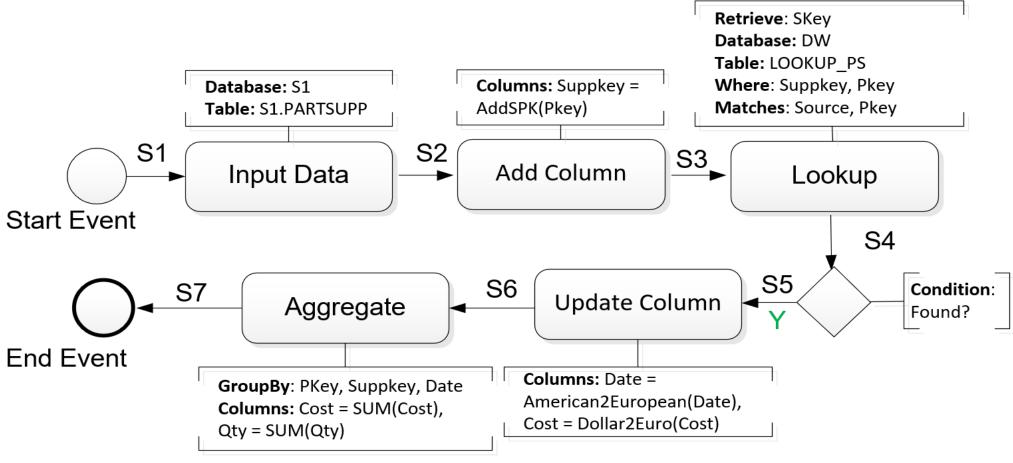
1. ETL Modelling



I. Scenario

	Attributes			
S1.PARTSUPP	Pkey, Qty, Date, Department, Cost			
S2.PARTSUPP	Pkey, SuppKey, Qty, Cost			
DW.PARTSUPP	Pkey, SuppKey, Date, Qty, Cost			

II. BPMN4ETL



III. Extended RA^[2]

Operator	Notation	Operator	Notation	
Selection	$\sigma_C(R)$	Aggregate	$ \mathcal{A}_{A_1,,A_m C_1=F_1(B_1),,C_n=F_n(B_n)}(R) $	
Projection	$\left \pi_{A_{1},,A_{n}}\left(R\right)\right $	Delete	$R \leftarrow R - \sigma_C(R)$	
Cartesian Product	$R_1 imes R_2$	Extend	$\mathcal{E}_{A_1 = Expr_1, \dots, A_n = Expr_n}(R)$	
Union	$R_1 \cup R_2$	Input	$R \leftarrow \mathcal{I}_{A_1, \dots, A_n}(F)$	
Intersection	$R_1 \cap R_2$	Insert	$R \leftarrow R \cup S$ or $R \leftarrow S$	
Difference	$R_1 - R_2$	Lookup	$R \leftarrow \pi_{A_1, \dots A_n} (R_1 \bowtie_C R_2)$	
Join	$R_1 \bowtie_C R_2$	Remove duplicates		
Natural Join	$R_1 * R_2$	Rename	$\rho_{A_1 \leftarrow B_1, \dots, A_n \leftarrow B_n}(R)$ or $\rho_S(R)$	
Left Outer Join	$R_1 \bowtie_C R_2$	Sort	$ au_A(R)$	
Right Outer Join	$R_1\bowtie_C R_2$	Update	$\mathcal{U}_{A_1 \equiv Expr_1, \dots, A_n \equiv Expr_n \mid C}(R)$	
Full Outer Join	$R_1 \bowtie_C R_2$	Update Set	$ \begin{vmatrix} \mathcal{U}_{A_1 = Expr_1, \dots, A_n = Expr_n \mid C}(R) \\ R \leftarrow \mathcal{U}(R)_{A_1 = Expr_1, \dots, A_n = Expr_n \mid C}(S) \end{vmatrix} $	
Semijoin	$R_1 \ltimes_C R_2$			
Division	$R_1 \div R_2$			

```
 \begin{aligned} & \mathsf{Temp1} \leftarrow \mathcal{I}_{\mathsf{Pkey,\ Qty,\ Department,\ Cost}}(\mathsf{SI.PARTSUPP}) & (1) \\ & \mathsf{Temp2} \leftarrow \mathcal{E}_{\mathsf{Suppkey} = \ \mathsf{AddSPK}(\mathsf{Pkey})}(\mathsf{Temp1}) & (2) \\ & \mathsf{Temp3} \leftarrow \pi_{\mathsf{Skey,\ Pkey,\ Suppkey,\ Qty,\ Department,\ Cost}(\mathsf{Temp2} \bowtie_{\mathsf{Pkey} = \ \mathsf{Pkey} \land \ \mathsf{Suppkey} = \ \mathsf{Source}} \ \mathsf{LOOKUP\_PS}) & (3) \\ & \mathsf{Temp4} \leftarrow \mathcal{U}_{\mathsf{Date} = \ \mathsf{American2European}(\mathsf{Date}),\ \mathsf{Cost} = \ \mathsf{Dollar2Euro}(\mathsf{Cost})}(\mathsf{Temp3}) & (4) \end{aligned}
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(5)

IV. BEXF^[3]

Temp5 $\leftarrow A_{Pkey, Suppkey, Date|Cost} = SUM(Cost), Qty = SUM(Qty) (Temp4)$

```
<ETLProcess id="_idProcess" name="Load of DW.PARTSUPP dimension table">
  <ETLTask id="_idInputData" name="Input Data" type="Input Data">
    <Database name="S1"/>
    <Table name="S1.PARTSUPP"/>
    <inputs>
    <inputColumn name="Pkey"/>
    <inputColumn name="Qty"/>
    <inputColumn name="Date"/>
    <inputColumn name="Department"/>
    <inputColumn name="Cost"/>
    </inputs>
    <inRefId>_idS1</inRefId>
    <outRefId>_idS2</outRefId>
  </ETLTask>
  <ETLTask id="_idAggregate" name="Aggregate" type="Aggregate">
    <AggColumn name="Pkey" order="1"/>
    <AggColumn name="Suppkey" order="2"/>
    <AggColumn name="Date" order="3"/>
    <NewColumn name="Cost" function="SUM(Cost)"/>
    <NewColumn name="Qty" function="SUM(Qty)"/>
    <inRefId>_idS6</inRefId>
    <outRefId>_idS7</outRefId>
  </ETLTask>
</ETLProcess>
```

V. Performance Evaluation with TPC-DI benchmark [4].

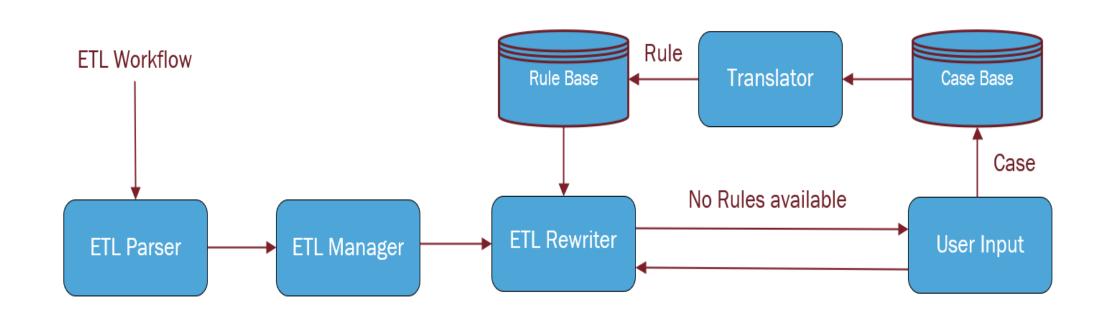
PLSQL: Our approach (BPMN4ETL to Extended RA - SQL)

PDI: BPMN4ETL to Pentaho Data Integration tool

Execution time: hours:minutes:seconds

Exception time: Hears.Himates.seconas								
		Historical	Incremental 1	Incremental 2				
SF-3	PLSQL	00:12:50	00:00:09	00:00:07				
	PDI	11:23:52	00:01:32	00:01:40				
SF-5	PLSQL		00:00:15	00:00:14				
	PDI	20:25:32	00:03:03	00:03:11				
SF-10	PLSQL	02:11:15	00:00:39	00:00:36				
	PDI	25:08:13	00:11:35	00:12:38				

2. ETL Evolution E3TL Framework



ETL Parser: Parses each command of the an ETL workflow. ETL workflow format (RA or SQLs).

ETL Manager: Assesses the impact of the data source change on each command of the ETL workflow and takes these decisions by applying rules stored in a the rule base.

ETL Rewriter: Rewrites the commands in the ETL workflow by applying recommendations from the ETL manager.

Rule Base: Contains distinct rules based on conditions.

User Input: Request the user's input if no rule is available in the rule base to deal with the problem or several solutions are applicable to solve the problem.

Case Base: Repository to store cases.

Translator: Applies algorithms to develop rules from cases.

References

[1] Curino, C. A., Tanca, L., Moon, H. J., & Zaniolo, C. (2008). Schema Evolution in Wikipedia: Toward a Web Information System Benchmark. In: Proc. Of the 10th International Conference on Enterprise Information Systems(ICEIS), Barcelona, Spain.
[2] Awiti, J., Vaisman, A., Zimányi, E.: From Conceptual to Logical ETL Design Using BPMN and Relational Algebra. In: Proc. of the 21st ACM International Conference on Big Data Analytics and Knowledge Discovery, DAWAK 2019. Springer, Linz, Austria (2019), forthcoming

[3] Awiti, J., Zimányi, E.: An XML Interchange Format for ETL Models. In: Proc. of the 23rd European Conference on Advances in Databases and Information Systems, ser. Workshop on BI & Big Data Applications, ADBIS 2019, forthcoming.
[4] Poess, M., Rabl, T., Jacobsen, H. A., & Caufield, B. (2014). TPC-DI: The First Industry Benchmark for Data Integration. Proceedings of the VLDB Endowment, 7(13), 1367-1378.





