

Motivation

- In a distributed business scenario, where multiple partner companies/organizations cooperate towards a common goal, traditional BI systems are no longer sufficient to maximize the effectiveness of decision making processes
- Two main requirements arise:
 - Cross-organization monitoring and decision making Accessing local information is no more enough, users need to transparently and uniformly access information scattered across several heterogeneous BI platforms
 - Pervasive and personalized access to information Users require that information can be easily and timely accessed through devices with different computation and visualization capabilities, and with sophisticated and customizable presentations

1st European Business Intelligence Summer School (eBISS 2011)

Collaboration [Wiki]

- Collaboration is working together to achieve a goal
 - It is a recursive process where two or more people or organizations work together to realize shared goals —this is more than the intersection of common goals, but a deep, collective, determination to reach an identical objective— by sharing knowledge, learning and building consensus
- Most collaborations require leadership
- Teams that work collaboratively can obtain greater resources, recognition and reward when facing competition for finite resources

5

Collaborative Bl Collaboration is seen today by companies as one of the major means for increasing flexibility and innovating so as to survive in today uncertain and changing market Companies need strategic information about the outer world, for instance about trading partners and related business areas Users need to access information anywhere it can be found, by locating it through a semantic process and performing integration on the fly This is particularly relevant in inter-business collaborative contexts where companies organize and coordinate themselves to share opportunities, respecting their own autonomy and heterogeneity but pursuing a common goal 7 1st European Business Intelligence Summer School (eBISS 2011)

But...

- ...most information systems were devised for individual companies and for operating on internal information, and they give limited support to intercompany cooperation
- ...traditional BI applications are aimed at serving individual companies, and they cannot operate over networks of companies characterized by an organizational, lexical, and semantic heterogeneity



need for innovative approaches and architectures

Data warehouse integration

- Data warehouse integration is an enabling technique for collaborative BI; it provides a broader base for decision-support and knowledge discovery than each single data warehouse could offer
 - Large corporations integrate their separately-developed departmental data warehouses
 - Newly merged companies integrate their data warehouses into a central data warehouse
 - Autonomous but related organizations join together their data warehouses to enforce the decision making process

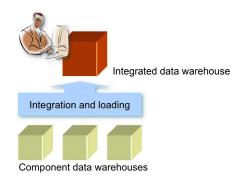


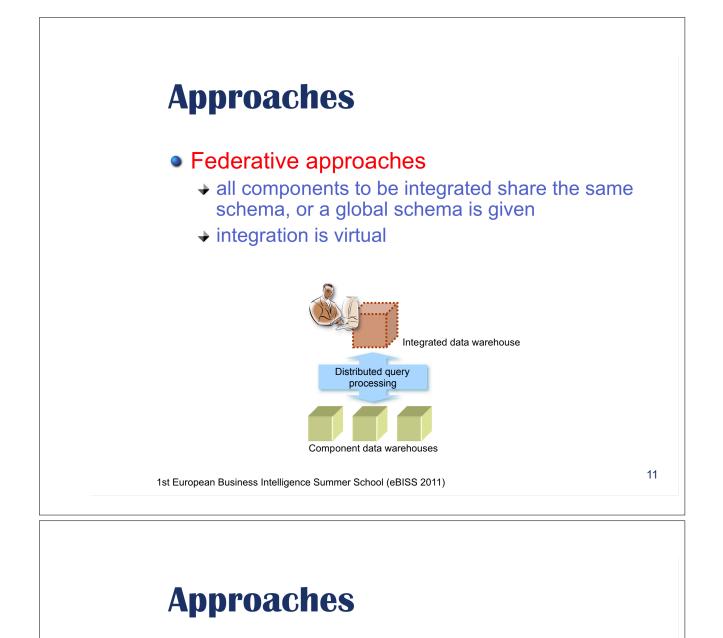
1st European Business Intelligence Summer School (eBISS 2011)

Approaches

Warehousing approaches

- all components to be integrated share the same schema, or a global schema is given

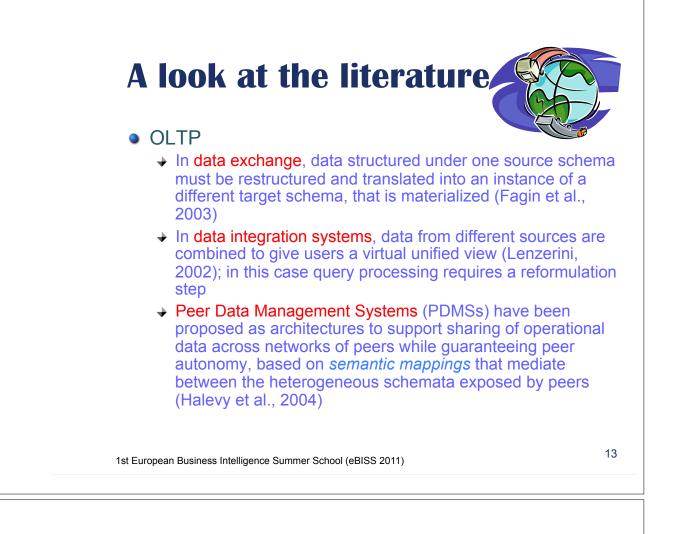




P2P approaches

- they do not rely on a global schema to integrate the component data warehouses
- necessary in contexts where the different parties have a common interest in collaborating while fully preserving their autonomy and their view of business
- each peer can formulate queries also involving the other peers, typically based on a set of mappings that establish semantic relationships between the peers' schemata

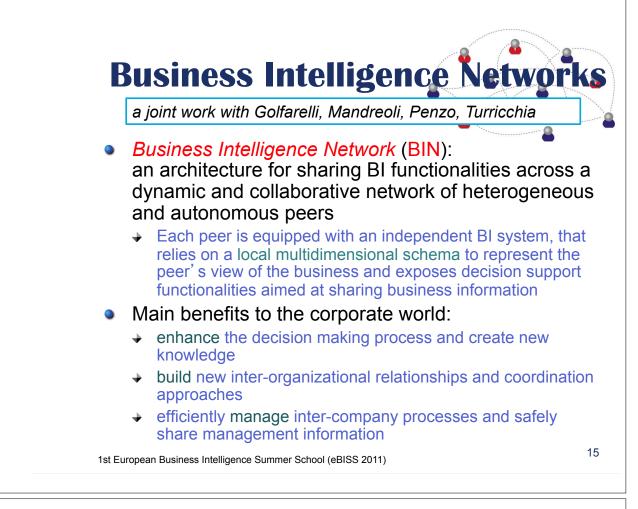


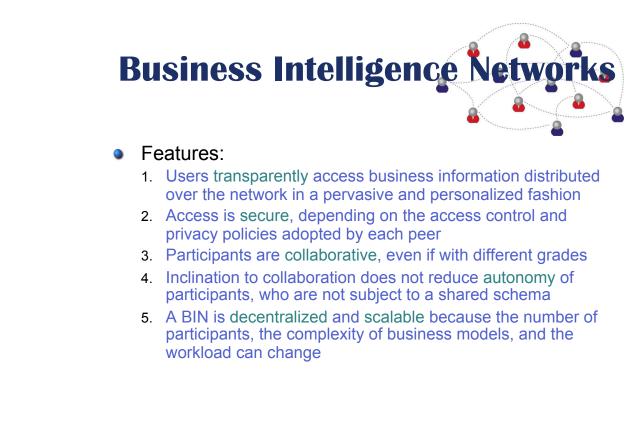


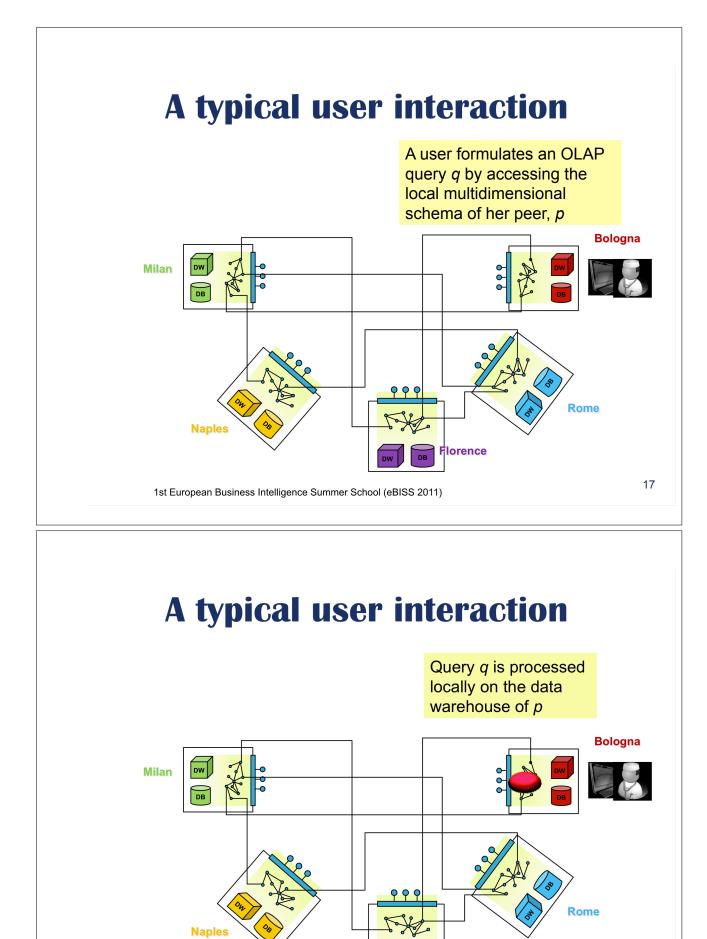
A look at the literature



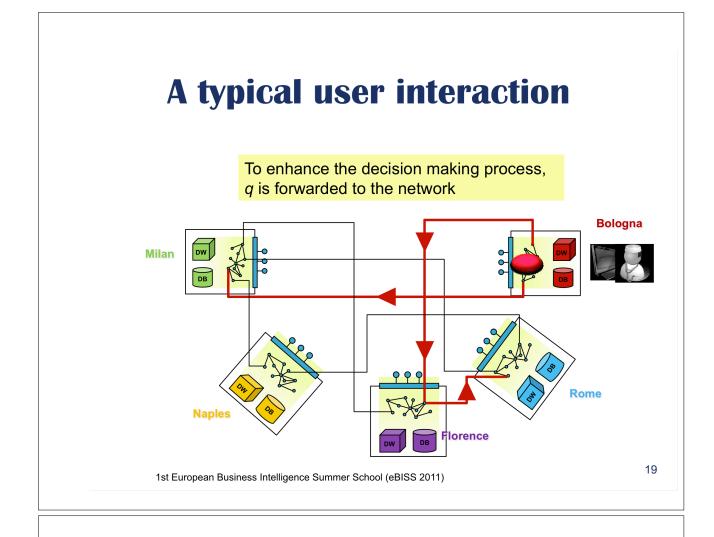
- In this context, problems related to data heterogeneity are usually solved by ETL processes that load data into a single multidimensional repository, but some works are specifically focused on strategies for data warehouse integration and federation:
 - Check of *coherence*, *soundness*, and *consistency* of two dimensions (Torlone, 2008)
 - CubeStar (Albrecht & Lehner, 1998); Skalla (Akinde et al, 2003); Multi data warehouse (Berger & Schrefl, 2006)
 - Discovery of inter-hierarchy mappings (Banek et al., 2008)
 - P2P warehousing of XML content (Abiteboul, 2003)
 - Multidimensional P2P network and OLAP query rewriting (Vaisman, 2009)
 - Semantic layer to enable communication among different components (Kehlenbeck & Breitner, 2009)



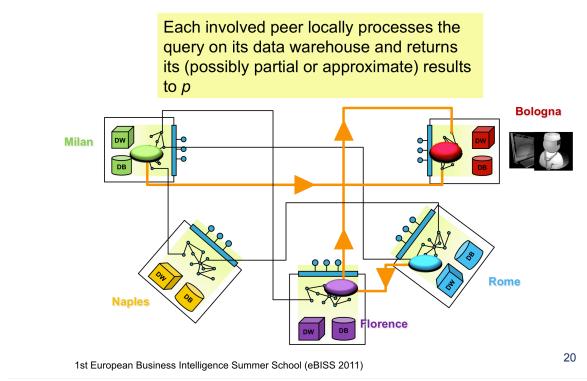


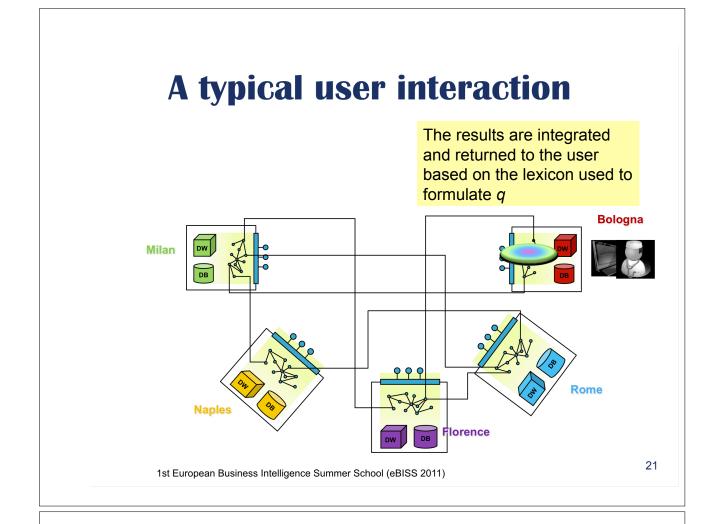






A typical user interaction

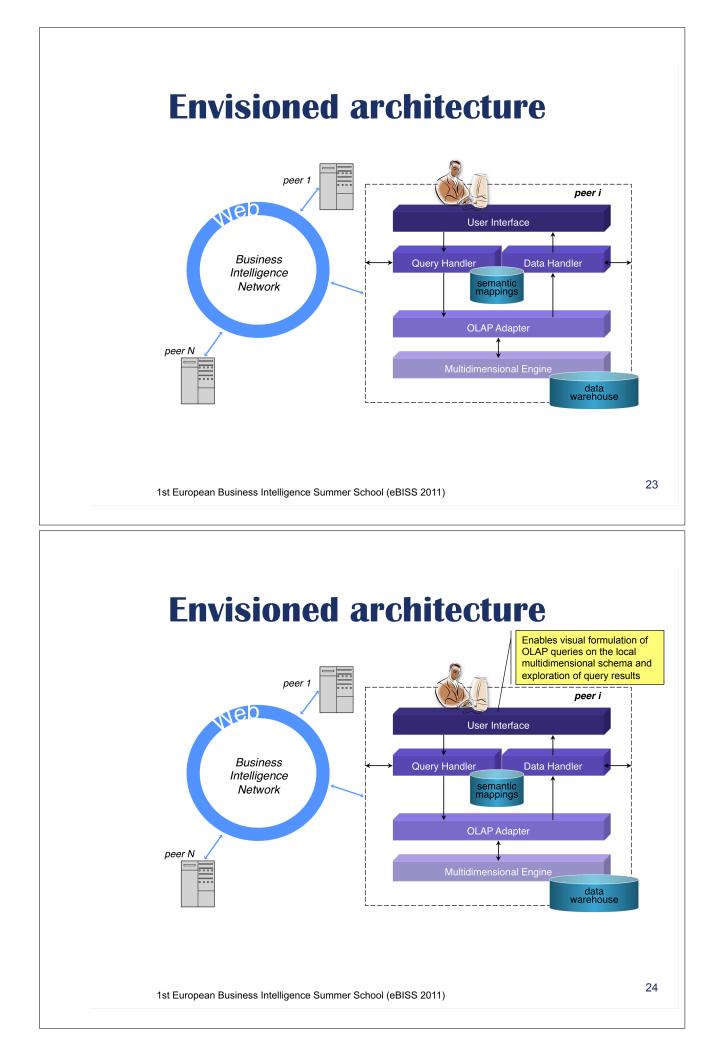


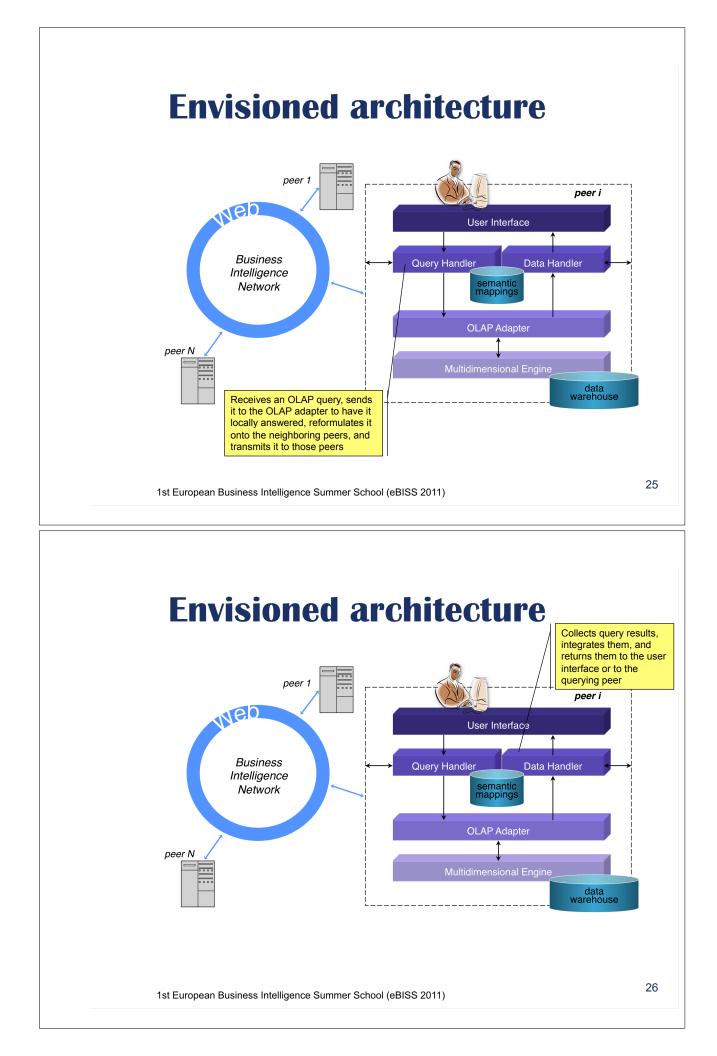


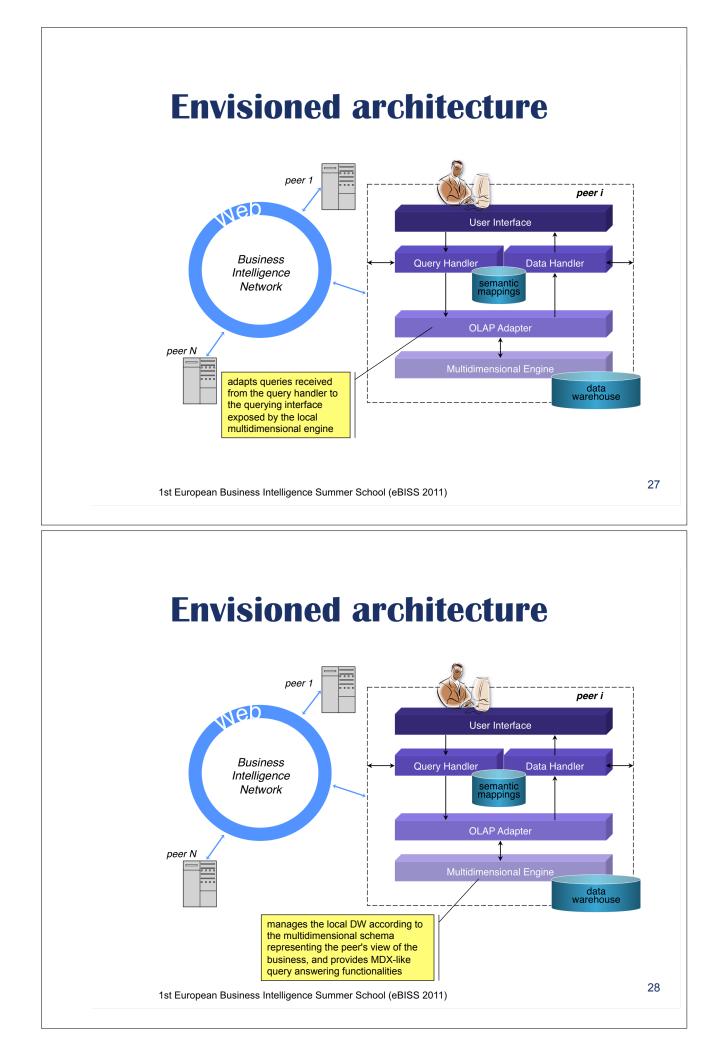
How to deal with heterogeneity?

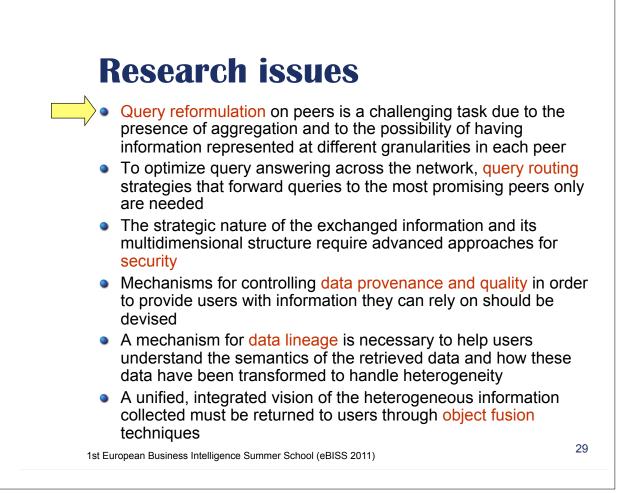
 Before a query issued on a peer can be forwarded to the network it must be first reformulated according to the multidimensional schemata of the destination peers





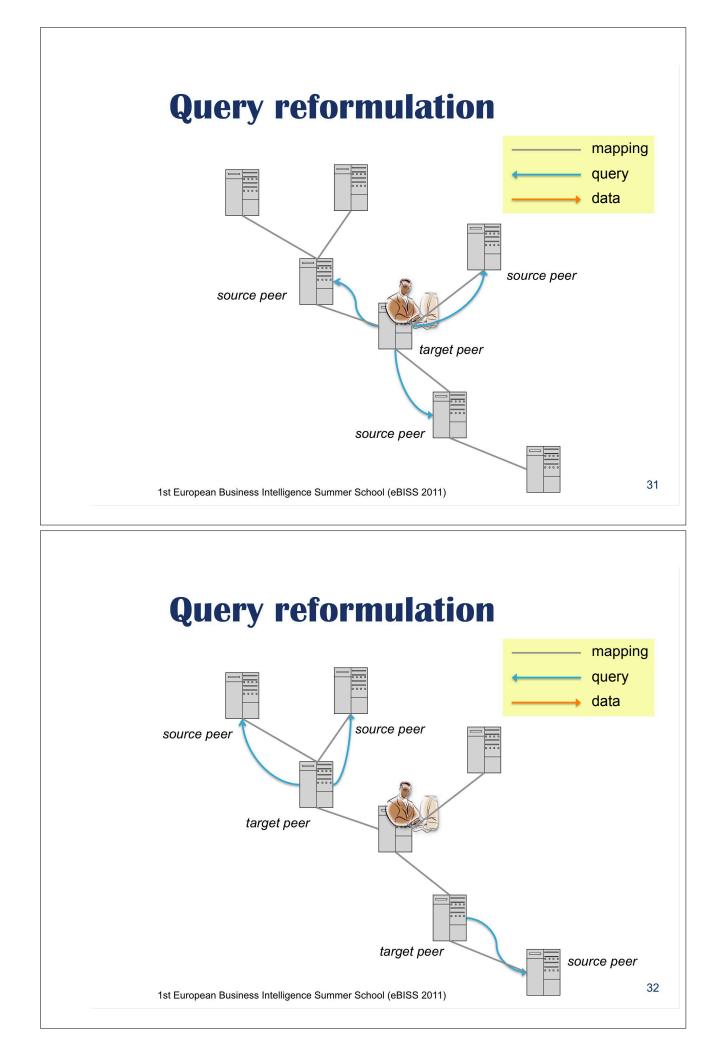


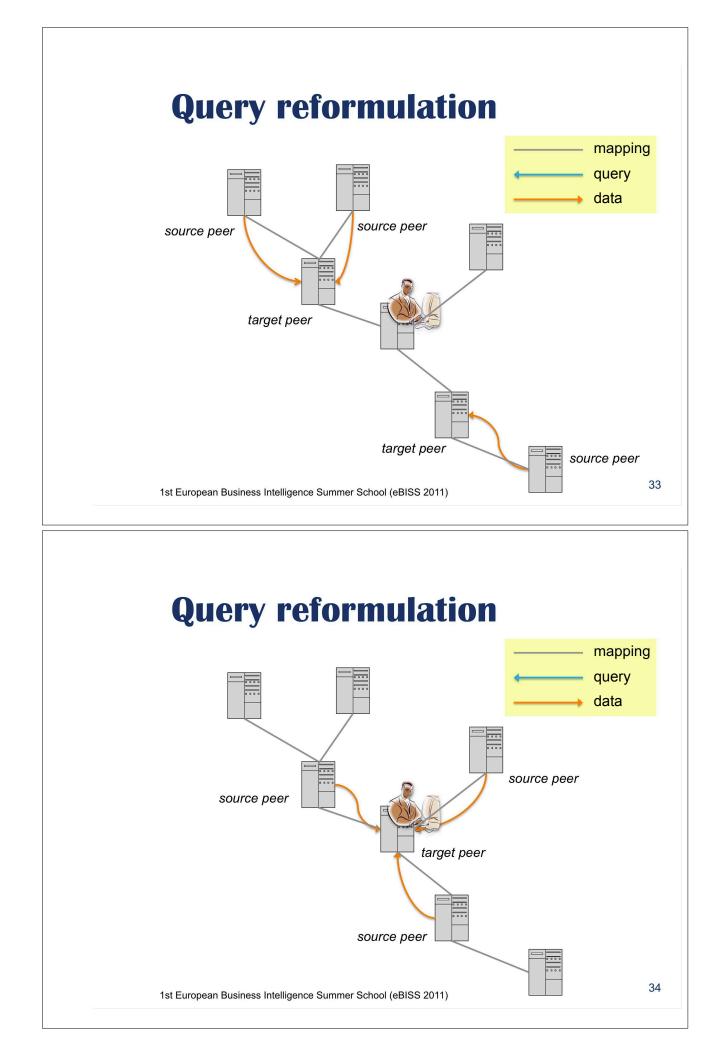




Query reformulation

- Like in PDMSs, query reformulation in a BIN is based on *semantic mappings* that mediate between the different multidimensional schemata exposed by two peers
- Direct mappings cannot be realistically defined for all the possible couples of peers; so, a query issued on p is forwarded to the network by first sending it to the immediate neighbors of p, then to their immediate neighbors, and so on
 - In this way, the query undergoes a chain of reformulations along the peers it reaches, and results are collected from any peer that is connected to p through a path of semantic mappings

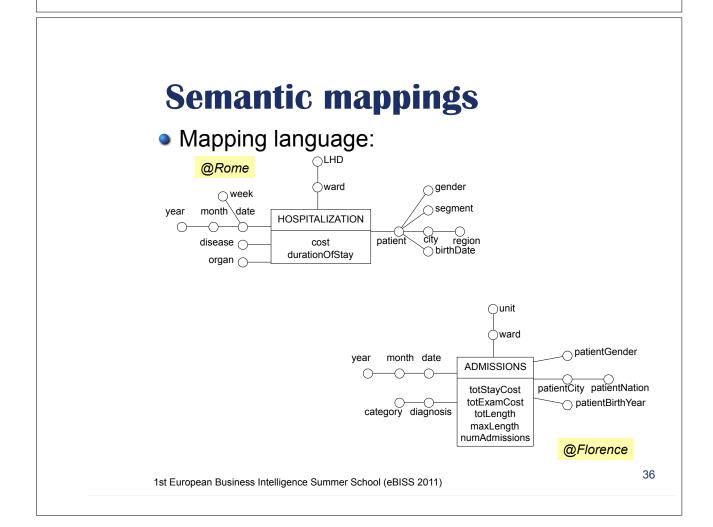


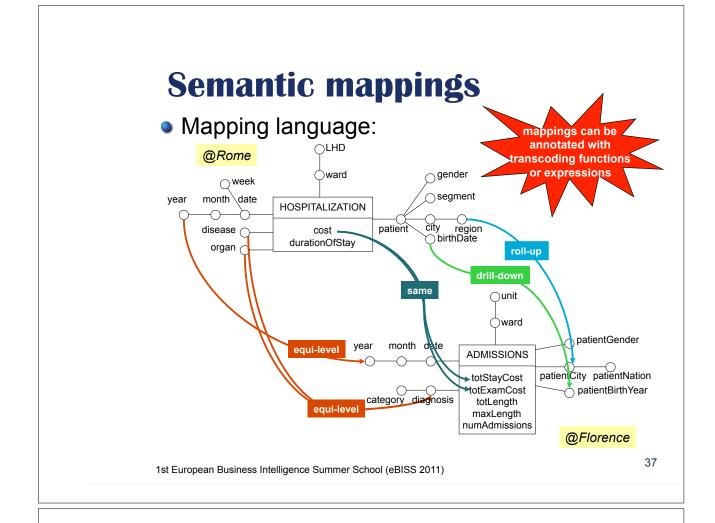




- They describe how the concepts in the multidimensional schema of one peer map onto those of another peer
- Requirements:
 - Handling the asymmetry between dimensions and measures
 - Specifying the relationship between two attributes of different multidimensional schemata in terms of their granularity
 - Considering aggregation operators to avoid the risk of inconsistent query reformulations
 - Expressing also mappings at the instance level to transcode data







Example: mappings

- $\omega 1 \quad \langle \text{ cost,sum } \rangle \text{ same } \{ \text{ totStayCost, totExamCost } \}$
- $\omega^2 \quad \langle \text{ cost,avg } \rangle \text{ same } \{ \text{ totStayCost, totExamCost } \}$
- $\omega 3 \quad \langle \text{ durationOfStay,sum } \rangle \text{ same } \{ \text{ totLength } \}$
- $\omega 4 \quad \langle \text{ durationOfStay,avg } \rangle \text{ same } \{ \text{ totLength } \}$
- $\omega 5 \quad \langle \text{ durationOfStay,max} \rangle \text{ same } \{ \text{ maxLength } \}$
- $\omega 6 \qquad \{ LHD \} \text{ roll-up } \{ \text{ unit } \}$
- $\omega 7 \quad \{ \text{ ward } \} \text{ equi-level } \{ \text{ ward } \}$
- $\omega 8$ { year } equi-level { year }
- $\omega 9 \quad \{ \text{ month } \} \text{ equi-level } \{ \text{ month } \}$
- $\omega 10 \quad \{ date \} equi-level \{ date \}$
- $\omega 11 \quad \{ week \} roll-up \{ date \} \}$
- $\omega 12 \quad \{ disease, organ \} equi-level \{ diagnosis \}$
- ω13 { disease } drill-down { category }
- $\omega 14$ { patient } drill-down { patientGender, patientCity, patientBirthYear }
- $\omega 15 \quad \{ \text{ gender } \} \text{ equi-level } \{ \text{ patientGender } \}$
- $\omega 16 \quad \{ \text{ segment } \} \text{ related } \{ \text{ patientGender, patientCity, patientBirthYear } \}$
- ω17 { birthDate } drill-down { patientBirthYear }
- $\omega 18$ { city } equi-level { patientCity }
- $\omega 19 \quad \{ region \} roll-up \{ patientCity \}$

Example: transcodings

	$\omega 1$	<pre>cost = totStayCost+totExamCost, segment in { 'NH','EU' }</pre>	
	$\omega 2$	<pre>cost = totStayCost+totExamCost, segment in { 'NH','EU' }</pre>	
	$\omega 3$	durationOfStay = totLength, segment in { 'NH', 'EU' }	
	$\omega 4$	durationOfStay = totLength, segment in { 'NH', 'EU' }	
	$\omega 5$	durationOfStay = maxLength, segment in { 'NH', 'EU' }	
	$\omega 6$	LHD = 'LHD39 - Florence'	
	$\omega 7$	ward = ward	
	$\omega 8$	year = year	
	$\omega 9$	month = month	
	$\omega 10$	date = date	
	$\omega 11$	week = weekOf(date)	
	$\omega 12$	disease = substring(diagnosis, 1, 40), organ = substring(diagnosis, 41, 80))
	$\omega 13$	categoryOf(disease) = category	
	$\omega 14$	_	
	$\omega 15$	gender = completeGender(patientGender)	
	$\omega 16$	_	
	$\omega 17$	yearOf(birthDate) = patientBirthYear	
	$\omega 18$	city = patientCity	
	$\omega 19$	region = regionOf(patientCity)	
1st European Business Intelligence Summer School (eBISS 2011) 39			39

Mapping accuracy

- A mapping is exact when it is either an equilevel or a roll-up mapping and it has an associated transcoding, or it is a same mapping
- A mapping is loose when it is either a drilldown or a related mapping
- An attribute mapping is approximate when it has no associated transcoding

Mapping accuracy

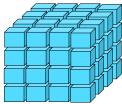
- The accuracy of a reformulation of q on peer s depends on the accuracy of the mappings involved
 - → When (i) for each attribute mentioned in *q* there is an exact mapping from *s*, and (ii) for each metric required by q there is a same mapping from *s*, there exists a *compatible* reformulation of *q* on *s*, i.e., one that fully preserves the semantics of *q*
 - when a compatible reformulation is used, the results returned by *s* do exactly match with *q* so they can be seamlessly integrated with those returned by *t*
 - In all the other cases, the results returned by s match q with some approximation
 - value mismatch
 - granularity mismatch
 - no reformulation

1st European Business Intelligence Summer School (eBISS 2011)

41

Expressiveness

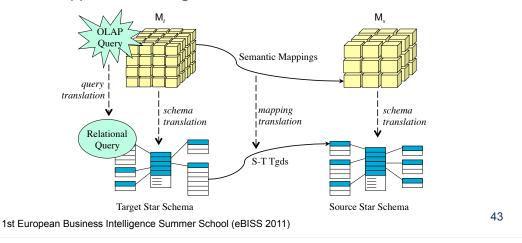
- BIN query:
 - → a group-by clause
 - → an (optional) selection conjunctive predicate
 - a numerical expression involving measures to be computed
 - an aggregation operator to be used for each measure



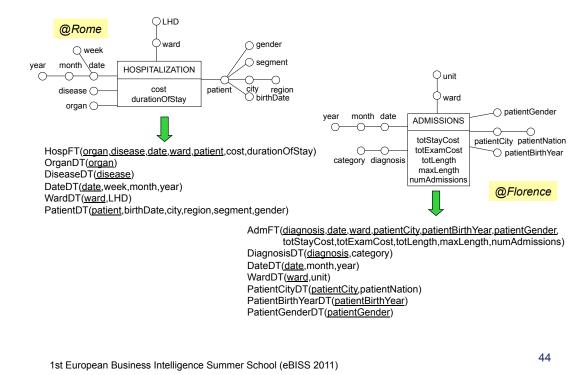
Query reformulation

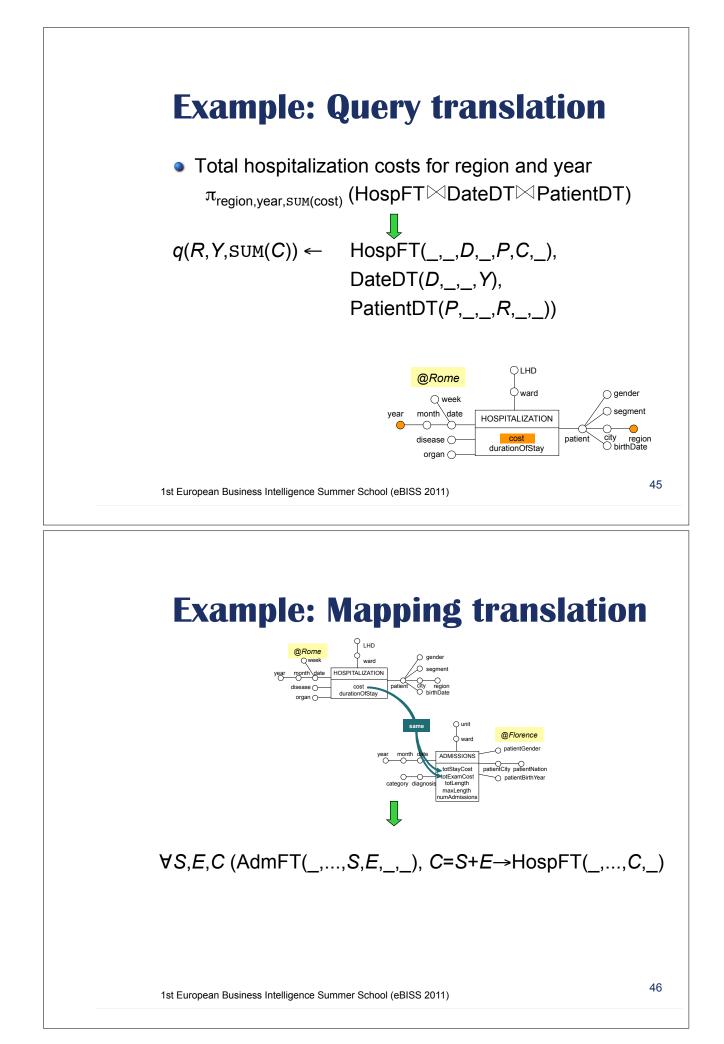
Framework:

→ To translate semantic mappings we use a logical formalism called source-to-target tuple generating dependencies (ten Cate & Kolaitis, 2010), asserting that if a pattern of facts appears in the source, then another pattern of facts must appear in the target



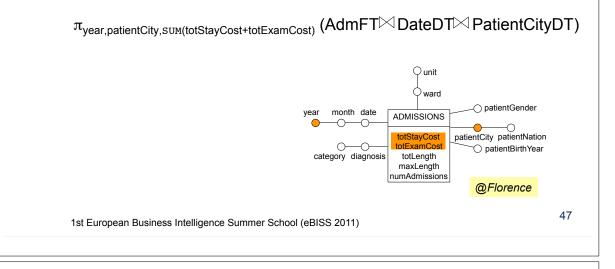
Example: Schema translation





Example: Reformulation

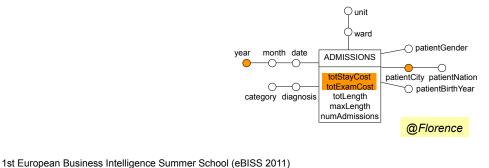
• The group-by is reformulated using the roll-up mapping from *region* to *patientCity* and the equilevel mapping from *year* to *year*, while measure *cost* is derived using the same mapping



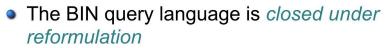
Example: Reformulation

 In presence of a transcoding from *region* to *patientCity*, a compatible reformulation can be done

π_{year,}regionOf(patientCity),suм(totStayCost+totExamCost)</sub> (AdmFT⊠DateDT⊠ PatientCityDT)



Theoretical results



- Our query reformulation algorithm can be used by each peer in a BIN to implement chains of reformulations. In this way, any query formulated over a peer schema can be safely distributed across the network, and answers can come from any other peer in the network which is connected to the queried peer through a chain of semantic mappings
- Our reformulation algorithm is proved to be sound and complete with respect to the semantics of query answering, that in data sharing settings is usually given in terms of certain answers

1st European Business Intelligence Summer School (eBISS 2011)

Implementation issues

- How to bridge the language gap between the query handler and the local multidimensional engine?
 - A BIN query cannot be directly executed on the peer local multidimensional engine
 - Intra-peer reformulation must deal with the presence of transcodings in the query group-by set, and must properly manage non-distributive aggregation operators

49

Implementation issues

How to share transcodings among peers?

- Public transcodings are standard database functions that are shared by all peers
- Protected transcodings are owned by a peer, that will make them available to its neighboring peers by attaching them to query messages
 - If protected transcodings are expressed as procedures, a shared programming language must be available in the BIN
 - Otherwise, transcodings can be expressed as look-up tables to be applied by a relational engine; in this case, an obvious drawback is the quantity of information to be transmitted over the network

1st European Business Intelligence Summer School (eBISS 2011)

51

Summary

- We have surveyed the basic approaches to collaborative BI
- We have outlined a peer-to-peer architecture for supporting distributed and collaborative decision-making scenarios
- We have shown how an OLAP query formulated on one peer can be reformulated on a different peer, based on a set of interpeer semantic mappings



Open issues

 Devising multidimensional-aware object fusion techniques for integrating data returned by different peers



- Finding smart algorithms for routing queries to the most "promising" peers in the BI network
- Designing smart user interfaces for emphasizing the differences and relationships between the returned data
- Using preferences to rank the returned data depending on how compliant they are with the original local query
- Studying mechanisms for controlling data provenance and quality to provide users with reliable information
- Devising advanced approaches for security, especially data sharing policies that depend on the degree of trust between participants

1st European Business Intelligence Summer School (eBISS 2011)

53

Related readings

- Abiteboul, S. Managing an XML warehouse in a P2P context. In *Proc. CAISE*, 2003
- Albrecht, J., & Lehner, W. On-line analytical processing in distributed data warehouses. In Proc. IDEAS, 1998
- Akinde, M.O., Bohlen, M.H., Johnson, T., Lakshmanan, L.V.S., & Srivastava, D. Efficient OLAP query processing in distributed data warehouses. *Inf. Syst.* 28(1-2), 2003
- Banek, M., Vrdoljak, V., Min Tjoa, A., & Skocir, Z. Automated integration of heterogeneous data warehouse schemata. *IJDWM*, 4(4), 2008
- Berger, S., & Schrefl, M. From federated databases to a federated data warehouse system. In Proc. HICSS, 2008
- Chang, K.C., Garcia-Molina, H.: Mind your vocabulary: Query mapping across heterogeneous information sources. In Proc. SIGMOD, 1999
- Dubois, D., & Prade, H. On the use of aggregation operations in information fusion processes. International Journal on Fuzzy Sets and Systems, 142(1), 2004
- Georgiadis, P., Kapantaidakis, I., Christophides, V., Nguer, E. M., & Spyratos, N. Efficient rewriting algorithms for preference queries. In *Proc. ICDE*, 2008
- Golfarelli, M., Mandreoli, F., Penzo, W., Rizzi, S., & Turricchia, E. BIN: Business intelligence networks. In Business Intelligence Applications and the Web: Models, Systems and Technologies, IGI Global, 2011 (to appear).
- Golfarelli, M., Mandreoli, F., Penzo, W., Rizzi, S., & Turricchia, E. OLAP Query Reformulation in Peer-to-Peer Data Warehousing. *Information Systems*, 2011 (to appear)
- Halevy, A. Y., Ives, Z. G., Madhavan, J., Mork, P., Suciu, D., & Tatarinov, I. The Piazza Peer Data Management System. IEEE TKDE, 16(7), 2004
- Hoang, T. A. D., & Binh Nguyen, T. State of the art and emerging rule-driven perspectives towards servicebased business process interoperability. In Proc. Int. Conf. on Computing and Communication Technologie, 2009

Related readings

- Kalnis, P., Siong Ng, W., Chin Ooi, B., Papadias, D., & Tan, K.-L. An adaptive peer-to-peer network for distributed caching of OLAP results. In Proc. SIGMOD Conference, 2002
- Kehlenbeck, M., & Breitner, M. H. Ontology-based exchange and immediate application of business calculation definitions for online analytical processing. In *Proc. DAWAK*, 2009
- Kießling, W. Foundations of preferences in database systems. In Proc. VLDB, 2002
- Mandreoli, F., Martoglia, R., Penzo, W., & Sassatelli S. SRI: exploiting semantic information for effective query routing in a PDMS. In Proc. ACM Int. Workshop on Web Information and Data Management, 2006
- Mecca, G., Papotti, P., & Raunich, S. Core Schema Mappings. In Proc. ACM SIGMOD Int. Conf. on Management of Data, 2009
- Papakonstantinou, Y., Abiteboul, S., & Garcia-Molina, H. Object fusion in mediator systems. In Proc. VLDB, 1996
- Schneider, M. Integrated vision of federated data warehouses. In Proc. DISWEB, 2006
- Sung, S., Liu, Y., Xiong, H., & Ng, P. Privacy preservation for data cubes. *Knowledge and Information* Systems, 9(1), 2006
- Tatarinov, I. & Halevy, A.Y. Efficient Query Reformulation in Peer-Data Management Systems. In Proc. ACM SIGMOD Int. Conf. on Management of Data, 2004
- ten Cate, B. & Kolaitis, P. G. Structural characterizations of schema-mapping languages. Comm. ACM, 53 (1), 2010
- Torlone, R. Two approaches to the integration of heterogeneous data warehouses. Int. Journ. on Distributed and Parallel Databases, 23(1), 2008
- Vaisman, A., Espil, M.M., & Paradela, M. P2P OLAP: Data model, implementation and case study. *Information Systems*, 34(2), 2009

1st European Business Intelligence Summer School (eBISS 2011)

55

Thank you for you attention

Questions?